

1, 2, 3; a, b, c; ...

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Diophantine equations

Basically:

- Unknowns are integers, or rational numbers
- Equations are given by polynomial relations between the unknowns.

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Roughly, three parameters:

- The number of variables;
- The number of equations;
- The degree of the equations.

Warm-up

Diophantine equations in *one* variable are easy to solve.

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Theorem

Let $a_0, \dots, a_{n-1}, a_n \in \mathbf{C}$, with $a_n \neq 0$. Every root x of the n th degree equation:

$$a_n x^n + \dots + a_1 x + a_0 = 0$$

satisfies

$$|x| \leq \sup \left(1, \frac{|a_0| + \dots + |a_{n-1}|}{|a_n|} \right).$$

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If we seek for integer solutions, it then suffices to try one by one all integers in the interval that is described by the theorem.

Geometry of diophantine equations

Equations of degree 1

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This “Chinese problem” is due to Sūnzǐ, and was published in the Sūnzǐ Suàngjīng 孫子算經, *The Mathematical Classic of Master Sun*



孫子算經卷上

唐劉義慶行參上輕軍都尉景淳撰算奉勅注釋



度之所起起於忽欲知其忽蠶吐絲為忽十忽
 為一絲十絲為一毫十毫為一釐十釐為一分
 十分為一寸十寸為一尺十尺為一丈十丈為
 一引五十尺為一端四十尺為一疋六尺為一
 步二百四十步為一畝三百步為一里
 稱之所起起於黍十黍為一絫十絫為一銖二
 十四銖為一兩十六兩為一斤三十斤為一鈞

Sūnzǐ Suànjīng,
 reproduction of a page
 from a Qing dynasty edition

Source: *Wikipedia*

The Chinese problem

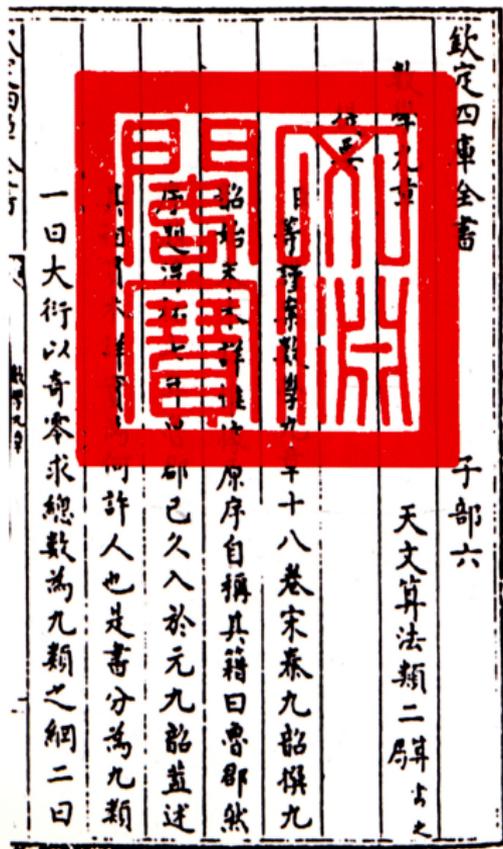
Such exercises were then reproduced in other manuals, such as the

Shùshū Jiǔzhāng — 數書九章,
Mathematical treatise in nine sections, 1247,

itself included in the

Sìkù quánhū — 四庫全書,
Complete Library of the Four Treasuries, XIXth c.,

a kind of encyclopaedia commissioned by the Qing emperors to attest their supremacy over the former Ming encyclopaedia (ca. 1403)



*A Mathematical Book in
Nine Chapters*
(數書九章, Shùshū
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Source: *Wikipedia*

Le problème chinois

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Or,

again, solve

$$n = 3x + 2 = 5y + 3 = 7z + 2,$$

in the main unknown n .

The “Chinese theorem” teaches us that the smallest solution is $n = 23$, and that all other solutions are obtained by adding a multiple of $105 = 3 \times 5 \times 7$.

Πρόβλημα,

ὄπερ Ἀρχιμήδης ἐν ἐπιγράμμασιν εὐφών τοῖς ἐν Ἀλεξανδρείᾳ περὶ ταῦτα πραγματευομένοις ζητεῖν ἐπίστειλεν ἐν τῇ πρὸς Ἐρατοσθένην τὸν Κυρηναῖον ἐπιστολῇ.

- 1 Πληθὺν Ἡελίοιο βοῶν, ᾧ ξεῖνε, μέτρησον
φροντίδ' ἐπιστήσας, εἰ μετέχεις σοφίης,
πόσση ἄρ' ἐν πεδίοις Σικελῆς ποτ' ἐβόσκετο νήσου
Θρινακίης τετραχῆ στίφεια δασσαμένη
- 5 χροίην ἀλλάσσοντα· τὸ μὲν λευκοῖο γάλακτος,
κυανέω δ' ἕτερον χρώματι λαμπόμενον,
ἄλλο γε μὲν ξανθόν, τὸ δὲ ποικίλον. ἐν δὲ ἐκάστῳ
στίφει ἔσαν ταῦροι πλήθει βριθόμενοι
συμμετρίας τοιῆσδε τετευχότες· ἀργότριχας μὲν
- 10 κυανέων τὰύρων ἡμίσει ἠδὲ τρίτῳ
καὶ ξανθοῖς σύμπασιν ἴσους, ᾧ ξεῖνε, νόησον,

Equations of degree 2

Πρόβλημα,

ἕκασ' Ἀρχιμήδης ἐν ἐπιγράμματι κτῶν τοῖς ἐν' Ἀλεξανδρείᾳ παρὶ τούτου πραγματευομένου ὄντιν ἀπέστειλεν ἐν τῇ πρὸς Ἐρατοσθένην τὸν Κυρηνεὸν ἐπιστολῇ.

1 Πληθὺν Ἥλιου βοῶν, ἃ ἔειν, μέγιστον
φρονεῖν ἑαυτοῦ, εἰ μέγιστος σφῶν·
πίσις ἥ' ἐν καθόλου Σικελίᾳ κατ' ἑβέτατο νήσου
Θρινακίᾳ τετραπλῆ στίραι δασυμένη
5 χωρὶν ἀλλήλοισιν· τὸ μὴ λευκοὶ γάλατος,
κυνῶν δ' ἕτερον χροῖμα λαμπρόν,
ἄλλο γὰρ μὴν ἔκον, τὸ δὲ κοινὸν. ἐν δὲ ἑσῆσιν
στίραι ἴσων ταύρων πλήθει· βριθόμενοι
συμμετρίας ταύρων τετραπλοῦς· ἀργύρευχερ μὴν
10 κυνῶν ταύρων ἡμίσει ἤδη τρίτη
καὶ ἑνθόλις στίραισιν ἴσους, ἃ ἔειν, νόησον,
αὐτῶν κυνῶν τῇ τετραπλῆ τὸ μῆκος
μικτοῦσιν καὶ πέμπτῃ, ἔτι ἑνθόλις τε καὶ
τοῖς δ' ἑπολειομένους κοινολόγησας ἔθρει
15 ἀργυρῶν ταύρων ἄτερ μῆκος ἰσομήτρῃ τε
καὶ ἑνθόλις αὐτῶν πᾶσιν ἰσομήτρους.
Θηλείαις δὲ βοῦσι τῶν ἑκλετο' λευκότευχερ μὴν
ἴσων συμπίσις κυνῶν ἀγέλης
τῇ τετραπλῆ τὸ μῆκος καὶ τετραπλῆ ἀτραχὺ ἴσων·
20 αὐτῶν κυνῶν τῇ τετραπλῆ τὸ πᾶσι
μικτοῦσιν καὶ πέμπτῃ ἰσοῦ μῆκος ἰσίζοντο
σὺν ταύροις· πίσις δ' εἰς νομὸν ἀρχαίαν
ἑνθόριον ἀγέλης πέμπτῃ μῆκος ἤδη καὶ ἄτερ
κοινῶν ἰσομήτρους πλῆθος ἔχειν τετραπλῆ.
25 ἑνθόλι δ' ἡμιβρωτῶν μῆκος τρίτου ἡμίσει ἴσων
ἀργυρῶν ἀγέλης ἰσομήτρῃ τὸ μῆκος.
ἔειν, οὐ δ' Ἥλιου βοῶν πίσις ἀτραχὺ εἰσῆν,
χωρὶς μὴν ταύρων ἑταίρων ἀφθῶν,
χωρὶς δ' οὐ θήλειαι ἴσων κατὰ χροῖμα ἕσονται,
30 οὐκ ἄνδρῶν κε λέγει· οὐδ' ἀριθμῶν ἀσθηῆς,
οὐ μὴν καὶ γὰρ σοφοῦ ἐναριθμοῦ. ἀλλ' ἴθι φράξαι
καὶ τὰς πᾶσαι βοῶν Ἥλιου πᾶσιν.
ἀργύρευχερ ταύρων μὴν ἑσῆ μάλιστα πληθύν
κυνῶν, ἴσων' ἡμῶν ἰσομήτρους
35 εἰς βῆθος εἰς κτῶν τε, τὸ δ' οὐ περιήρημα πᾶσιν
κίρῳλετο πλῆθος Θρινακίᾳ πᾶσι.
ἑνθόλι δ' αὐτ' εἰς ἔν καὶ κοινῶν ἀποσπῶντις
ἴσων' ἀμβολίδων ἔξ ἑνοῦ ἀρχόμενοι
σχῆμα τελειοῦντις τὸ τετραπλοῦσιν οὐκ προσόντων
40 ἀλλοτροῦσιν ταύρων οὐκ ἑπολειομένων.
ταῦτα συζητηρῶν καὶ ἐνὶ πραγμασίῃσιν ἀφροῦσις
καὶ κληθῶν ἀποδοῦς, ἃ ἔειν, πᾶσι μέγας
ἔργο πᾶσιον νικηφόρος, ἴσθι τε πᾶσι
κεκρίμενοι ταύτῃ ἡμῶν ἐν σοφῇ.

Archimedes's cattle problem
Archimedis Opera omnia, cum
commentariis Eutocii
Edited by J. L. Heiberg
B. G. Teubner, Leibzig, Volume 2
(1881), pp. 448–450

The Brahmagupta (Pell–Fermat) equation

If n is a (non-square) parameter, find the solutions in rational integers to the equation

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In Archimedes's problem: $n = 4 \times 609 \times 7766 \times 4657^2 \dots$

Solving the Brahmagupta equation

Brahmagupta (628 c.E.): if (x, y) and (x', y') are solutions, one may build a third one (x'', y'') by the formula:

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All solutions (more or less) are obtained from a minimal one.

Modern explanation :

$$x^2 - ny^2 = (x + \sqrt{ny})(x - \sqrt{ny}),$$

is the *norm* of the quadratic number $x + \sqrt{ny}$; it is multiplicative and one has

$$(x + \sqrt{ny})(x' + \sqrt{ny}') = (xx' + nyy') + \sqrt{n}(xy' + x'y).$$

The Pythagoras equation

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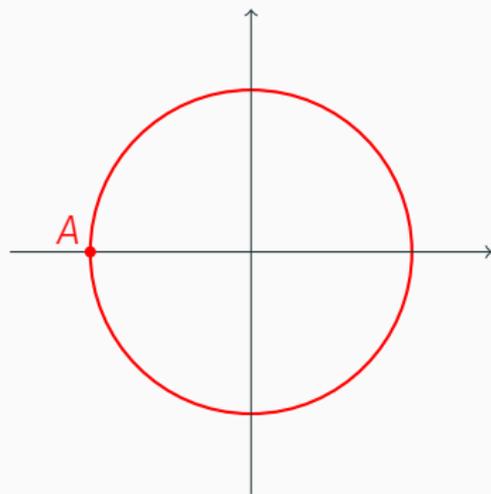
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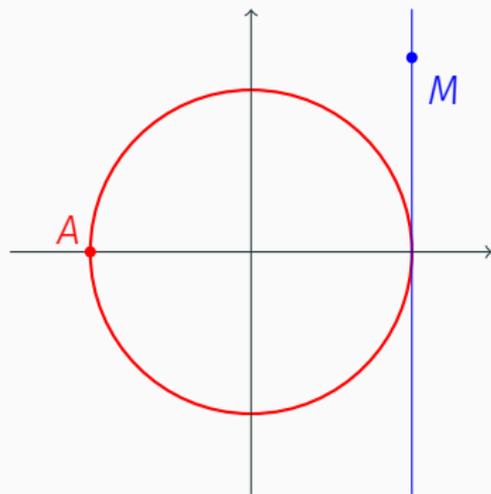
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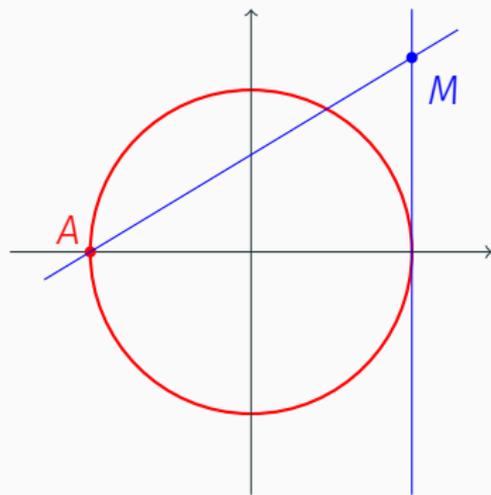
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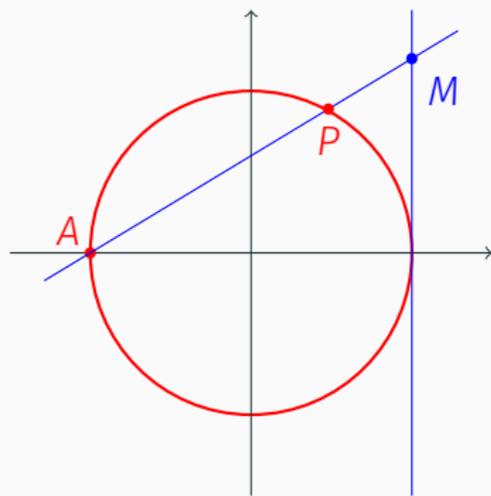
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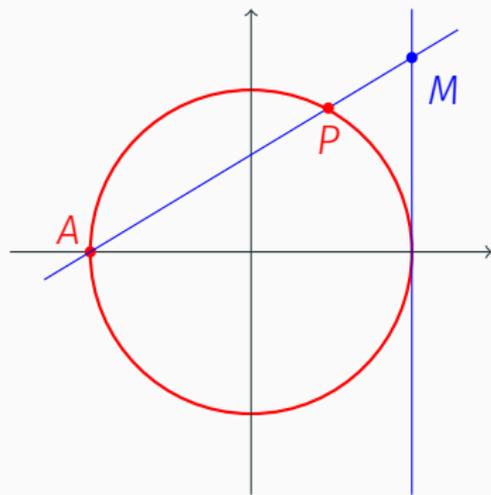
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If M has coordinates $(1, t)$,
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with

$$\begin{cases} x = \frac{1-t^2}{1+t^2} \\ y = \frac{2t}{1+t^2} \end{cases}$$

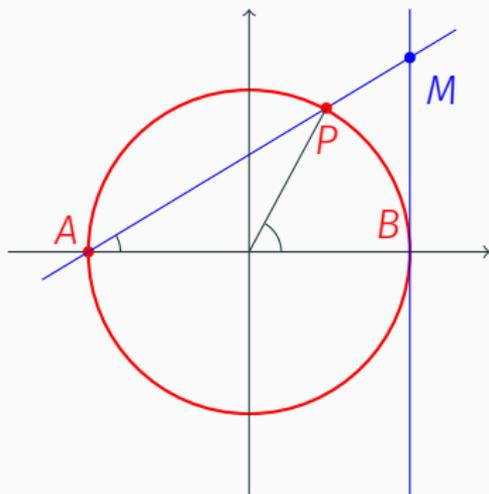
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$$t = \arctan(\widehat{BAP}) = \arctan(\widehat{BOP}/2)$$

Rational parameterizing of conics

The procedure explained for the circle works for every *conic* (given by a degree 2 equation in two variables) *provided* there exists at least one solution in rational numbers.

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More generally: *If a quadratic form $q(x, y, z, \dots)$ with coefficients in \mathbf{Q} has non-trivial solutions in all p -adic fields \mathbf{Q}_p , as well as in \mathbf{R} , then it has a non-trivial solution*

QVÆSTIO VIII.

PROPOSITVM quadratum diuidere in duos quadratos. Imperatum fit vt 16. diuidatur in duos quadratos. Ponatur primus 1 Q. Oportet igitur $16 - 1$ Q. æquales esse quadrato. Fingo quadratum à numeris quotquot libuerit, cum defectu tot unitatum quod continet latus ipsius 16. esto à 2 N. - 4. ipse igitur quadratus erit 4 Q. + 16. - 16 N. hæc æquabuntur unitatibus 16 - 1 Q. Communis adiiciatur vtriusque defectus, & à similibus auferantur similia, fient 5 Q. æquales 16 N. & fit 1 N. $\frac{16}{5}$ Erit igitur alter quadratorum $\frac{16}{5}$. alter verò $\frac{144}{25}$ & vtriusque summa est $\frac{176}{25}$ seu 16. & vterque quadratus est.

ἢ εἰκοσὸπμπτζ, ἦτοι μνάδας 15. καὶ ἔστιν ἐκείνους τετράγωνοι.

TON ἑπιτετραγώνου τετραγώνου διελὼν εἰς δύο τετραγώνους. ἐπιτετράγωνο δὴ τὸ 15² διελὼν εἰς δύο τετραγώνους. καὶ τετράγωνο ὁ πρῶτος διωάμειος μνάς. δέησει ἄρα μονάδας 15² λείψει διωάμειος μνάς 16ας ἢ τετραγώνου. πλάσσω τὸ τετράγωνον ἀπὸ 5. ὅσων δὴ ποτε λείπει ποσῶτων μ² ὅσων ὅστιν ἢ τὸ 15² μ² πλάσσω. ἔστω 5² β² λείψει μ² δ². αὐτὸς ἄρα ὁ τετράγωνος ἔσται διωάμειος δ² μ² 15² λείψει 5² 15². ταῦτα ἴσα μονάσει 15² λείψει διωάμειος μνάς. κοινὴ πρῶτος εἰδὼ ἢ λείψει. καὶ ἀπὸ ὁμοίων ὁμοία. διωάμειος ἄρα ἔσται ἀβθμοῖς 15². καὶ γίνεται ὁ ἀβθμοῖς 15². πέντε πέντων. ἔσται ὁ ἀβθμοῖς εἰκοσὸπμπτων. ὁ δὲ μνάς εἰκοσὸπμπτων. ἔσται οἱ δύο συντεθέντες ποσῶσι.

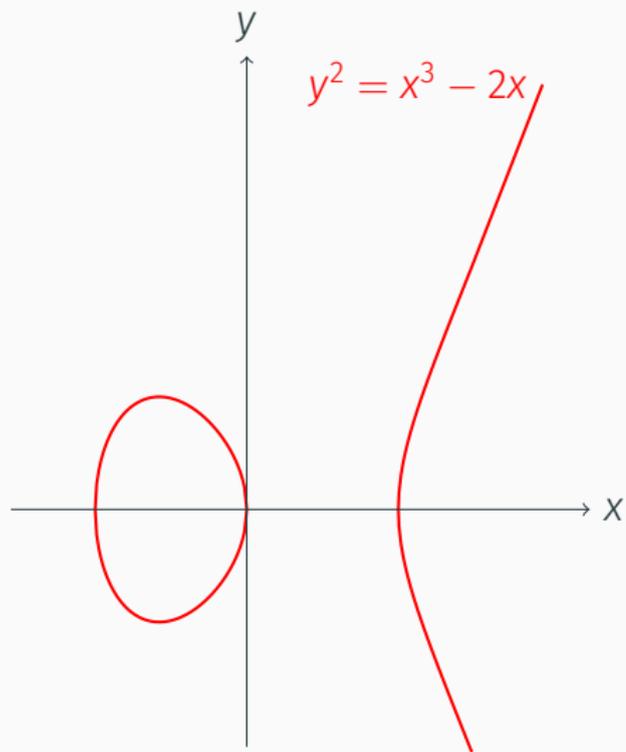
OBSERVATIO DOMINI PETRI DE FERMAT.

Cubum autem in duos cubes, aut quadratoquadratum in duos quadratoquadratos & generatim nullam in infinitum ultra quadratum potestatem in duos eiusdem nominis fas est diuidere cuius rei demonstrationem mirabilem sane detexi. Hanc marginis exiguitas non caperet.

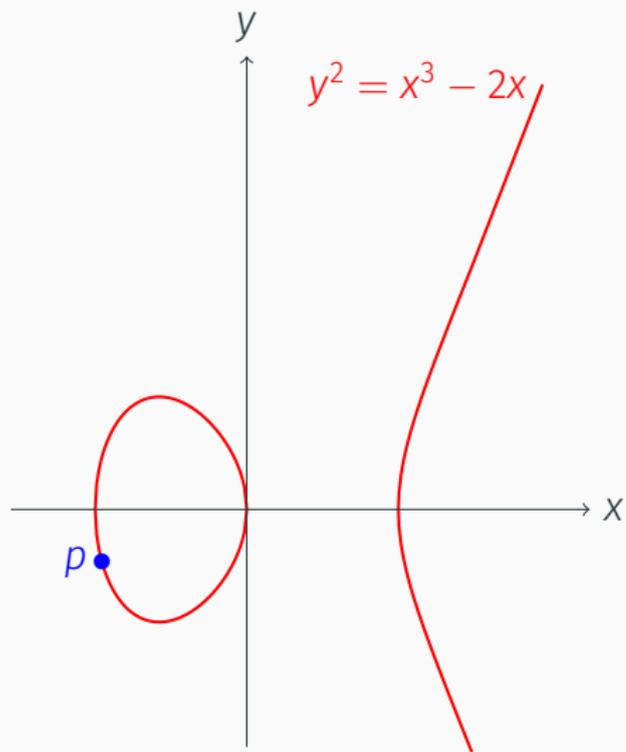
Diophante, *Arithmetica*. Bachet de Méziriac edition, 1670.

Source: Wikipedia

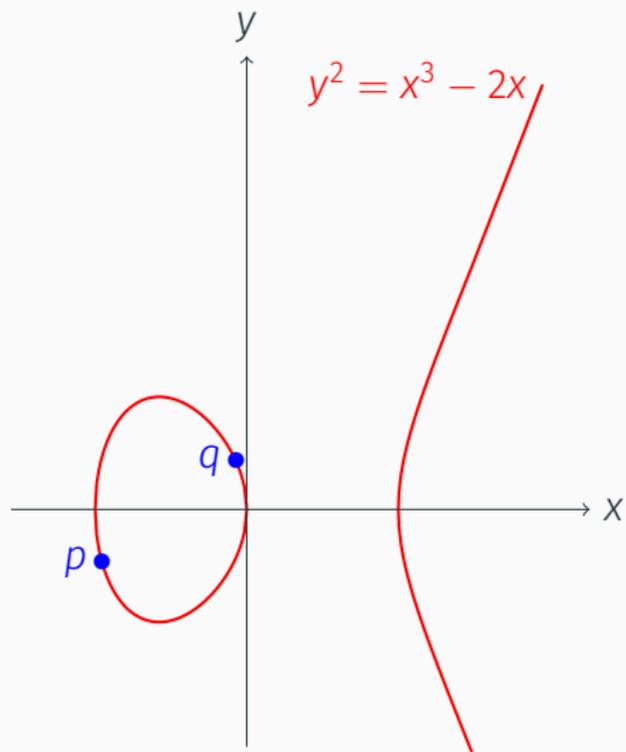
A cubic equation



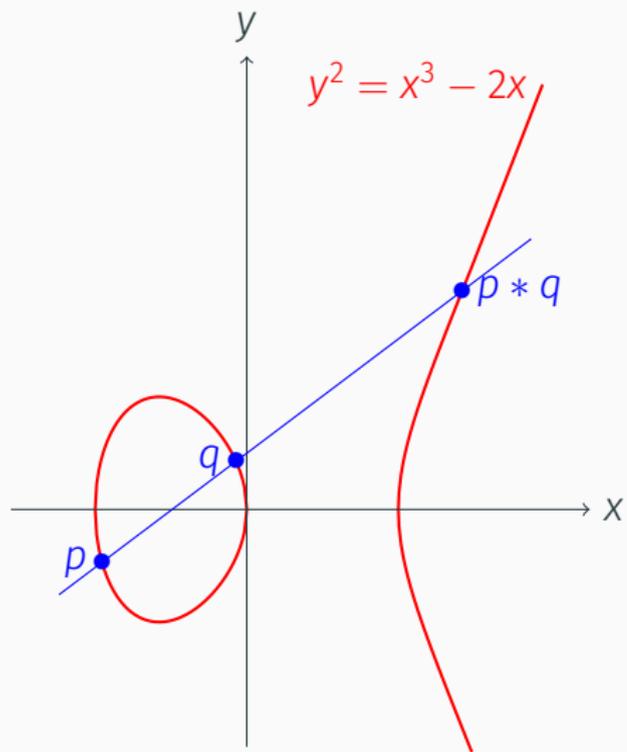
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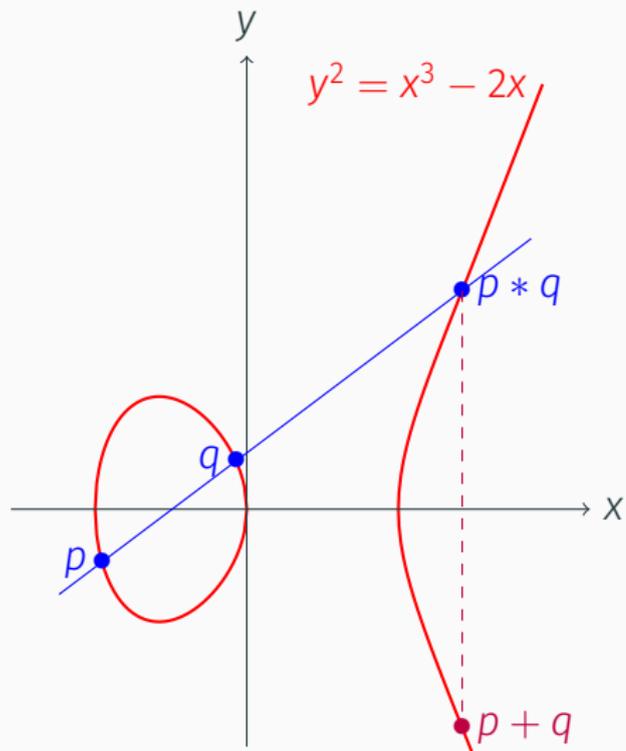
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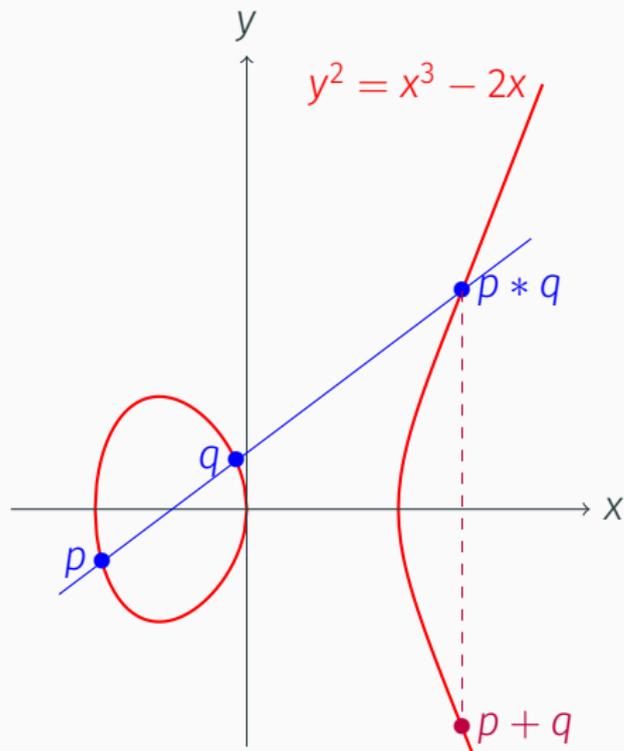
A cubic equation



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A cubic equation



This defines a
commutative group law
on the set of rational
solutions.

The neutral element is
the “point at infinity”;

Elliptic curves

Elliptic curves are those curves defined by a cubic equation of the form

$$f(x, y) = y^2 - x^3 - ax - b = 0, \quad \Delta = -4a^3 - 27b^2 \neq 0.$$

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The condition on the discriminant Δ states that the curve is *non singular*: if (x, y) is a singular point,

$$\frac{\partial}{\partial x} f(x, y) = \frac{\partial}{\partial y} f(x, y) = 0$$

implies that $y = 0$ and x is a multiple root of $x^3 + ax + b$.

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Contrary to conics, elliptic curves cannot be parameterized by rational functions.

Elliptic curves — Another Poincaré conjecture

Let an elliptic be given by a cubic equation of the form

$$y^2 = x^3 + ax + b, \quad \Delta = -4a^3 - 27b^2 \neq 0$$

where a and b are rational numbers.

Theorem (Mordell, 1922)

The group of rational solutions is a finitely generated abelian group.

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Theorem (Mordell, 1922)

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Theorem (Siegel, 1929)

There are only finitely many integer solutions.

Equations of higher degree

Let us now consider an equation of degree ≥ 1 , defining a *nonsingular curve*. It is important to work in the context of *projective geometry* and to forbid singularities at infinity, or complex.

Theorem (Faltings, 1983; conjectured by Mordell)

There are only finitely many rational solutions.

Equations of higher degree

Let us now consider an equation of degree ≥ 1 , defining a *nonsingular curve*. It is important to work in the context of *projective geometry* and to forbid singularities at infinity, or complex.

Theorem (Faltings, 1983; conjectured by Mordell)

There are only finitely many rational solutions.

Consequence : for every integer $n \geq 4$, the Fermat equation has only finitely many solutions.

Geometric trichotomy

Up to now, we saw three classes of equations:

- degree 1 or 2 (conics): rational parameterizations, sometimes infinitely many integer solutions;
- degree 3 (elliptic curves) : no rational parameterization, sometimes infinitely many rational solutions, finitely many integer solutions;
- degree 4 or higher : finitely many rational solutions.

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The correct way of understanding this trichotomy requires to consider the **complex solutions** — they form a Riemann surface and the distinction is then

- genus 0 (Riemann sphere, positive curvature);
- genus 1 (zero curvature);
- genus 2 or higher (negative curvature).

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Only (?) known cases: subvarieties of abelian varieties (Faltings, 1991).

Deciding the solvability of a diophantine equation

Hilbert's 10th problem

David Hilbert, 1900, International congress of mathematicians

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Entscheidung der Lösbarkeit einer diophantischen Gleichung.

Eine diophantische Gleichung mit irgendwelchen Unbekannten und mit ganzen rationalen Zahlkoeffizienten sei vorgelegt: man soll ein Verfahren angeben, nach welchem sich mittels einer endlichen Anzahl von Operationen entscheiden lässt, ob die Gleichung in ganzen rationalen Zahlen lösbar ist.

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On the possibility of solving a diophantine equation.

Let be given a diophantine equation in any number of unknowns and with coefficients in rational integers: one asks to find a method by which, using finitely many operations, one will distinguish whether the equation is solvable in rational integers

The Entscheidungsproblem (Hilbert, 1928)

In 1928, Hilbert generalizes his 10th problem and states **the decision problem** (Entscheidungsproblem): the question is to prove (or disprove) the existence of an algorithm that correctly answers by yes or no **every mathematical question** (suitably formalized, in first order logic).

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1936: Gödel, Turing, Church prove that no such algorithm exists.

But what about diophantine equations?

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Strong version: There exists a polynomial

$$f(t, x_1, \dots, x_9) \in \mathbf{Z}[t, x_1, \dots, x_9]$$

in 10 variables for which no algorithm can tell, given an integer $a \in \mathbf{Z}$, whether or not the equation $f(a, x_1, \dots, x_9) = 0$ has a solution in \mathbf{Z}^9 .

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For rational solutions: the question is still open!

But what about equations in 2 variables?

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In genus ≥ 2 , finding an effective version of Mordell conjecture is a completely open question.

The ABC conjecture

Conjecture (Masser, Oesterlé, 1985)

For every $\theta > 1$, there exists $K_\theta > 0$ so that the following holds:

If A, B, C are three coprime integers such that $A + B = C$, then

$$\max(|A|, |B|, |C|) \leq K_\theta (\text{rad}(ABC))^\theta.$$

The *radical* $\text{rad}(ABC)$ is the product of the prime numbers that divide it.

In other words, this conjecture predicts that the multiplicities of the prime factors of A, B, C are not too large.

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Record (Reyssat, 1987) :

$$2 + 3^{10} \times 109 = 23^5$$
$$\theta(a, b, c) := \frac{\log(\sup(a, b, c))}{\log(\text{rad}(abc))} \approx 1.63$$

Fermat's Last Theorem as a consequence the *ABC* conjecture

Let (x, y, z) be a nontrivial solution ($xyz \neq 0$) of the Fermat equation

$$x^n + y^n = z^n.$$

One may assume that x, y, z are coprime.

Set $A = x^n$, $B = y^n$, $C = z^n$.

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Since $\text{rad}(x^n y^n z^n) = \text{rad}(xyz)$, assuming the truth of the *ABC* conjecture, one has

$$\sup(|x|, |y|, |z|)^n \leq K_\theta (\text{rad}(xyz))^\theta.$$

On the other hand, it is obvious that

$\text{rad}(xyz) \leq |x| |y| |z| \leq \sup(|x|, |y|, |z|)^3$. Consequently,

$$\sup(|x|, |y|, |z|)^n \leq K_\theta \sup(|x|, |y|, |z|)^{3\theta},$$

hence, if $n > 3\theta$,

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The ABC conjecture for polynomials

Theorem (Stothers, 1981; Mason, 1984)

Let $A, B, C \in \mathbf{C}[t]$ be three coprime polynomials such that $A + B = C$. Then

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If x is a root of A , B , or C , with multiplicity m , it is a root of multiplicity $\geq m - 1$ of D , hence

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Then $\deg(C) \leq \nu(ABC) - 1$ and similarly for $\deg(A)$ and $\deg(B)$.

Application to geometric irrationality

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Let $P^n + Q^n = R^n$, for three coprime polynomials $P, Q, R \in \mathbb{C}[t]$.
Then

$$\begin{aligned}n \sup(\deg(P), \deg(Q), \deg(R)) &\leq \nu(PQR) - 1 \\ &\leq \deg(P) + \deg(Q) + \deg(R) - 1 \\ &< 3 \sup(\deg(P), \deg(Q), \deg(R)),\end{aligned}$$

hence $n < 3$.

Theorem (Elkies, 1991)

If the ABC conjecture is true, the “effective” version of Mordell conjecture is true as well: one can give an explicit bound for the size of the solutions.

Principle: Construct a rational function ϕ of x and y whose restriction to the plane curve is unramified everywhere but possibly above $0, 1, \infty$.

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Conversely:

Theorem (Moret-Bailly, Szpiro, 1990)

An effective version of Mordell conjecture would imply the ABC conjecture (with some exponent rather than $1 + \epsilon$).

Proving the *ABC* conjecture?

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August 2018, Peter Scholze, Jakob Stix :

Why abc is still a conjecture