Algorithmic and algebraic combinatorial aspects of polylogarithms with application on the computation of Drinfel'd associators

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Polylogarithms and polyzetas occur in

- Control theory (Lille).
- Analysis combinatorics (Flajolet, Labelle, Laforest, Salvy, Vallée, ...).
- Vassiliev knot invariants & Drinfel'd associator (Kontsevich, González-Lorka, Lê, Murakami, Furusho, Racinet, ...).
- Perturbative quantum field theory (Broadhurst, Kreimer, ...).
- Chern classes of a manifold (Hoffman, ...).
- K-theory (Gangl, Wojtkowiak, Zagier, ...).
- Irrationality & transcendence of $\zeta(2k+1)$ (Borwein, Ecalle, Goncharov, Zagier, ...).

Knizhnik-Zamolodchikov equation KZ_3

Drinfel'd constructed the solutions (80s) for the Knizhnik–Zamolodchikov equation KZ_3

$$\frac{dG(z)}{dz} = \frac{1}{2i\pi} \left(\frac{A}{z} + \frac{B}{z-1} \right) G(z), \quad 0 < z < 1,$$

where A,B are noncommuting symbols, and G(z) is a formal power series in A,B with coefficients that are analytic function of z.

$$G_1(z) \sim z^{A/2i\pi} = e^{A/2i\pi \log(z)}, \qquad z \to 0,$$

 $G_2(z) \sim (1-z)^{B/2i\pi} = e^{B/2i\pi \log(1-z)}, \qquad z \to 1,$

$$\Phi_{KZ}(A, B) = G_2(z)^{-1}G_1(z).$$

$$\Rightarrow G_1(1-z) \sim z^{B/2i\pi}\Phi_{KZ}(A, B), \quad z \to 0.$$

Question 1 How to compute the associator $\Phi_{KZ}(A,B)$?

Drinfel'd associators & $\Phi_{KZ}(A,B)$

Drinfel'd defined associator $\Phi(A,B)$ as a Lie exponential satisfying 3 relations :

- 1. Duality : $\Phi(B, A) = \Phi^{-1}(A, B)$.
- 2. Hexagonal: ...
- 3. Pentagonal: ...

$$\begin{split} & \Phi_{KZ} \in \mathsf{MZV}\langle\!\langle A, B \rangle\!\rangle \text{ (Lê & Murakami), where} \\ & \mathsf{MZV} = \left\{ \zeta(s_1, \dots, s_k) = \sum_{n_1 > \dots > n_k > 0} \frac{1}{n_1^{s_1} \cdots n_k^{s_k}}. \right\}. \end{split}$$

Drinfel'd proved also the existence of associators with *rational coefficients*.

Question 2 How to compute the rational associators?

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Non-commutative formal power series

 X^* : the free monoide generated by an alphabet X for the concatenation with ϵ (the empty word) as the neutral element.

A formal power series S is an infinite sum

$$S = \sum_{w \in X^*} \langle S | w \rangle \ w.$$

A finite FPS is called polynomial.

Let $x, y \in X, u, v \in X^*, xu \sqcup yv$ is the polynomial defined recursively as follows

$$xu \sqcup \epsilon = \epsilon \sqcup xu = xu,$$

 $xu \sqcup yv = y[(xu) \sqcup v] + x[u \sqcup (yv)].$

Example - $x_0x_1 \sqcup x_0x_1 = 4 x_0^2x_1^2 + 2 x_0x_1x_0x_1$.

 $\mathbb{C}\langle\langle X \rangle\rangle$, $\mathbb{C}\langle X \rangle$ denote the sets of FPS and polynomials over X and with coefficients in \mathbb{C} .

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Operations on formal power series

For $S, T \in \mathbb{C}\langle\langle X \rangle\rangle$, one defines

$$\begin{array}{rcl} \forall w \in X^*, \ \langle S+T|w \rangle & = & \langle S|w \rangle + \langle T|w \rangle, \\ \forall w \in X^*, & \langle ST|w \rangle & = & \sum\limits_{u,v \in X^*,uv = w} \langle S|v \rangle \langle T|u \rangle, \\ S \sqcup T & = & \sum\limits_{u,v \in X^*} \langle S|u \rangle \langle T|v \rangle u \sqcup v. \end{array}$$

 $\operatorname{Sh}_{\mathbb C}\langle X \rangle$ denotes the polynomial algebra equipped the *shuffle* product ul .

The exponential of S is the sum

$$\exp(S) = \sum_{k>0}^{\infty} \frac{S^k}{k!}.$$

The logarithm of 1+S is the sum

$$\log(1+S) = \sum_{k\geq 0}^{\infty} (-1)^{k+1} \frac{S^k}{k}.$$

Lyndon words and Standard factorization*

A Lyndon word is a non empty word that is less than each of its strict right factors (for the lexicographical ordering).

Example – Let $X = \{x_0, x_1\}, x_0 < x_1$. The Lyndon words of length ≤ 5 are the following (in lexicographically decreasing order):

$$\{x_0, x_0^4 x_1, x_0^3 x_1, x_0^3 x_1^2, x_0^2 x_1, x_0^2 x_1 x_0 x_1, x_0^2 x_1^2, x_0^2 x_1^3, x_0 x_1, x_0 x_1 x_0 x_1^2, x_0 x_1^2, x_0 x_1^3, x_0 x_1^4, x_1\}$$

 $\mathcal{L}yn(X)$ denotes the set of Lyndon words.

Let $l \in \mathcal{L}yn(X) \setminus X$. A standard factorization of l, noted by $\mathrm{st}(l)$, is the sole couple (u,v), where u,v are Lyndon words and v is the longest strict right factor of l verifying u < uv < v.

Example -
$$st(x_0^2x_1x_0x_1) = (x_0^2x_1, x_0x_1)$$
.

^{*}M. Lothaire.— Combinatorics on Words, Encyclopedia of Mathematics and its Applications, Addison-Wesley, 1983.

Sirsov lemma and Radford theorem

Lemma 1 (Širšov) For any
$$w \in X^*$$
,

$$w = l_1^{i_1} \dots l_k^{i_k}, \quad l_1 > \dots > l_k.$$

Example - Let
$$X = \{x_0, x_1\}, x_0 < x_1$$
.

$$x_1x_0x_1x_1x_0x_1x_1x_0x_0x_1 = x_1.x_0x_1x_1.x_0x_1x_1.x_0x_0x_1$$

= $x_1.x_0x_1x_1^2.x_0x_0x_1$,

here $x_1>x_0x_1x_1>x_0x_0x_1$ are Lyndon words. \Box

Theorem 1 (Radford) The \mathbb{C} -algebra $\operatorname{Sh}_{\mathbb{C}}\langle X \rangle$ is the polynomial algebra generated by $\mathcal{L}yn(X)$.

Example – Let
$$X = \{x_0, x_1\}, x_0 < x_1$$
.

$$\begin{aligned} x_0 x_1 x_0^2 x_1 &= x_0 x_1 \sqcup x_0^2 x_1 - 3 \ x_0^2 x_1 x_0 x_1 - 6 \ x_0^3 x_1^2, \\ x_0^3 x_1 x_0^4 x_1 &= x_0^3 x_1 \sqcup x_0^4 x_1 - 5 \ x_0^4 x_1 x_0^3 x_1 \\ &- 15 \ x_0^5 x_1 x_0^2 x_1 - 35 \ x_0^6 x_1 x_0 x_1 - 70 \ x_0^7 x_1^2. \end{aligned}$$

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Free Lie algebra

The free Lie algebra, noted by $\mathcal{L}ie_{\mathbb{C}}\langle X\rangle$, is the \mathbb{C} -algebra of polynomials, over X, equipped the bracket [.,.] defined as follows

$$\forall P, Q \in \mathcal{L}ie_{\mathbb{C}}\langle X \rangle, \quad [P, Q] = PQ - QP$$

and verifying the following properties

$$[P, P] = 0,$$

$$[P, [Q, R]] + [Q, [R, P]] + [R, [P, Q]] = 0.$$

An element of $\mathcal{L}ie_{\mathbb{C}}\langle X\rangle$ is called *Lie polynomial*.

Let $S \in \mathbb{C}\langle\!\langle X \rangle\!\rangle$, S is called Lie series if it can be written as follows

$$S = \sum_{k \ge 1} P_k,$$

where P_k is a homogenous Lie polynomial of degree k. $\mathcal{L}ie_{\mathbb{C}}\langle\!\langle X \rangle\!\rangle$ denotes the set of Lie series over X.

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PBW basis and dual basis

The bracket form P_l of a Lyndon word l is defined recursively by

$$\begin{cases} P_l = [P_u, P_v] = P_u P_v - P_v P_u & \text{ for } uv = \operatorname{st}(l), \\ P_x = x & \text{ for } x \in X, \end{cases}$$

The set $\{P_l;\ l\in\mathcal{L}yn(X)\}$ is a basis for the free Lie algebra $\mathcal{L}ie_{\mathbb{C}}\langle X\rangle$.

The PBW basis $\mathcal{B} = \{P_w; w \in X^*\}$ is obtained by setting

$$P_v = P_{l_1}^{i_1} P_{l_2}^{i_2} \dots P_{l_k}^{i_k}, \quad v = l_1^{i_1} \dots l_k^{i_k}, l_1 > \dots > l_k,$$

and its dual basis \mathcal{B}^* is obtained by setting

$$\begin{cases} S_{l} = xS_{w}, & \forall l \in \mathcal{L}yn(X), l = xw, x \in X, w \in X^{*}, \\ S_{w} = \frac{S_{l_{1}}^{\sqcup \sqcup i_{1}} \sqcup \ldots \sqcup S_{l_{k}}^{\sqcup \sqcup i_{k}}}{i_{1}! \ldots i_{k}!}, w = l_{1}^{i_{1}} \ldots l_{k}^{i_{k}}, l_{1} > \ldots > l_{k}. \end{cases}$$

Example

l	P_l	$S_l = P_l^* = [l]$
x_0	x_0	x_0
x_1	x_1	x_1
x_0x_1	$[x_0, x_1]$	x_0x_1
$x_0^2 x_1$	$[x_0, [x_0, x_1]]$	$x_0^2 x_1$
$x_0x_1^2$	$[[x_0,x_1],x_1]$	$x_0x_1^2$
$x_0^3 x_1$	$[x_0, [x_0, [x_0, x_1]]]$	$x_0^3x_1$
$x_0^2 x_1^2$	$[x_0, [[x_0, x_1], x_1]]$	$x_0^2 x_1^2$
$x_0x_1^3$	$[[[x_0, x_1], x_1], x_1]$	$x_0x_1^3$
$x_0^4 x_1$	$[x_0, [x_0, [x_0, [x_0, x_1]]]]$	$x_0^4 x_1$
$x_0^3 x_1^2$	$[x_0, [x_0, [[x_0, x_1], x_1]]]$	$x_0^3 x_1^2$
$x_0^2 x_1 x_0 x_1$	$[[x_0, [x_0, x_1]], [x_0, x_1]]$	$2x_0^3x_1^2 + x_0^2x_1x_0x_1$
$x_0^2 x_1^3$	$[x_0,[[[x_0,x_1],x_1],x_1]]$	$x_0^2 x_1^3$
$x_0x_1x_0x_1^2$	$[[x_0,x_1],[[x_0,x_1],x_1]]$	$3x_0^2x_1^3 + x_0x_1x_0x_1^2$
$x_0x_1^4$	$[[[[x_0,x_1],x_1],x_1],x_1]$	$x_0x_1^4$
$x_0^5 x_1$	$[x_0, [x_0, [x_0, [x_0, [x_0, x_1]]]]]$	$x_0^5 x_1$
$x_0^4 x_1^2$	$[x_0, [x_0, [x_0, [[x_0, x_1], x_1]]]]$	$x_0^4 x_1^2$
$x_0^3 x_1 x_0 x_1$	$[x_0, [[x_0, [x_0, x_1]], [x_0, x_1]]]$	$2x_0^4x_1^2 + x_0^3x_1x_0x_1$
$x_0^3 x_1^3$	$[x_0, [x_0, [[[x_0, x_1], x_1], x_1]]]$	$x_0^3 x_1^3$
$x_0^2 x_1 x_0 x_1^2$	$[x_0, [[x_0, x_1], [[x_0, x_1], x_1]]]$	$3x_0^3x_1^3 + x_0^2x_1x_0x_1^2$
$x_0^2 x_1^2 x_0 x_1$	$[[x_0,[[x_0,x_1],x_1]],[x_0,x_1]]$	$6x_0^3x_1^3 + 3x_0^2x_1x_0x_1^2 + x_0^2x_1^2x_0x_1$
$x_0^2 x_1^4$	$[x_0,[[[[x_0,x_1],x_1],x_1],x_1]]$	$x_0^2 x_1^4$
$x_0x_1x_0x_1^3$	$[[x_0, x_1], [[[x_0, x_1], x_1], x_1]]$	$4x_0^2x_1^4 + x_0x_1x_0x_1^3$
$x_0x_1^5$	$[[[[[x_0, x_1], x_1], x_1], x_1], x_1]$	$x_0x_1^5$

Factorization Hopf algebra

 $\mathbb{C}\langle X \rangle \otimes \mathbb{C}\langle X \rangle$ denotes the tensorial product of $\mathbb{C}\langle X \rangle$ with itself. The co-product Φ of the concatenation is defined as follows

$$\forall u, v \in X^*, \quad \langle \Phi w | u \otimes v \rangle = \langle uv | w \rangle \\ \iff \Phi w = \sum_{u,v \in X^*, uv = w} u \otimes v.$$

 Φ is an morphism for the shuffle algebra:

$$\forall u, v \in X^*, \quad \Phi(u \sqcup v) = \Phi(u) \sqcup \Phi(v),$$

A co-unity e is defined by :

$$\begin{array}{ccc} e: \mathbb{C}\langle X\rangle & \longrightarrow & \mathbb{C}\langle X\rangle, \\ P & \longmapsto & e(P) = \langle P|\epsilon\rangle. \end{array}$$

For $S \in \mathbb{C}\langle\!\langle X \rangle\!\rangle$, the antipode of S is the following FPS (\tilde{w} denotes the *miror* of w)

$$a(S) = \sum_{w \in X^*} (-1)^{|w|} \langle S|w \rangle \ \tilde{w}.$$

 $(\mathbb{C}\langle\langle X \rangle\rangle, \sqcup, 1, \Phi, e, a)$ is the factorization Hopf algebra.

Decomposition Hopf algebra

The map Γ_2 is defined as follows

$$\forall u, v, w \in X^*, \quad \langle \Gamma_2 w | u \otimes v \rangle = \langle w | u \sqcup v \rangle.$$

In particular

$$\forall x \in X, \quad \Gamma_2 x = 1 \otimes x + x \otimes 1.$$

It is extended to $\mathbb{C}\langle\langle X \rangle\rangle$ as follows

$$\langle \Gamma_2 S | u \otimes v \rangle = \sum_{w \in X*} \langle S | w \rangle \Gamma_2 w = \langle S | u \sqcup v \rangle.$$

 Γ_2 is an morphism for the associative algebra :

$$\forall u, v \in X^*, \quad \Gamma_2(uv) = \Gamma_2(u)\Gamma_2(v).$$

And it is a co-associative coproduct.

 $(\mathbb{C}\langle\langle X\rangle\rangle,.,1,\Gamma_2,e,a)$ is the decomposition Hopf algebra.

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Primitive and group-like

Let $S \in \mathbb{C}\langle\langle X \rangle\rangle$, S is called *primitive* if

$$\Gamma_2 S = 1 \otimes S + S \otimes 1.$$

S is called group-like if

$$\Gamma_2 S = S \otimes S$$
.

S verifies the *Friedrichs criterion* if

$$\forall u, v \in X^*, \quad \langle S|u \sqcup v \rangle = \langle S|u \rangle \langle S|v \rangle.$$

Theorem 2 (Ree)

$$S \in \mathcal{L}ie_{\mathbb{C}}\langle\langle X \rangle\rangle$$

 \iff S is primitive

 \iff e^S is group-like

 \iff e^S verifies the Friedrichs criterion.

Diagonal series & Schützenberger factorization*

Let us consider, in the completed tensorial product $\mathbb{C}\langle X\rangle\widehat{\otimes}\mathbb{C}\langle X\rangle$, the following operation : the shuffle product for the left factor, the concatenation for right factor (for $u_1, v_1, u_2, v_2 \in X^*$):

$$(u_1 \otimes v_1)(u_2 \otimes v_2) = (u_1 \sqcup u_2) \otimes (v_1 v_2).$$

By a Schützenberger factorization, the following diagonal series in $\mathbb{C}\langle X
angle\widehat{\otimes}\,\mathbb{C}\langle X
angle$

$$\mathcal{D} = \sum_{w \in X^*} w \otimes w$$

can be factorized in an infinite product, indexed by the Lyndon words:

$$\mathcal{D} = e^{x_1 \otimes x_1} \left[\prod_{l \in \mathcal{L}yn(X) \setminus X} e^{P_l^* \otimes P_l} \right] e^{x_0 \otimes x_0}$$

$$\prod_{l \in \mathcal{L}yn(X) \setminus X} e^{P_l^* \otimes P_l} = e^{-x_1 \otimes x_1} \mathcal{D} e^{-x_0 \otimes x_0}.$$

$$\iff \prod_{l \in \mathcal{L}yn(X) \setminus X} e^{P_l^* \otimes P_l} = e^{-x_1 \otimes x_1} \mathcal{D} e^{-x_0 \otimes x_0}$$

*C. Reutenauer.— Free Lie Algebras, London Math. Soc. Monog. 7 (new series), Clarendon Press-Oxford Sciences Publications, 1993.

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From iterated integrals to words

The iterated integrals over the differential forms $\{\omega_0(z),\ldots,\omega_n(z)\}$ can be encoded by the words $w=x_{i_1}\ldots x_{i_k}$ over $X=\{x_0,\ldots,x_n\}$ (Fliess):

$$\begin{split} \alpha_{z_0}^z(w) &= \int_{z_0}^z \omega_{i_1}(t_1) \dots \omega_{i_k}(t_k) \\ &= \left\{ \begin{array}{ccc} 1 & \text{if } w = \epsilon, \\ \int_{z_0}^z \omega_i(t) \alpha_{z_0}^t(v) & \text{if } w = x_i v. \end{array} \right. \end{split}$$

Remark — In control theory, Fliess takes the differential forms ω_i that are in the form $a_i(t)dt$, where $a_i(t)$ are real piecewise continous. \square

 α is a \mathbb{C} -algebra morphism for " \square " (**Chen**): $\alpha: \mathsf{Sh}_{\mathbb{C}}\langle X \rangle \to \{\mathsf{comb. of iterated integrals}, +, .\}.$

$$\begin{array}{rcl} \forall u,v \in X^* \setminus \{\epsilon\}, & \alpha_{z_0}^z(u+v) & = & \alpha_{z_0}^z(u) + \alpha_{z_0}^z(v), \\ \forall \lambda \in \mathbb{C}, u \in X^*, & \alpha_{z_0}^z(\lambda u) & = & \lambda \alpha_{z_0}^z(u), \\ \forall u,v \in X^*, & \alpha_{z_0}^z(u \sqcup v) & = & \alpha_{z_0}^z(u) \alpha_{z_0}^z(v). \end{array}$$

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Iterated integral and shuffle algebra

Let us consider the following differential forms

$$\omega_0(z) = \frac{dz}{z}$$
 and $\omega_2(z) = dz$.

Example - Note that

$$x_0 \sqcup x_2 = x_0 x_2 + x_2 x_0.$$

But $\alpha_0^z(x_0 \sqcup x_2) = \alpha_0^z(x_0)\alpha_0^z(x_2)$ and $\alpha_0^z(x_2x_0)$ diverge while $\alpha_0^z(x_2) = z$. \square

Example – For any n > 0, one has

$$\alpha_0^z(x_0^n x_2) = \alpha_0^z(x_2) = z.$$

Theorem 3 (FPSAC98) For $\omega_0 = dz/z$, $\omega_1 = dz/(1-z)$, α is injective from $\operatorname{Sh}_{\mathbb{C}}\langle x_0, x_1 \rangle$ to the smallest algebra that contains \mathbb{C} and that is stable under integration with respect to ω_0, ω_1 .

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From word to polylogarithm

Definition 1 For any word $v \in X^*x_1$. Let us define the polylogarithms as follows:

$$\begin{split} \operatorname{Li}_{x_0 v}(z) &= \alpha_0^z(x_0 v) = \int_0^z \omega_0(t) \operatorname{Li}_v(t), \\ \operatorname{Li}_{x_1 v}(z) &= \alpha_0^z(x_1 v) = \int_0^z \omega_1(t) \operatorname{Li}_v(t). \end{split}$$

And for any $v \in x_0 X^* x_1$, the polyzetas as

$$\zeta(v) = \operatorname{Li}_v(1)$$
.

Fact 1 For $v = x_0^{s_1-1} x_1 \dots x_0^{s_k-1} x_1$, one has :

$$\mathsf{Li}_v(z) = \sum_{n_1 > \ldots > n_k > 0} \frac{z^{n_1}}{n_1^{s_1} \ldots n_k^{s_k}}$$

Fact 2
$$v = x_0^{s_1-1}x_1 \dots x_0^{s_k-1}x_1 \longleftrightarrow (s_1, \dots, s_k).$$

 $\text{Li}_v(z) = \text{Li}_{s_1, \dots, s_k}(z)$ and $\zeta(v) = \zeta(s_1, \dots, s_k).$

We extend the definion 1 over X^* by putting

$$\operatorname{Li}_{x_0^n}(z) = \frac{\log^n(z)}{n!}, \ \operatorname{Li}_{x_1 x_0^n}(z) = \int_0^z \omega_1(t) \frac{\log^n(t)}{n!}.$$

Non-commutative g.s. of $Li_w(z)$

Definition 2
$$L(z) = \sum_{w \in X^*} \operatorname{Li}_w(z) w.$$

Proposition 1 (FPSAC98) L(z) satisfies the differential equation (Drinfel'd equation):

$$dL(z) = [x_0\omega_0(z) + x_1\omega_1(z)]L(z)$$

with the boundary condition

$$L(\varepsilon) = e^{x_0 \log \varepsilon} + o(\sqrt{\varepsilon})$$
 for $\varepsilon \to 0^+$.

Proof - (sketched) Observing that

$$L(z) = 1 + \sum_{u \in X^*} \operatorname{Li}_{x_0 u}(z) x_0 u + \sum_{v \in X^*} \operatorname{Li}_{x_1 v}(z) x_1 v.$$

The exponential term $e^{\log \varepsilon x_0}$ comes from the definition of $\text{Li}_{x_0^n}, n \geq 1$.

The coefficient of each other word w in $L(\varepsilon)$ is easily seen to be bounded by $o(\varepsilon^n \log^m \varepsilon)$, where n is the number of x_1 's in w. \square

Solutions of Drinfel'd equation

Proposition 2 If G(z) and H(z) are solutions of Drinfel'd equation then

$$d[H(z)^{-1}G(z)] = 0.$$

Proof – Since $H(z)H(z)^{-1}=1$ then

$$[dH(z)]H(z)^{-1} = -H(z)[dH(z)^{-1}].$$

Therefore if H(z) is solution then

$$[dH(z)^{-1}] = -H(z)^{-1}[dH(z)]H(z)^{-1}$$

$$= -H(z)^{-1}[x_0\omega_0(z) + x_1\omega_1(z)],$$

$$d[H(z)^{-1}G(z)] = [dH(z)^{-1}]G(z) + H(z)^{-1}[dG(z)]$$

$$= -H(z)^{-1}[x_0\omega_0(z) + x_1\omega_1(z)]G(z)$$

$$+H(z)^{-1}[x_0\omega_0(z) + x_1\omega_1(z)]G(z).$$

We get then the expected result. \square

Corollary 1 Let g_* be the substitution morphism defined by $g_*x_0=-x_1, g_*x_1=-x_0$. If H(z) is solution of Drinfel'd equation then

$$d[H(z)^{-1}g_*H(1-z)] = 0.$$

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Factorization of the g.s. L(z)

Corollary 2

$$\forall u, v \in X^*, \ \mathsf{Li}_{u+1}(z) = \mathsf{Li}_u(z) \, \mathsf{Li}_v(z).$$

Proof – Use the Friedrichs criterion. \square

Corollary 3

$$L(z) = e^{-\log(1-z)x_1} L_{\text{reg}}(z) e^{\log(z)x_0},$$

where

$$L_{\operatorname{reg}}(z) = \prod_{l \in \mathcal{L}yn(X) \setminus \{x_0, x_1\}}^{\searrow} \exp(\operatorname{Li}_{P_l^*}(z)P_l).$$

Proof – Use the Schützenberger factorization. $\hfill\Box$

L(z) is groupe like

Theorem 4 (FPSAC98) $\Delta L(z) = L(z) \otimes L(z)$.

Proof – (sketched) Intuitively speaking, it follows from the boundary condition and thus the limit at 0 of L(z) is a Lie exponential, and L(z) is a Lie exponential for any z.

We have to prove $T(z) = \Delta L(z) - L(z) \otimes L(z)$ vanishes for all z. We claim that T satisfies

$$dT(z) = (\Delta V(z)) T(z)dz,$$

$$\lim_{\varepsilon \to 0^+} T(\varepsilon) = 0,$$

where $V(z)=[x_0\omega_0(z)+x_1\omega_1(z)]$. Thus we have a recursive formula for the coefficients of T(z) by means of differential equations with limit conditions in 0. Since these limits all vanish in 0, it follows by induction that the coefficients of T all vanish globally. \square

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Asymptotic behaviour at z=1

Corollary 4 The asymptotic expansion of L(z) at z=1 is given by :

$$L(1-\varepsilon) \sim e^{-x_1 \log \varepsilon} L_{\text{reg}}(1) e^{x_0 \varepsilon}, \quad \text{for} \quad \varepsilon \to 0^+.$$

Example – For $\varepsilon \to 0^+$, we have

$$\operatorname{Li}_{x_0}(1-\varepsilon)\sim \varepsilon$$
 and $\operatorname{Li}_{x_1}(1-\varepsilon)\sim -\log \varepsilon$.

The Radford theorem gives

$$x_1^2 x_0 = x_0 x_1^2 - x_0 x_1 \sqcup x_1 + 1/2 x_0 \sqcup x_1^{\sqcup 2}.$$

Therefore

$$\operatorname{Li}_{x_1^2 x_0}(1-\varepsilon) \sim \zeta(2,1) + \zeta(2) \log \varepsilon - \frac{1}{2}\varepsilon \log^2 \varepsilon + \dots$$
$$\sim \zeta(3) + \zeta(2) \log \varepsilon - \frac{1}{2}\varepsilon \log^2 \varepsilon + \dots$$

The last expression is obtained by use of the Euler's identity $\zeta(2,1)=\zeta(3)$. \square

In the other words, for any $w \in X^*$, for $\varepsilon \to 0^+$,

$$\mathsf{Li}_w(1-arepsilon) \sim \sum_{i>1} Q_{w,i}(\log arepsilon) arepsilon^i.$$

Non-commutative g.s. of polyzetas

Let
$$\zeta_{\coprod} = \zeta \circ \operatorname{reg}_{\coprod}$$
, where

$$\begin{array}{ll} \operatorname{reg}_{\sqcup \sqcup} & : & \mathbb{C}\langle\langle X \rangle\rangle \to \mathbb{C}\langle\langle X \rangle\rangle, \\ \operatorname{such that} & \operatorname{reg}_{\sqcup \sqcup} x_0 = \operatorname{reg}_{\sqcup \sqcup} x_1 = 0, \\ \forall w \in x_0 X^* x_1, & \operatorname{reg}_{\sqcup \sqcup} w = w, \\ \forall u, v \in X^*, & \operatorname{reg}_{\sqcup \sqcup} u \sqcup v = \operatorname{reg}_{\sqcup \sqcup} u \sqcup \operatorname{reg}_{\sqcup \sqcup} v. \end{array}$$

Definition 3
$$Z = \sum_{w \in \{x_0, x_1\}^*} \zeta_{\coprod}(w) \ w.$$

Theorem 5 (FPSAC98)

$$Z = L_{\text{reg}}(1) = \prod_{l \in \mathcal{L}yn(X) \setminus \{x_0, x_1\}}^{\searrow} \exp[\zeta(P_l^*) P_l].$$

Proof - Z is the image by $\zeta_{\square \square} \otimes id$ of \mathcal{D} . \square

Corollary 5 $\forall u, v \in X^*, \zeta_{\coprod}(u \coprod v) = \zeta_{\coprod}(u)\zeta_{\coprod}(v).$ Therefore, for any convergent words u and v, $\zeta(u \coprod v) = \zeta(u)\zeta(v).$

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Z and $\log Z$ up to order 4 by computer

$$\begin{split} Z &= \cdots e^{\frac{2}{5}\zeta(2)^2[x_0,[x_0,[x_0,x_1]]]} \cdots e^{\zeta(3)[[x_0,x_1],x_1]} \cdots \\ & \cdots e^{\zeta(2)[x_0,x_1]} \cdots e^{\frac{1}{10}\zeta(2)^2[x_0,[[x_0,x_1],x_1]]} \cdots \\ & \cdots e^{\zeta(3)[x_0,[x_0,x_1]]} \cdots e^{\frac{2}{5}\zeta(2)^2[x_0,[x_0,[x_0,x_1]]]} \cdots \\ & = 1 + \zeta(2)[x_0,x_1] \\ & + \zeta(3)([x_0,[x_0,x_1]] + [[x_0,x_1],x_1]) \\ & + \frac{2}{5}\zeta(2)^2([x_0,[x_0,[x_0,x_1]]] + [[[x_0,x_1],x_1],x_1] \\ & + \frac{5}{8}[x_0,x_1]^2 + \frac{1}{4}[x_0,[[x_0,x_1],x_1]]) + \cdots, \\ \log Z &= \zeta(2)[x_0,x_1] \\ & + \zeta(3)([x_0,[x_0,x_1]] + [[x_0,x_1],x_1]) \\ & + \frac{2}{5}\zeta(2)^2([x_0,[x_0,[x_0,x_1]]] + [[[x_0,x_1],x_1],x_1] \\ & + \frac{1}{4}[x_0,[[x_0,x_1],x_1]]) + \cdots \end{split}$$

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Chen series & analytic continuation of L(z)

For a differentiable path $\gamma:[0,1]\to\mathbb{C}\backslash\{0,1\}$ between a and b, let S_γ be the evaluation at b of the solution of the differential equation

$$\begin{cases} dS_{\gamma}(z) = [x_0\omega_0(z) + x_1\omega_1(z)]S_{\gamma}(z), \\ S_{\gamma}(a) = 1. \end{cases}$$

 $S_{\gamma} \in \mathbb{C}\langle X \rangle$ is called the *Chen series* along γ . S_{γ} is a Lie exponential and it depends only on the homotopy class of γ (**Chen**).

Proposition 3 (FPSAC98) Let $z_0 \rightsquigarrow z$ be a differentiable path on $\mathbb{C}\setminus\{0,1\}$ s.t. L admits an analytic continuation. Then $L(z) = S_{z_0 \rightsquigarrow z} L(z_0)$.

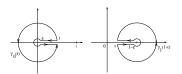
Proof – L(z) and $S_{z_0 \leadsto z} L(z_0)$ satisfy the Drinfel'd equation taking the same value at z_0 . \square

Corollary 6

$$S_{\varepsilon \to 1-\varepsilon} \sim e^{-x_1 \log \varepsilon} Z e^{-x_0 \log \varepsilon}$$
 for $\varepsilon \to 0^+$.

Proof – $S_{\varepsilon \to 1-\varepsilon} = L(1-\varepsilon)L(\varepsilon)^{-1}$ and the behaviour of L lead to the expected result. \square

Monodromy of the g.s. L(z)



Paths of integration

Theorem 6 (FPSAC98) The monodromy of L(t) for $t \in]0,1[$ around 0 and 1 is given by

$$\mathcal{M}_0 L(t) = L(t) e^{2i\pi x_0},$$

 $\mathcal{M}_1 L(t) = L(t) Z^{-1} e^{-2i\pi x_1} Z = L(t) e^{2i\pi m_1},$

where \mathfrak{m}_1 is a Lie series given by the formula

$$\mathfrak{m}_1 = \prod_{l \in \mathcal{L}yn(X) \backslash \{x_0, x_1\}}^{\searrow} e^{-\zeta_{P_l^*} \operatorname{ad} P_l} (-x_1).$$

Proof of the monodromy theorem

• Monodromy of L(z) around 0

$$\begin{split} \mathcal{M}_0 L(t) &= S_{\varepsilon \to t} S_{\gamma_0(\varepsilon)} S_{t \to \varepsilon} L(t), \\ &= L(t) L^{-1}(\varepsilon) S_{\gamma_0(\varepsilon)} L(\varepsilon), \\ &= L(t) \lim_{\varepsilon \to 0^+} L^{-1}(\varepsilon) S_{\gamma_0(\varepsilon)} L(\varepsilon), \\ &= L(t) \lim_{\varepsilon \to 0^+} e^{-x_0 \log \varepsilon} e^{2i\pi x_0} e^{x_0 \log \varepsilon} \\ &= L(t) e^{2i\pi x_0}. \end{split}$$

ullet Monodromy of L(z) around 1

$$\begin{split} \mathcal{M}_1 L(t) &= S_{1-\varepsilon \to t} S_{\gamma_1(\varepsilon)} S_{t \to 1-\varepsilon} L(t), \\ &= L(t) L^{-1} (1-\varepsilon) S_{\gamma_1(\varepsilon)} L(1-\varepsilon), \\ &= L(t) \lim_{\varepsilon \to 0^+} Z^{-1} e^{x_1 \log \varepsilon} e^{-2i\pi x_1} e^{-x_1 \log \varepsilon} Z \\ &= L(t) Z^{-1} e^{-2i\pi x_1} Z. \end{split}$$

Using the expression of Z and the formula $e^a e^b e^{-a} = e^{a \operatorname{d}_a b}$, we get finally the expression for \mathfrak{m}_1 .

The series \mathfrak{m}_1 up to order 6 by computer

$$\begin{array}{ll} \mathfrak{m}_{1} &=& -[x_{1}] + \zeta(x_{0}x_{1})[x_{0}x_{1}^{2}] + \zeta(x_{0}^{2}x_{1})[x_{0}^{2}x_{1}^{2}] \\ &+ \zeta(x_{0}x_{1}^{2})[x_{0}x_{1}^{3}] + \zeta(x_{0}^{3}x_{1})[x_{0}^{2}x_{1}^{2}] \\ &- \zeta(x_{0}^{3}x_{1})[x_{0}^{2}x_{1}x_{0}x_{1}] + \zeta(x_{0}^{2}x_{1}^{2})[x_{0}^{2}x_{1}^{3}] \\ &+ (\zeta(x_{0}^{2}x_{1}^{2}) - \frac{1}{2}\zeta(x_{0}x_{1})^{2})[x_{0}x_{1}x_{0}x_{1}^{2}] \\ &+ \zeta(x_{0}x_{1}^{3})[x_{0}x_{1}^{4}] + \zeta(x_{0}^{4}x_{1})[x_{0}^{4}x_{1}^{2}] \\ &- 2\zeta(x_{0}^{4}x_{1})[x_{0}^{3}x_{1}x_{0}x_{1}] + \zeta(x_{0}^{3}x_{1}^{2})[x_{0}^{3}x_{1}^{3}] \\ &+ (3\zeta(x_{0}^{3}x_{1}^{2}) + \zeta(x_{0}^{2}x_{1}x_{0}x_{1}))[x_{0}^{2}x_{1}x_{0}x_{1}^{2}] \\ &+ (3\zeta(x_{0}^{3}x_{1}^{2}) + \zeta(x_{0}x_{1})\zeta(x_{0}^{2}x_{1}) \\ &+ 2\zeta(x_{0}^{2}x_{1}x_{0}x_{1}))[x_{0}^{2}x_{1}^{2}x_{0}x_{1}] + \zeta(x_{0}^{2}x_{1}^{3})[x_{0}^{2}x_{1}^{4}] \\ &+ (4\zeta(x_{0}^{2}x_{1}^{3}) + \zeta(x_{0}x_{1}x_{0}x_{1}^{2}))[x_{0}x_{1}x_{0}x_{1}^{3}] \\ &+ \zeta(x_{0}x_{1}^{4})[x_{0}x_{1}^{5}] \end{array}$$

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Monodromy around z=1 (for $p=2i\pi$)

$$\begin{array}{rcl} \mathcal{M}_1 \operatorname{Li}_{x_0} &=& \operatorname{Li}_{x_0} \\ \mathcal{M}_1 \operatorname{Li}_{x_1} &=& \operatorname{Li}_{x_1} - p \\ \mathcal{M}_1 \operatorname{Li}_{x_0 x_1} &=& \operatorname{Li}_{x_0 x_1} - p \operatorname{Li}_{x_0} \\ \mathcal{M}_1 \operatorname{Li}_{x_0^2 x_1} &=& \operatorname{Li}_{x_0^2 x_1} - \frac{p}{2} \operatorname{Li}_{x_0}^2 \\ \mathcal{M}_1 \operatorname{Li}_{x_0 x_1^2} &=& \operatorname{Li}_{x_0 x_1^2} - p \operatorname{Li}_{x_0 x_1} + \frac{p^2}{2} \operatorname{Li}_{x_0} + p \zeta(x_0 x_1) \\ \mathcal{M}_1 \operatorname{Li}_{x_0^3 x_1} &=& \operatorname{Li}_{x_0^3 x_1} - \frac{p}{6} \operatorname{Li}_{x_0^3}^3 \\ \mathcal{M}_1 \operatorname{Li}_{x_0^2 x_1^2} &=& \operatorname{Li}_{x_0^2 x_1^2} - p \operatorname{Li}_{x_0^2 x_1} + \frac{p^2}{4} \operatorname{Li}_{x_0}^2 \\ &&& + p \zeta(x_0 x_1) \operatorname{Li}_{x_0} + p \zeta(x_0^2 x_1) \\ \mathcal{M}_1 \operatorname{Li}_{x_0 x_1^3} &=& \operatorname{Li}_{x_0 x_1^3} - p \operatorname{Li}_{x_0 x_1^2} + \frac{p^2}{2} \operatorname{Li}_{x_0 x_1} \\ &&&& - \frac{p^3}{6} \operatorname{Li}_{x_0} + p \zeta(x_0 x_1^2) - \frac{p^2}{2} \zeta(x_0 x_1) \\ \mathcal{M}_1 \operatorname{Li}_{x_0^4 x_1} &=& \operatorname{Li}_{x_0^4 x_1} - \frac{p}{24} \operatorname{Li}_{x_0}^4 \\ &&&& \vdots \end{array}$$

Structure of the monodromy group

Corollary 7 Monodromy of Li_w is given by

$$\forall w \in X^*, \quad \mathcal{M}_0 \sqcup_{wx_0} = \sqcup_{wx_0} + 2i\pi \sqcup_w + \cdots$$

$$\mathcal{M}_1 \sqcup_{wx_1} = \sqcup_{wx_1} - 2i\pi \sqcup_w + \cdots,$$

The remaining terms are combinations of polylogarithms encoded by words of length < |w|.

Proof - The monodromy theorem implies

$$M_0=e^{2i\pi\mathfrak{m}_0}=1+2i\pi x_0+$$
 words of length >1 $M_1=e^{2i\pi\mathfrak{m}_1}=1-2i\pi x_1+$ words of length >1

Corollary 8 The monodromy group of Li_w for $|w| \leq n$ is nilpotent at order n+1.

Proof $-M_0=e^{2i\pi x_0}$ and $M_1=e^{-2i\pi x_1+\cdots}$. From $e^Ae^Be^{-A}e^{-B}=e^{[A,B]+\cdots}$, it follows that the commutator $M_0M_1M_0^{-1}M_1^{-1}$ does not contain any Lie brackets of length 1. Iterating this computation, the brackets of lengths 2, next 3, etc. until n disappear. \square

A structure theorem

Theorem 7 (FPSAC98) The polylogarithms are lenearly indepedant.

Proof – This is trivial for n=0. Assume that we have proved our assertion for all k, 0 < k < n-1. For k=n,

$$\begin{split} & \sum_{|w| \leqslant n} \lambda_w \operatorname{Li}_w = \mathbf{0} \\ \iff & \lambda_1 + \sum_{|u| < n} \lambda_{ux_0} \operatorname{Li}_{ux_0} + \sum_{|u| < n} \lambda_{ux_1} \operatorname{Li}_{ux_1} = \mathbf{0}. \end{split}$$

(the λ_w are elements of \mathbb{C}). Applying $(\mathcal{M}_0 - Id)$ and $(Id - \mathcal{M}_1)$, we have

$$\begin{cases} 2i\pi \sum_{|u|=n-1} \lambda_{ux_0} \operatorname{Li}_u + \sum_{|u|< n-1} \mu_u \operatorname{Li}_u &= 0, \\ 2i\pi \sum_{|u|=n-1} \lambda_{ux_1} \operatorname{Li}_u + \sum_{|u|< n-1} \nu_u \operatorname{Li}_u &= 0. \end{cases}$$

By the induction hypothesis, we get the expected result. $\hfill\Box$

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 $\mathsf{Li}_s(1-t)$ by computer

$$\begin{array}{rcl} \operatorname{Li}_{1}(1-t) &=& -\log{(t)} \\ \operatorname{Li}_{2}(1-t) &=& -\operatorname{Li}_{2}(t) + \log(t)\operatorname{Li}_{1}(t) + \zeta(2) \\ \operatorname{Li}_{3}(1-t) &=& -\operatorname{Li}_{2,1}(t) + \operatorname{Li}_{1}(t)\operatorname{Li}_{2}(t) \\ && -\frac{1}{2}\log(t)\operatorname{Li}_{1}(t)^{2} \\ && -\zeta(2)\operatorname{Li}_{1}(t) + \zeta(3) \\ \operatorname{Li}_{2,1}(1-t) &=& -\operatorname{Li}_{3}(t) + \log(t)\operatorname{Li}_{2}(t) \\ && -\frac{1}{2}\log{(t)^{2}}\operatorname{Li}_{1}(t) + \zeta(3) \\ \operatorname{Li}_{4}(1-t) &=& -\operatorname{Li}_{2,1,1}(t) + \operatorname{Li}_{1}(t)\operatorname{Li}_{2,1}(t) \\ && -\frac{1}{2}\operatorname{Li}_{1}(t)^{2}\operatorname{Li}_{2}(t) \\ && +\frac{1}{6}\log(t)\operatorname{Li}_{1}(t)^{3} + \frac{1}{2}\zeta(2)\operatorname{Li}_{1}(t)^{2} \\ && -\zeta(3)\operatorname{Li}_{1}(t) + \frac{2}{5}\zeta(2)^{2} \end{array}$$

$$L(1-t)$$

Proposition 4 (FPSAC98) For any $t \in]0,1[$,

$$L(1-t) = g_*[L(t)]Z,$$

and g_* is defined by $g_*x_0 = -x_1, g_*x_1 = -x_0$.

Proof - (sketched) One has firstly

$$L(1-t) = S_{1-\varepsilon \leadsto 1-t} L(1-\varepsilon)$$

$$= g_* S_{\varepsilon \leadsto t} L(1-\varepsilon)$$

$$= g_* [L(t) L^{-1}(\varepsilon)] L(1-\varepsilon)$$

$$= g_* L(t) g_* L^{-1}(\varepsilon) L(1-\varepsilon).$$

and secondly

$$L(1-t) \sim g_*L(t) g_*e^{-x_0\log\varepsilon} e^{-x_1\log\varepsilon}Z.$$

If g(t)=1-t then $g_*\omega_0=-\omega_1$ and $g_*\omega_1=-\omega_0$. Therefore $g_*x_0=-x_1$ and $g_*x_1=-x_0$. Hence it follows the expected result. \square

Corollary 9 Let g_* be the morphism defined by $g_*x_0 = -x_1$ and $g_*x_1 = -x_0$. Then

$$L\left(\frac{1}{2}\right) = g_* \left[L\left(\frac{1}{2}\right)\right] Z.$$

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Duality relation

Proposition 5 Let τ be the composition of the mirror morphism and of the involutive substitution morphism $x_0 \to x_1$ and $x_1 \to x_0$. Then

$$Z = \tau(Z)$$
.

Proof – For $t \in]0,1[$, one has

$$S_{t \to 1-t}(x_0, x_1) = S_{1-t \to t}(-x_1, -x_0)$$

= $S_{t \to 1-t}^{-1}(-x_1, -x_0)$
= $\tau[S_{t \to 1-t}(x_0, x_1)].$

By the renormalisation

$$S_{t \leadsto 1-t} \sim e^{-x_1 \log t} Z e^{-x_0 \log t}, \ \text{for} \ t \to 0^+,$$
 and then

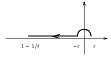
$$\tau(S_{t \sim 1-t}) \sim e^{-x_1 \log t} \tau(Z) e^{-x_0 \log t}$$
, for $t \to 0^+$, we get the expected result. \square

$$L(1 - 1/t)$$

Proposition 6 (FPSAC98) For any $t \in]0,1[$,

$$L(1-1/t) = g_*[L(t)]g_*(Z^{-1})e^{i\pi x_0},$$

and g_* is defined by $g_*x_0 = -x_0 + x_1, g_*x_1 = x_1$.



Proof $-L(1-1/t)=S_{-\varepsilon \leadsto 1-1/t}S_{\varepsilon \leadsto -\varepsilon}L(\varepsilon)=S_{-\varepsilon \leadsto 1-1/t}e^{i\pi x_0}e^{x_0\log \varepsilon}.$ For g(t)=1-1/t then $g_*\omega_0=-\omega_0+\omega_1$ and $g_*\omega_1=-\omega_0.$ This leads to $g_*x_0=-x_0+x_1$ and $g_*x_1=-x_0.$ Thus

$$S_{-\varepsilon \to 1-1/t} = g_* S_{1-\varepsilon \to t} = g_* (L(t)L^{-1}(1-\varepsilon))$$

= $g_* (L(t)Z^{-1}e^{x_1 \log \varepsilon}).$

Corollary 10 For any $w \in X^*$, for $\varepsilon \to 0^+$,

$$\mathsf{Li}_w(-1/arepsilon) \sim rac{(-1)^{|w|_{x_0}}}{|w|!} \mathsf{log}^{|w|}(arepsilon).$$

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$\operatorname{Li}_s(1-1/t)$ by computer

$$\begin{array}{rcl} \log(1-1/t) &=& (i\pi)-\mathrm{Li}_1(t)-\log(t) \\ \mathrm{Li}_1(1-1/t) &=& \log(t) \\ \mathrm{Li}_2(1-1/t) &=& \mathrm{Li}_2(t)-\log(t)\mathrm{Li}_1(t) \\ && -\zeta(2)-\frac{1}{2}\mathrm{log}(t)^2 \\ \mathrm{Li}_3(1-1/t) &=& \mathrm{Li}_{2,1}(t)-\mathrm{Li}_3(t)-\mathrm{Li}_1(t)\mathrm{Li}_2(t) \\ && +\frac{1}{2}\mathrm{log}(t)\mathrm{Li}_1(t)^2 \\ && +\left(\zeta(2)+\frac{1}{2}\mathrm{log}(t)^2\right)\mathrm{Li}_1(t) \\ && +\mathrm{log}(t)\zeta(2)+\frac{1}{6}\mathrm{log}(t)^3 \\ \mathrm{Li}_{2,1}(1-1/t) &=& -\mathrm{Li}_3(t)+\mathrm{log}(t)\mathrm{Li}_2(t) \\ && -\frac{1}{2}\mathrm{log}(t)^2\mathrm{Li}_1(t) \\ && +\zeta(3)-\frac{1}{6}\mathrm{log}(t)^3 \end{array}$$

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Hexagonal relation

Proposition 7 Let ρ be the substitution morphism $x_0 \to x_1$ and $x_1 \to x_0$. Then

$$Ze^{i\pi x_0}\rho(Z)e^{i\pi(-x_0+x_1)}\rho^2(Z)e^{-i\pi x_1}=1.$$

Proof – Let g(z)=1-1/z permuting the singularities 0, 1 and ∞ . Then $g_*\omega_0=-\omega_0+\omega_1$ and $g_*\omega_1=-\omega_0$. This leads to $g_*x_0=-x_0+x_1$ and $g_*x_1=-x_0$. Thus

$$S_{\varepsilon \to 1-\varepsilon} e^{i\pi x_0} g_* (S_{\varepsilon \to 1-\varepsilon} e^{i\pi x_0}) g_*^2 (S_{\varepsilon \to 1-\varepsilon} e^{i\pi x_0}) = 1.$$



By the renormalisation

$$S_{\varepsilon \to 1-\varepsilon} \sim e^{-x_1 \log \varepsilon} Z e^{-x_0 \log \varepsilon}$$
, for $\varepsilon \to 0^+$, we get the expected result. \square

By Campbell-Baker-Hausdorf formula, one has

Corollary 11 $\zeta(2) = \pi^2/6$.

Drinfel'd associator $\Phi_{KZ}(A,B)$ and non-commutative g.s. of polyzetas

By changing

$$x_0 := rac{A}{2i\pi}$$
 and $x_1 := -rac{B}{2i\pi},$

we have

$$\Phi_{KZ}(A,B) \equiv Z(x_0,x_1)$$

Thus

$$\log \Phi_{KZ}(A, B) = \frac{1}{24}[A, B]$$

$$+ \frac{\zeta(3)}{(2i\pi)^3}([[A, B], B] - [A, [A, B]])$$

$$+ \frac{1}{1440}([[[A, B], B], B] - [A, [A, [A, B]]])$$

$$+ \frac{1}{4}[A, [[A, B], B]]) + \cdots$$

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