

1 restart;maple_mode(1);cas_setup(0,0,0,1,0,1e-10,10,[1,50,0,25],0,0,0); #radians,pas de cmplx, pas de Sqrt

2 time(p:=nextprime(10^50));

0.005

3 time(isprime(p));

0.07

4 time(is_pseudoprime(p));

0.0026

5 le temps de isprime est trop long pour avoir ete utilise pour trouver le nombre premier avec nextprime. La doc confirme bien que nextprime donne un pseudopremier.

6 i:=20;n:=nextprime(10^i)*nextprime(20^i);

(20, 10485760000000000004089455500000000000000003549)

7 ifactor(n);

Evaluation time: 1.64

100000000000000000039 · 10485760000000000000000091

8 a partir de i=30 ca ne repond plus?

9 p:=nextprime(=nextprime(10^30)*nextprime(20^30));

1073741824000000000000000061203284109000000000000000000008409

10 ifactor(p-1); #ca marche souvent

$2^3 \cdot 3 \cdot 389 \cdot 733 \cdot 156904374622257604823879982847602392900751802349981470895277241$

11 pari();

All PARI functions are now defined with the pari_ prefix.
 PARI functions are also defined without prefix except:
 abs acos acosh arg asin asinh atan atanh binomial bitand bitor bitxor ceil charpoly concat conj content cos cosh divis
 Note that p-adic numbers must have O argument quoted e.g. 905/7+O('7^3')
 Type ?pari for short help
 Inside xcas, try Help->Manuals->PARI for HTML help

12 pari_isprime(p,1);

Evaluation time: 9.66

2	7	1
3	2	1
389	2	1
733	2	1
156904374622257604823879982847602392900751802349981470895277241	2	56467
		65530849258879746208

13 Les grands nombres premiers apparaissant dans le certificat sont eux m^eme certifi'es, ce qui explique les crochets emboit'es.

14 On a cree n avec de petits facteurs. on arrive a le factoriser

```

15 i:=0;l:=[];
   while i<1 do
   n:=1;for j from 1 to 10 do n:=n*(1+rand(10000)) od;
   if is_pseudoprime(n+1)>0 then p:=n+1; i++; fi;
   od;
   ( 0, [], 1 )
16 Divn:=transpose(factors(p-1)[2]);
   2 3 5 7 11 13 17 19 127 131 167 977 1091 4211 5009
   10 8 2 2 1 1 1 1 1 1 1 1 1 1 1
17 divn:=Divn[1];
   [ 2 3 5 7 11 13 17 19 127 131 167 977 1091 4211 5009 ]
18 expn:=Divn[2];
   [ 10 8 2 2 1 1 1 1 1 1 1 1 1 1 1 ]
19 On cherche maintenant le certificat. Pour chaque premier divisant p-1 on trouve
   un element d'ordre la puissance maximale de ce diviseur qui divise p-1. On
   utilisera factors et powmod
20 certifs:=[];
   for i in divn do
   a:=2;
   while powermod(a,(p-1)/i,p)=1 do a:=1+rand(p-1) od;
   certifs:=[op(certifs),a];
   od;
21 g:=1;
   1
22 for i from 1 to dim(divn) do g:=mods(g*powmod(certifs[i],(p-1)/(divn[i]^expn[i]),p),p) od;
   11700374211966552891421579217907138
23 # seq(powmod(g,i,p),i=divisors(p-1));# NON! trop de diviseurs!
   Syntax compatibility mode maple
   Parse error line 2 at /
   undef
24 p:=nextprime(10^13);i:=0;l:=[];
   ( 100000000000037 , 0, [] )
25 while i<10 do
   if is_pseudoprime(4*p+1)>0 then l:=[op(l),4*p+1]; i++; fi;
   p:=nextprime(p);
   od;l;
   Done, [ 40000000000733 40000000002629 40000000011053 40000000016309 40000000017917 40000000019117 ]
26 p:=nextprime(10^13+100);
   10000000000129
27 g:=pari_znprimroot(p);
   7 % 10000000000129
28 pari_znlog(3,Mod(g,p)); #il reussit
   2552475531946
29 ifactor(p-1);
   2^7 · 3 · 98867 · 263401
30 p-1 a de petits facteurs premiers, on peut donc resoudre le probleme
   du log discret. Dans le cas suivant ca ne marche plus.
31 p:=l[3];
   40000000011053
32 g:=pari_znprimroot(p);

```

```
33 #pari_znlog(3,Mod(g,p)); #depasse la memoire max de pari
Syntax compatibility mode maple
Parse error line 2 at /
undef Menu
```

```
34 -----EXERCICE-----
```

```
35 p:=1[3];
40000000011053 Menu
```

```
36 b,e pas trop petits par rapport a p, pour que le modulo p melange vraiment.
```

```
37 message:="ABCD A, Remarquez que des lettres identiques ne sont pas codees de la meme maniere!";asc(message)
ABCD A, Remarquez que des lettres identiques ne sont pas codees de la meme maniere! , [ 65 66 67 68
```

```
38 asc(messsage)-[seq(65,i=1..length(message))];
[-65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65 -65
```

```
39 b:=12345;e:=54321; igcd(e,p-1); B:=powmod(b,e,p);
( 12345, 54321, 1, 33473253573817 ) Menu
```

```
40 l:=asc(message);k:=seq(rand(50),i=1..length(message));
[ 65 66 67 68 65 44 32 82 101 109 97 114 113 117 101 122 32 113 117 101 32
```

```
41 C:=[seq([(b^k[i]) mod p, (l[i]*B^k[i]) mod p],i=1..length(message))];#le message crypte
```

1112550120652	1712465384039
9285222567600	20568860851465
10304518854412	5530469661717
12345	31144114074766
15764341295284	22853646682761
26072565355155	24716418195292
20640365499672	31403166870846
18276441706461	23740747885750
10304518854412	24131591326474
23119392340457	3004317636486
18689839150260	8253956815539
18689839150260	8253956815539
9008468196307	31350325693608
1881365963625	196945846564
25099884345130	15950384136745
18276441706461	23740747885750
1881365963625	36640389256496
10304518854412	37158880861940
5312023043230	11247923060337
18072155013312	7462318873458
20753577694059	3027232228382
15764341295284	29996329018557
3060797962643	38087767443751
3060797962643	38087767443751

42 ee:=p-1-e;#c'est -e [p-1] qui nous interesse

39999999956731

43 L'astuce est que $(B^k) = ((b^e)^k)$ qui vaut par commutativit'e des la composition des application $B:x \rightarrow x^k$ et $A:x \rightarrow x^e$ $((b^k)^e)$.
Le principe a retenir, est qu'en crypto on a juste besoin d'operations A,B qui commutent, et que la donnee de A de donne pas $A^{(-1)}$

44 decrypt:=[seq(irem(powmod(C[i][1],ee,p)*C[i][2],p),i=1..length(message))];

65 66 67 68 65 44 32 82 101 109 97 114 113 117 101 122 32 113 117 101 32

45 char(decrypt);

ABCD A, Remarquez que des lettres identiques ne sont pas codees de la meme maniere!

46 -----Exercice-----

47 $8^2+7 \bmod 7$; $(8^2+7) \bmod 7$; #Attention mod est prioritaire sous xcas

(16, 2)

48 Si a=1, on a $x_n = x_0 + n.c$ [m], on fait donc par exemple: c=7,x0=1

49 l:=[seq((n*7+1) mod 1001,n=1..100)];

8 15 22 29 36 43 50 57 64 71 78 85 92 99 106 113 120 127 134 141 148 155

50 histogram(classes(l,0,1001/20)); #C'est tout a fait plat. Trop!

51 Remplacer les crochets par des accolades cree un ensemble, ce qui simplifie automatiquement les elements egaux

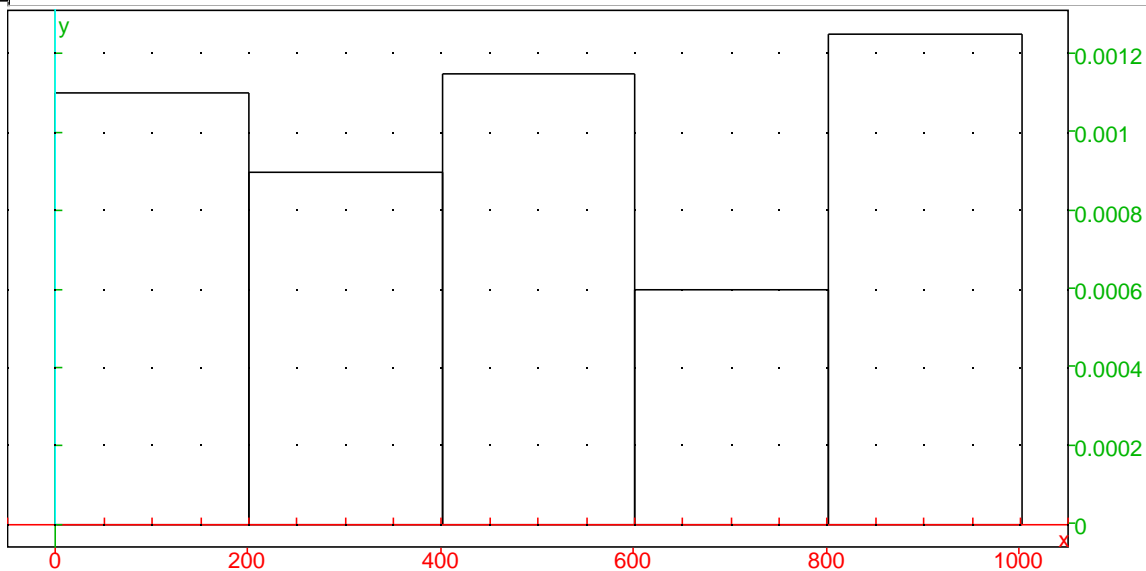
52 nops(l);nops({op(l)}); #ils sont tous distincts

(100, 100)

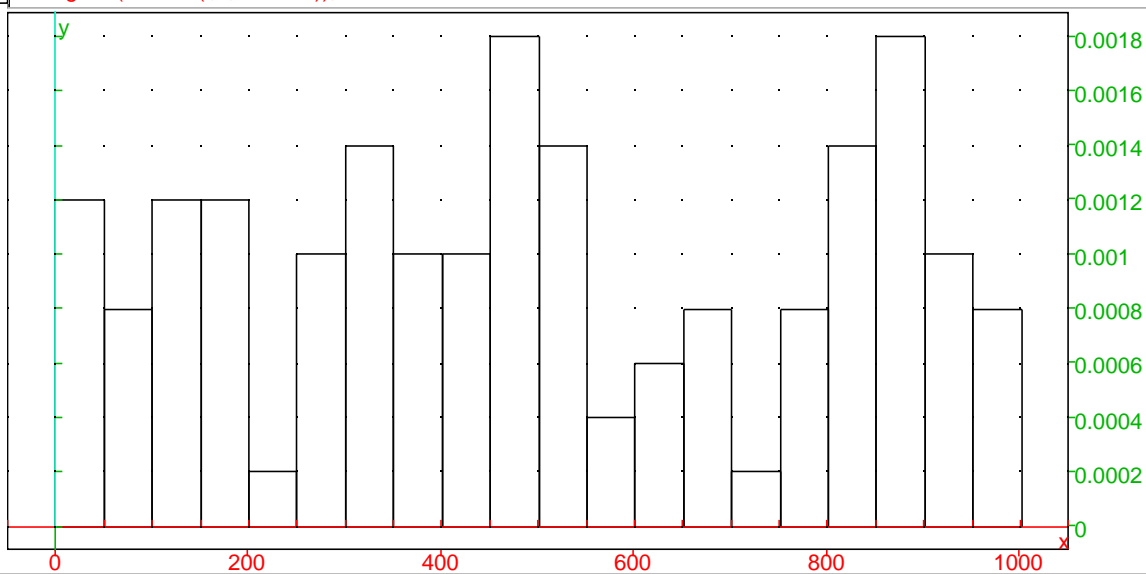
53 l:=[seq(rand(1001),n=1..100)];

896 774 541 303 453 464 177 281 74 885 147 182 25 939 66 814 975 112 76 333

```
54 histogram(classes(l,0,1001/5));
```



```
55 histogram(classes(l,0,1001/20));
```



```
56 nops(l);nops({op(l)});#il y a bien des anniversaires identiques
```

```
( 100, 93 )
```

```
57 On a en fait  $1001^{100}$  suites possibles, et  $1001!/901!$  suites dont  
tous les termes sont distincts.
```

```
58  $1001!/901!/1001.^{100}$ ;
```

```
0.00599058972
```

```
59 x:=1; a:=237;c:=54321;m:=10^4; l:=[];
```

```
( 1, 237, 54321, 10000, [] )
```

```
60 for i from 1 to 500 do x:=(a*x+c) mod m; l:=[op(l),x] ; od:
```

```
Done
```

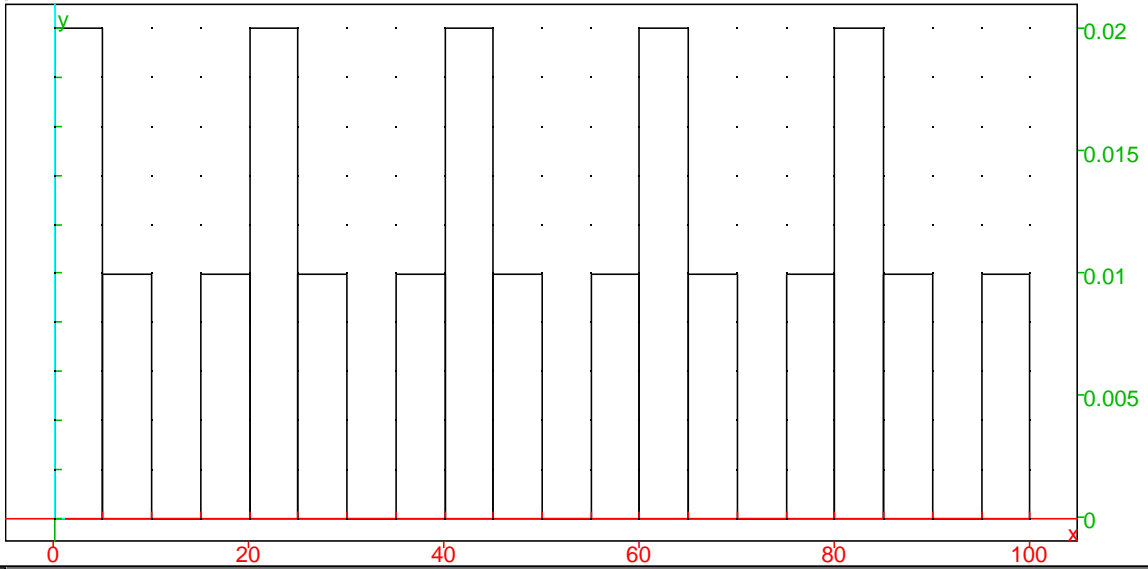
```
61 l1:=seq(i mod 10^2,i=1);
```

```
58 67 0 21 98 47 60 41 38 27 20 61 78 7 80 81 18 87 40 1 58 67 0 21 98
```

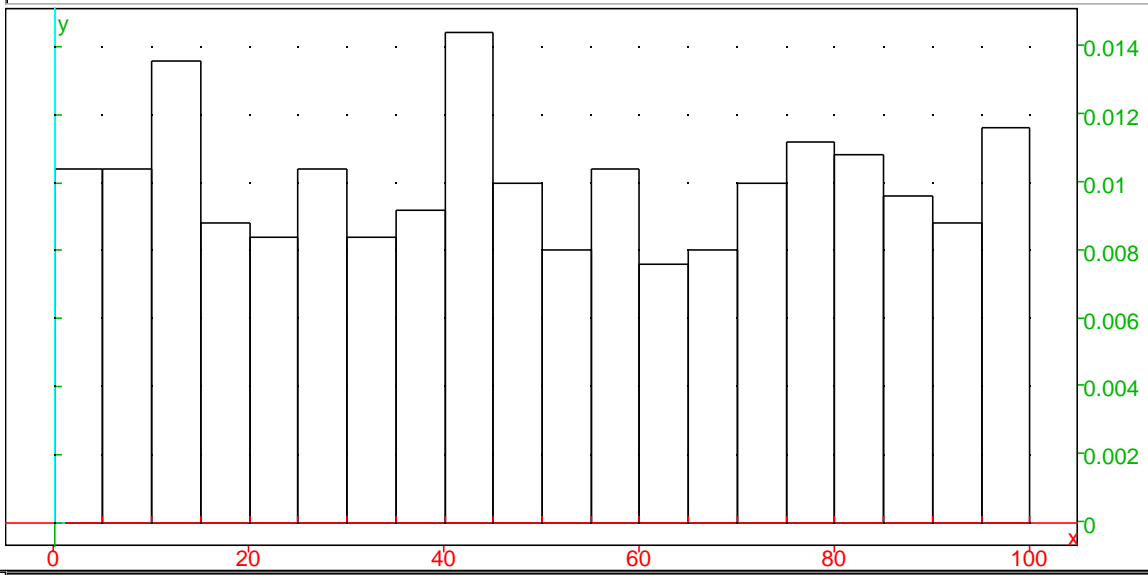
```
62 l2:=seq(trunc(i/10^2),i=1);
```

```
45 45 67 22 6 97 43 76 52 57 16 82 21 5 44 60 55 20 89 31 92 84 10 13 73
```

```
63 histogram(classes(1,0,5));
```



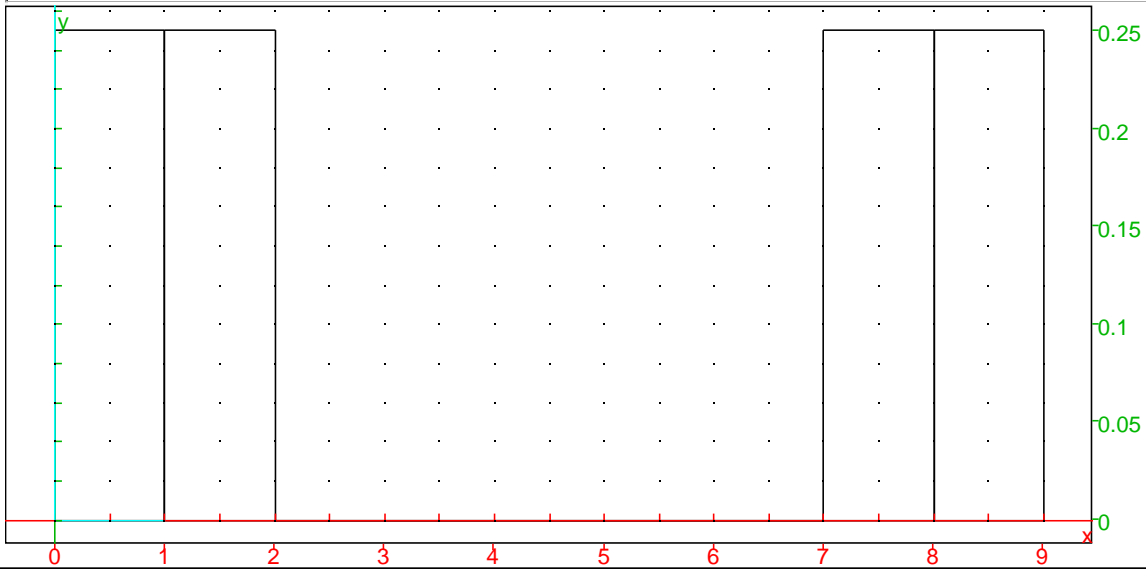
```
64 histogram(classes(12,0,5));
```



```
65 l1:=seq(i mod 10,i=1);
```

```
[ 8 7 0 1 8 7 0 1 8 7 0 1 8 7 0 1 8 7 0 1 8 7 0 1 8 7 0 1 8
```

66 histogram(classes(l1,0,1));



67 Si d est un diviseur de m et que $y_n = x_n [d]$, alors $y_{n+1} = ay_n + c [d]$
Donc si $d=10^2$, les 2 derniers chiffres de x_n ont p'eriode d'au plus d.

68 $x:=1; a:=237;c:=54321;m:=10^4-1; l:=[];$

(1, 237, 54321, 9999, [])

69 for i from 1 to 500 do $x:=(a*x+c) \bmod m; l:=\text{op}(l),x$; od:

Done

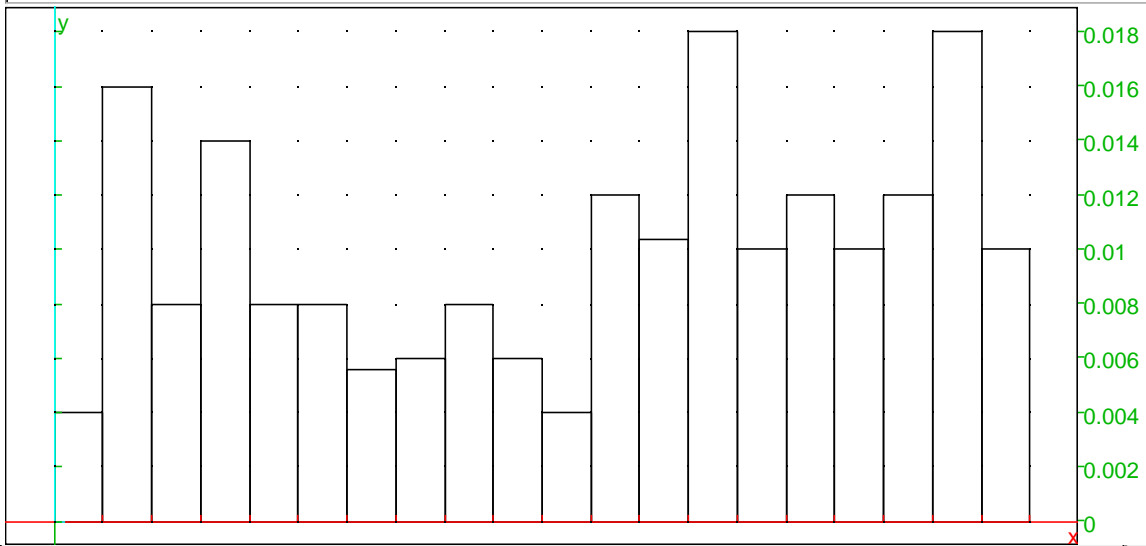
70 $l1:=\text{seq}(i \bmod 10^2, i=l);$

63 65 70 22 31 78 27 68 10 93 68 17 72 91 56 79 62 81 60 61 25 92 6 69

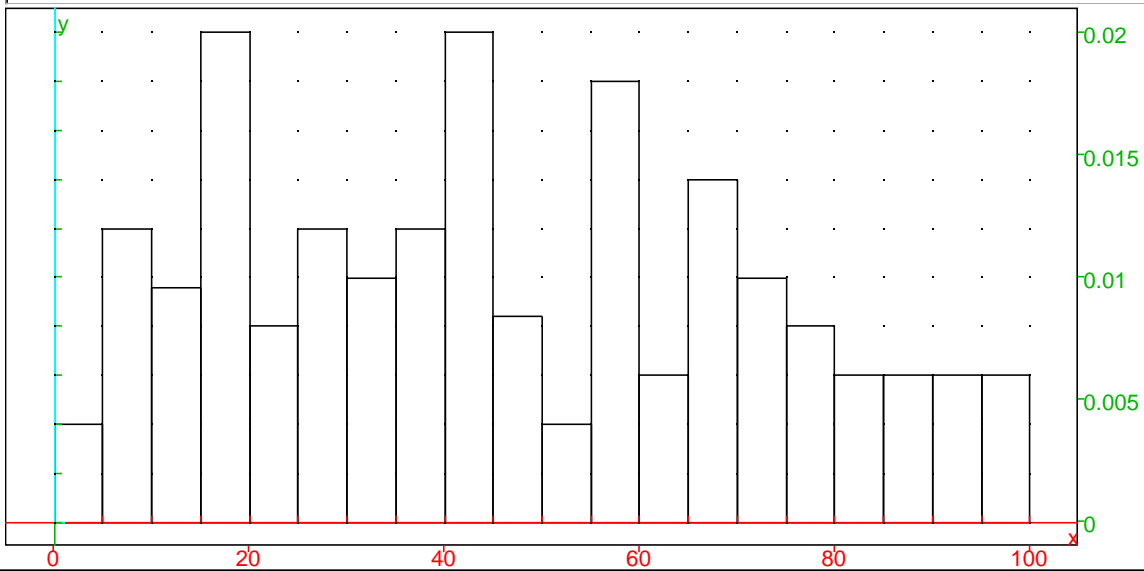
71 $l2:=\text{seq}(\text{trunc}(i/10^2), i=l);$

45 58 44 38 2 90 60 28 41 84 73 7 42 68 76 89 25 15 90 17 17 31 9 90 38

72 histogram(classes(l1,0,5));



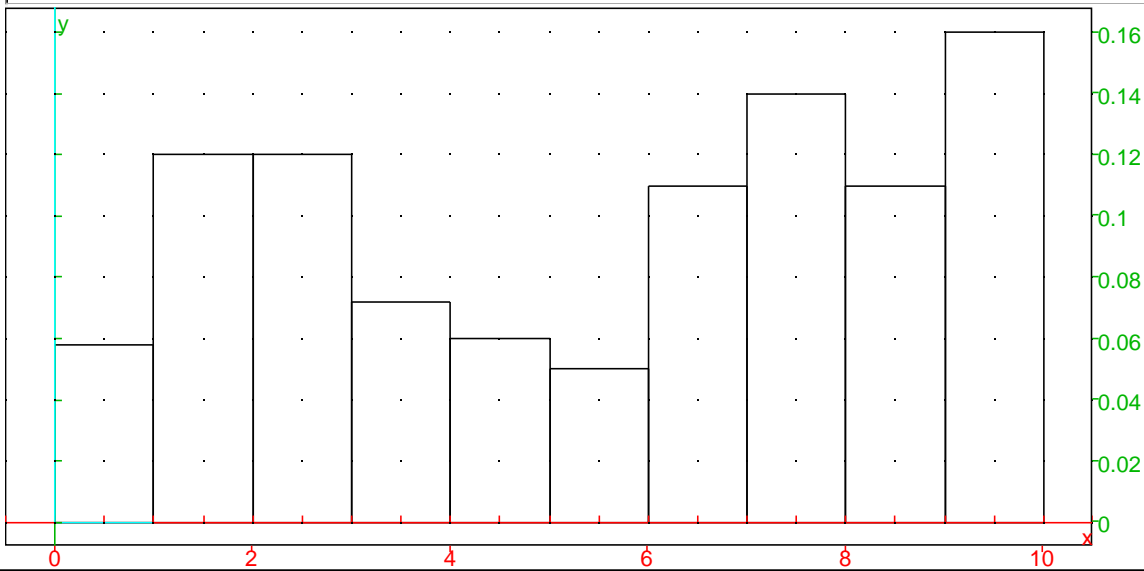
73 histogram(classes(l2,0,5));



74 l1:=seq(i mod 10,i=1);

3 5 0 2 1 8 7 8 0 3 8 7 2 1 6 9 2 1 0 1 5 2 6 9 4 6 1 9 9 5 4 1 8

75 histogram(classes(l1,0,1));



76 x:=1; a:=237;c:=54321;m:=prevprime(10^4); l:=[];

(1, 237, 54321, 9973, [])

77 for i from 1 to 500 do x:=(a*x+c) mod m; l:=op(l,x); od:

Done

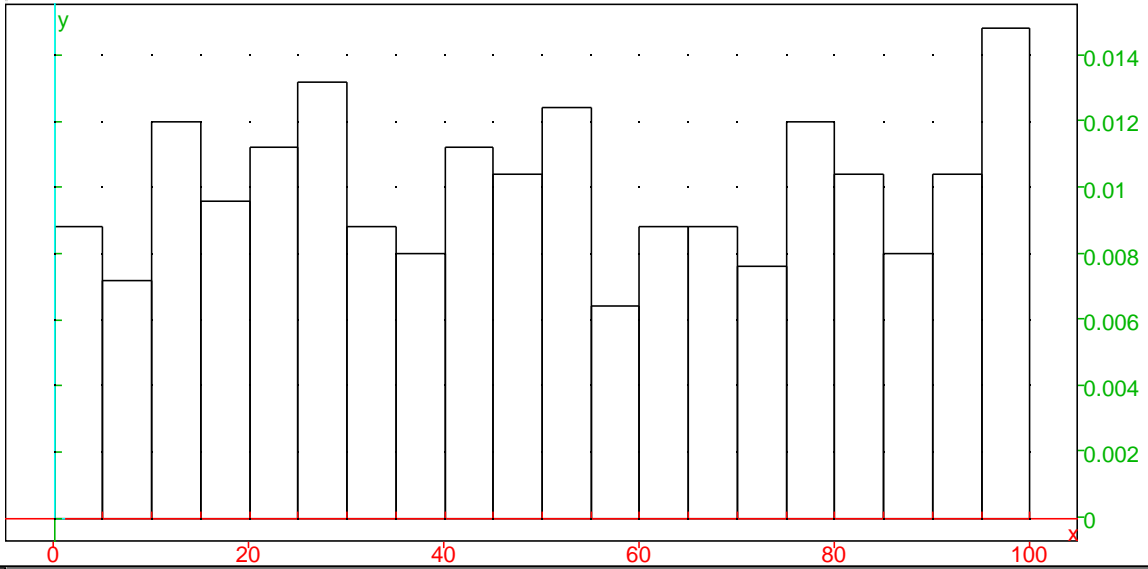
78 l1:=seq(i mod 10^2,i=1);

93 94 44 95 62 27 20 98 29 1 58 77 43 76 26 47 76 89 96 18 57 31 49 96

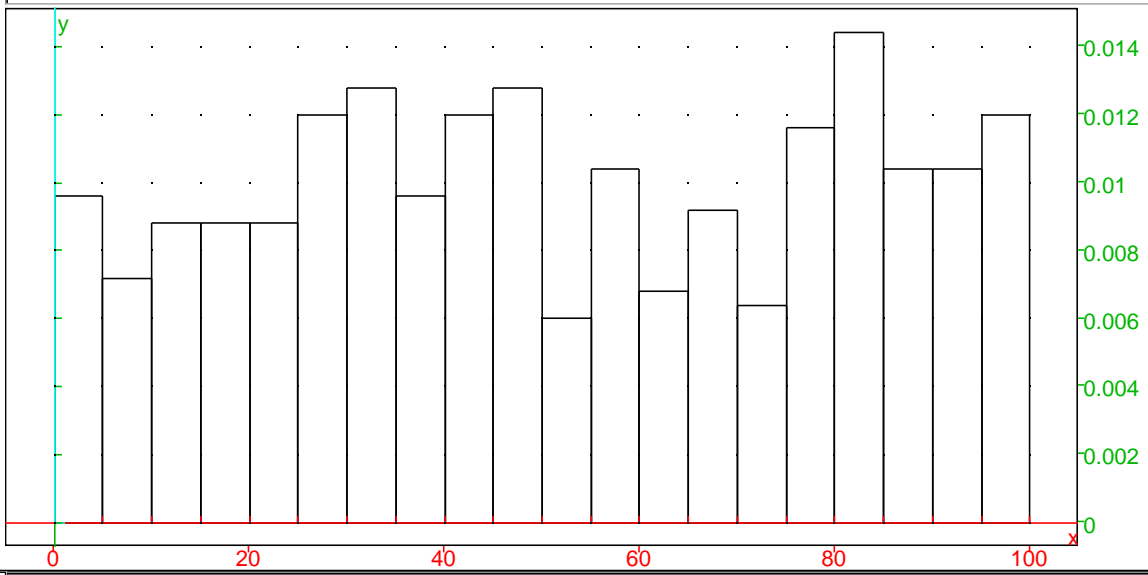
79 l2:=seq(trunc(i/10^2),i=1);

46 96 81 97 21 82 95 67 99 40 52 39 95 22 53 1 93 25 96 86 24 83 42 41


```
80 histogram(classes(l1,0,5));
```



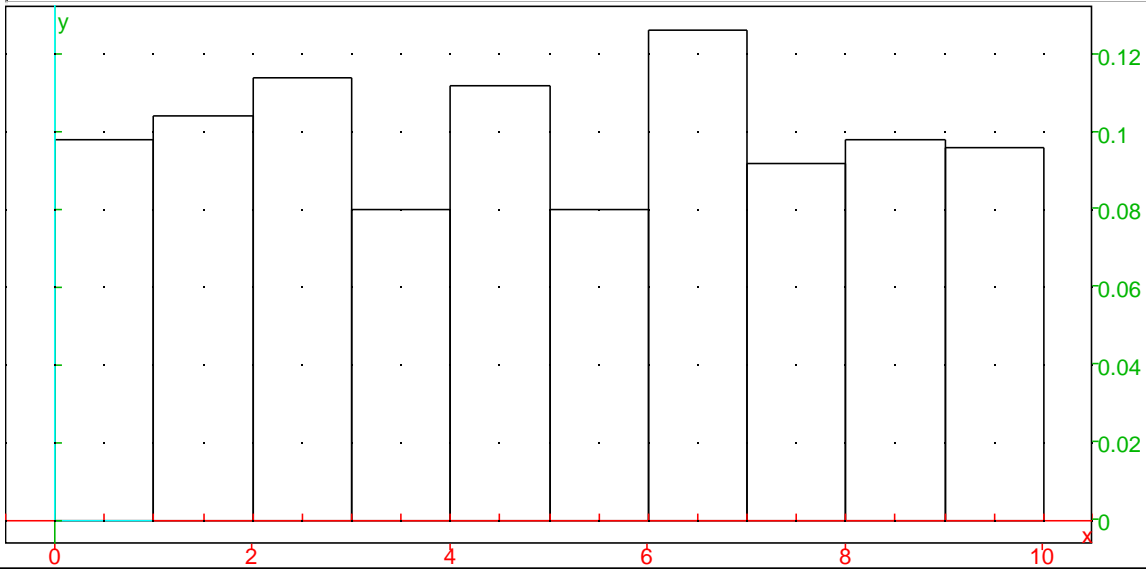
```
81 histogram(classes(l2,0,5));
```



```
82 l1:=seq(i mod 10,i=1);
```

```
[ 3 4 4 5 2 7 0 8 9 1 8 7 3 6 6 7 6 9 6 8 7 1 9 6 8 8 4 8 2 5 6 5 5
```

83 histogram(classes(l,0,1));



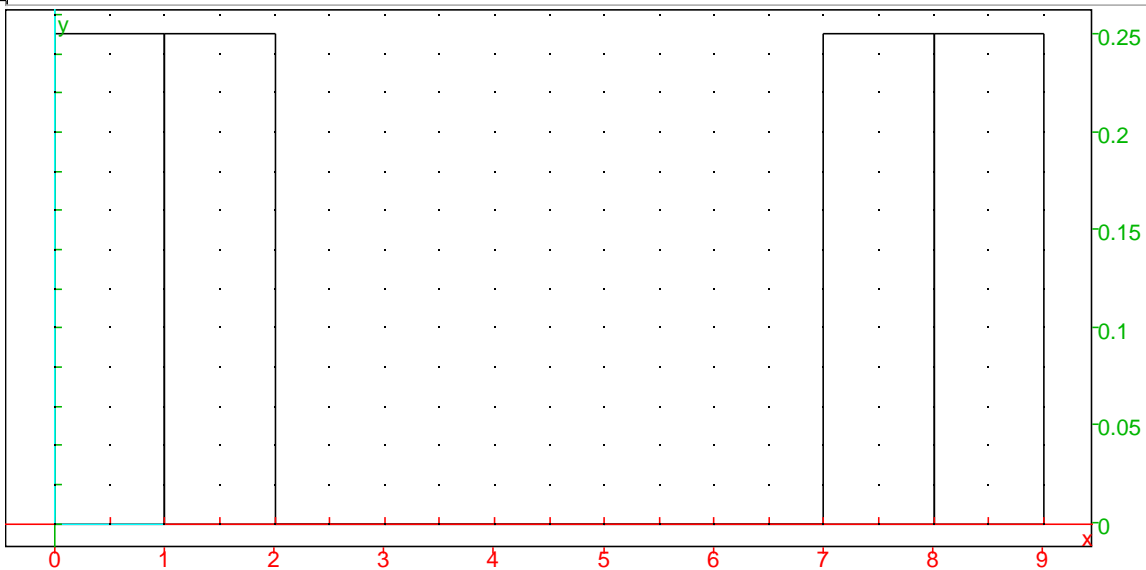
84 x:=1; a:=237;c:=54321;m:=10; l:=[];

(1, 237, 54321, 10, [])

85 for i from 1 to 100 do x:=(a*x+c) mod m; l:=[op(l),x] ; od;

Done

86 histogram(classes(l,0,1)); #certaines valeurs ne sont pas atteintes.



87

```
88 Prog Edit Add      |      |      |      |      |      |      |
   periode:= proc (a)
   x:=1;p:=1;
   if (a mod 5 <> 0) then while x>0 do p:=p+1;x:=(a*x+1) mod 125; od; fi;
   p; end proc;
```

// End defining periode

```
proc(a)
x:=1;
p:=1;
if (irem(a,5)<>0 then
while x>0 do
p:=p+1;
x:=irem(a*x+1,125);
od;
fi;
p;
```

```
89 periode(1),periode(5); #verification
      ( 125, 1 )
90 l:=seq(periode(i),i=2..124);max(l);
      100, 100, 50, 1, 125, 20, 100, 50, 1, 125, 100, 100, 50, 1, 125, 100, 20, 50, 1, 125, 100, 100
91 On a:  $x^n = (a^n - 1) / (a - 1) \pmod m$ . donc la periode est l'ordre de a.
```