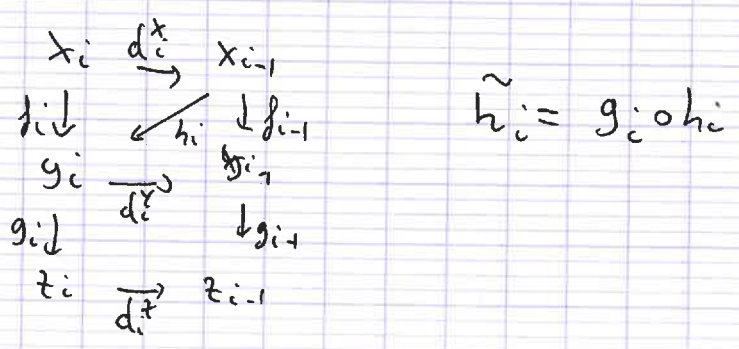


n=3  split as 3 2 matrices

More generally $h_n(\sigma) = \sum_{i=0}^n (-1)^i h_0(\sigma + Id)_{(v_0, \dots, v_i, w_{i+1}, \dots, w_n)}$ □

Lemma 4.20 $f: X_0 \rightarrow Y_0$ $g: Y_0 \rightarrow Z_0$ morphisms of $Ch(E)$. Then $f \circ g \Rightarrow g \circ f \circ g$

Proof



Prop/Def 4.21 E be an additive category. The homotopy category $k(E)$ is the category $ob(k(E)) = ob(Ch(E))$
 $Hom_{k(E)}(x, y) = Hom_{Ch(E)}(x, y) / \sim$

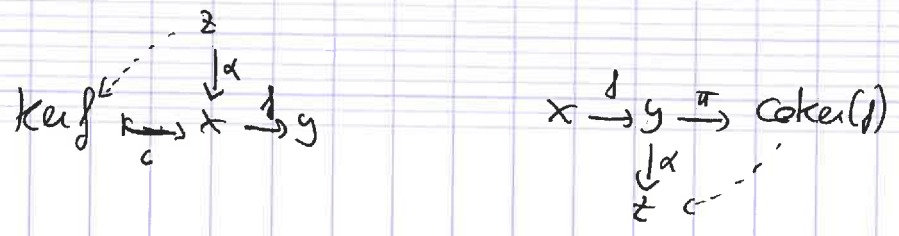
Proof one has to check that this is a category using 4.20 □

Rem (1) $k(E)$ is an additive category
 (2) Δ in general this is a complex object: it is a triangulated category.

Chap 5 Abelian categories

Def 5.1 E be an additive category. A kernel of $f \in Mor(E)$ is an equalizer of $(f, 0)$. Dually a cokernel of f is a coequalizer of $(f, 0)$.

Concretely



We assume that every morphism in \mathcal{E} has a kernel and a cokernel

Then

$$\ker f \xrightarrow{c} X \xrightarrow{f} Y \xrightarrow{\pi} \operatorname{Coker}(f)$$

$$\downarrow p \quad \uparrow$$

$$\operatorname{Coker}(\ker f) \quad \ker(\operatorname{Coker} f)$$

$$f \circ i = 0 \Rightarrow \ker f \rightarrow X \xrightarrow{f} Y \xrightarrow{\pi} \operatorname{Coker} f \quad \text{no } \pi \circ \tilde{f} \circ p = \pi \circ f = 0 \quad (*)$$

$\downarrow \tilde{f}$
Cokernel

lem 5.2

kernels are monomorphism and cokernels are epimorphisms

proof

$$\ker f \xrightarrow{i} X \xrightarrow{f} Y$$

\uparrow \uparrow

$a \uparrow b$ \uparrow

W W

so $\boxed{i \circ (b-a) = 0}$

$$\ker f \xrightarrow{i} X \xrightarrow{f} Y$$

\uparrow \uparrow \uparrow

$b-a$ 0 W

are two solutions for the universal problem so $b-a=0$ \square

hence $(*) \quad \pi \circ \tilde{f} \circ p = \pi \circ f = 0 \Rightarrow \pi \circ \tilde{f} \circ p = 0$ so \tilde{f} factors through $\ker(\operatorname{Coker} f)$

hence setting $\operatorname{Coim} f = \operatorname{Coker}(\ker f)$ and $\operatorname{Im} f = \ker(\operatorname{Coker} f)$

every morphism has a canonical factorisation

$$\ker f \rightarrow X \xrightarrow{f} Y \rightarrow \operatorname{Coker} f$$

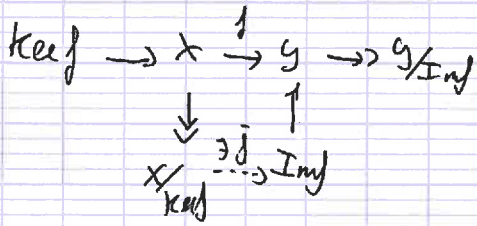
$$\downarrow a \quad \uparrow$$

$$\operatorname{Coim}(f) \rightarrow \operatorname{Im} f$$

$\exists \tilde{f}$

Example

$\mathcal{E} = \text{AbMod}$



First isomorphism theorem \bar{f} is an isomorphism.

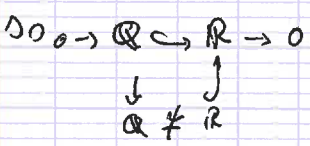
Def 5.3

Let \mathcal{E} be an additive category. Then \mathcal{E} is abelian if

- (1) Every morphism has a kernel and a cokernel in \mathcal{E}
- (2) $\forall f: X \rightarrow Y$ the canonical morphism $\bar{f}: \text{Coim} f \rightarrow \text{Im} f$ is an isomorphism.

Examples

- (1) A ring $\text{Mod} A$ A nontrivial $\text{mod} A$ of f module
- (2) \mathcal{E} abelian so is \mathcal{E}^{op}
- (3) There are examples of categories with (1) but not (2)
 eg Hausdorff topological abelian groups
 kernel = null
 coker = quotient by $\text{Im} f$



Prop 5.9

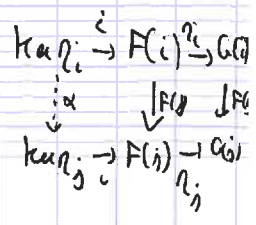
Let \mathcal{A} be an abelian category and \mathcal{J} small category. Then

- (1) $\text{Fun}(\mathcal{J}, \mathcal{A})$ is an abelian category.
- (2) $\text{Ch}(\mathcal{A})$ is an abelian category.

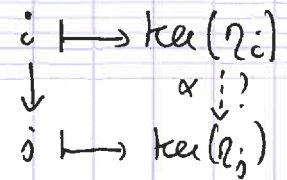
Sketch of proof

(1) $\eta: F \Rightarrow G$

$\ker(\eta): \mathcal{J} \rightarrow \mathcal{A}$

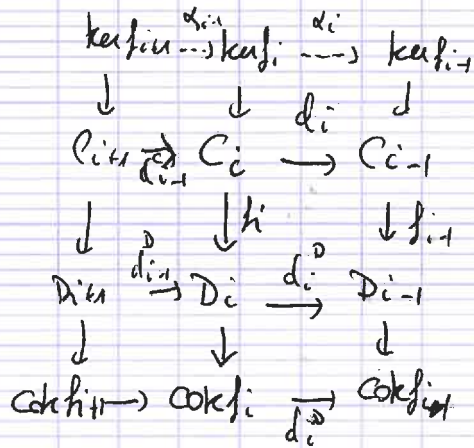


we have $\eta_j \circ F(\beta) \circ \alpha_i = \eta_j \circ \alpha_i \circ F(\beta) = 0$
 so $\exists! \alpha$



univ property of kernel \Rightarrow this is a factor + a kernel
 similarly got cokernel + $\text{Coim} f \Rightarrow \text{Im} f$ is an iso since it is
 at every evaluation.

(2) $\text{Ch}_0(\mathcal{A}) \subseteq \text{Fon}(\mathbb{Z}, \mathcal{A})$ hence $f: C_0 \rightarrow D_0$ has a
 kernel + cokernel in $\text{Fon}(\mathbb{Z}, \mathcal{A})$



$d_i \circ d_{i+1}$ is the unique
 application induced by $d_i \circ d_{i+1} = 0$

so $d_i \circ d_{i+1} = 0 \dots$

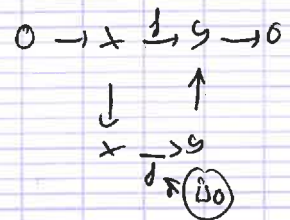
□

Rem

(0) There is another equivalent definition of abelian categories
 abt preabelian + every mono is a kernel and every epi is a cokernel

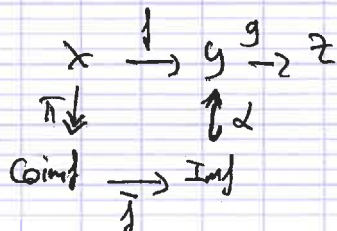
- (1) Abelian categories have finite limits and colimits
- (2) $f \in \text{Mor}(\mathcal{A})$ of abelian f mono $\Leftrightarrow \text{ker} f = 0$
 g epi $\Leftrightarrow \text{cok} g = 0$

Moreover epi + mono \Rightarrow isomorphism.



If

$x \xrightarrow{f} y \xrightarrow{g} z$ are two composable morphisms in an abelian
 category s.t. $gf = 0$



$$\begin{aligned}
 gf = 0 & \Rightarrow g \alpha f \pi = 0 \\
 & \Rightarrow g \alpha = 0
 \end{aligned}$$

So $\ker g \hookrightarrow Y \xrightarrow{g} Z$ $\bar{\alpha} : \text{Im} f \hookrightarrow \ker(g)$

$$\begin{array}{ccc} & \uparrow & \uparrow \\ & \alpha & \alpha \\ X & \xrightarrow{f} & Y \end{array}$$

Def 5.5 (1) $X \xrightarrow{f} Y \xrightarrow{g} Z$ s.t. $gf=0$ is exact if the canonical map $\text{Im} f \rightarrow \ker(g)$ is an isomorphism

(2) A complex (C_n, d_n) is exact if $\text{Im}(d_{n+1}) \xrightarrow{\sim} \ker(d_n) \forall i \in \mathbb{Z}$

(3) A short exact sequence is an exact complex of the form

$$0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0$$

Example in $\text{Mod} A$

$$X \xrightarrow{f} Y \xrightarrow{g} Z \quad gf=0 \Rightarrow \text{Im} f \subseteq \ker(g)$$

so exactness $\Leftrightarrow \text{Im} f = \ker(g)$

(1) $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z$ exact iff $f = \ker g$

(2) $\text{Seq} \Leftrightarrow \begin{cases} f \text{ mono} \\ g \text{ epi} \\ \text{Im} f \xrightarrow{\sim} \ker g \end{cases} \Leftrightarrow \begin{cases} g = \text{coker}(f) \\ f = \ker(g) \end{cases}$

Rem There is a difficult theorem of Freyd-Mitchell saying that any abelian category can be seen as a full subcategory of $\text{Mod} A$ for some ring A in such a way that the abelian structure is induced by the usual one in $\text{Mod} A$.

\leadsto we will most of the time use this result to simplify the exposition.

~~1/10/19~~

Def 5.6

Let \mathcal{A} be an abelian category and $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$

$$0 \rightarrow D \xrightarrow{h} E \xrightarrow{k} F \rightarrow 0$$

two seq. A morphism of seq is the data of 3 morphism α, β, γ

making the diagram a commutative diag

$$\begin{array}{ccccccc}
 0 & \rightarrow & A & \xrightarrow{m} & B & \xrightarrow{e} & C \rightarrow 0 \\
 & & \alpha \downarrow & & \downarrow \beta & & \downarrow \gamma \\
 0 & \rightarrow & D & \xrightarrow{m'} & E & \xrightarrow{e'} & F \rightarrow 0
 \end{array} \quad (*)$$

Lemma 5.7 [Short five lemma] In (*) $\begin{cases} f+h \text{ mono} \Rightarrow g \text{ mono} \\ f+h \text{ epi} \Rightarrow g \text{ epi} \\ f+h \text{ iso} \Rightarrow g \text{ iso} \end{cases}$

Proof (1) In Mod A $x \in B$; $\beta(x) = 0 \Rightarrow e' \beta(x) = 0 \Rightarrow \gamma \circ e(x) = 0$
 $\gamma \text{ mono} \Rightarrow e(x) = 0 \Rightarrow x \in \ker(e)$; $x \in \ker(e) \Rightarrow \text{Im}(m)$
 $\Rightarrow \exists a; x = m(a)$

$$\begin{aligned}
 0 = \beta m(a) &= m' \alpha(a) \Rightarrow \alpha(a) \in \ker m' = \{0\} \\
 &\Rightarrow \alpha(a) = 0 \Rightarrow a = 0 \Rightarrow x = 0.
 \end{aligned}$$

+ dual for the second statement.

(2) In an arbitrary abelian category

$$\begin{array}{ccccccc}
 & & \exists k' \text{ --- } \ker & & & & \\
 & & \swarrow & \downarrow k & & & \\
 0 & \rightarrow & A & \xrightarrow{m} & B & \xrightarrow{e} & C \rightarrow 0 \\
 & & \alpha \downarrow & & \downarrow \beta & & \downarrow \gamma \\
 0 & \rightarrow & D & \xrightarrow{m'} & E & \xrightarrow{e'} & F \rightarrow 0
 \end{array}$$

$\beta k = 0$
 give $\gamma e k = e' \beta k = 0$
 $\gamma \text{ mono} \Rightarrow e k = 0$ so k factorizes
 through $\ker(e) = A$
 \hookrightarrow gives k'

we have $\beta k = 0 \Rightarrow 0 = \beta m k' = \underbrace{m'}_{\text{mono}} \alpha k' \Rightarrow k' = 0$ so $k = 0$

so β mono. The rest is dual. \square

Thm 5.8 [Splitting lemma] Let $0 \rightarrow A \xrightarrow{q} B \xrightarrow{R} C \rightarrow 0$ be a short exact sequence in an abelian category \mathcal{A} . $\forall A \in \mathcal{A}$

- (1) $\exists t: B \rightarrow A$ s.t. $tq = Id_A$.
- (2) $\exists s: C \rightarrow B$ s.t. $Rs = Id_C$.
- (3) $\exists h: B \xrightarrow{\cong} A \oplus C$ s.t.

$$\begin{array}{ccccccc}
 0 & \rightarrow & A & \xrightarrow{q} & B & \xrightarrow{R} & C \rightarrow 0 \\
 & & \parallel & & \downarrow h & & \parallel \\
 0 & \rightarrow & A & \xrightarrow{i_A} & A \oplus C & \xrightarrow{\pi_C} & C \rightarrow 0
 \end{array}$$

is an isomorphism of ses.
 In this case we say that the ses splits.

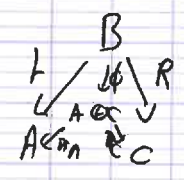
Proof

3 \Rightarrow 2 $s: C \xrightarrow{i_C} A \oplus C \xrightarrow{h^{-1}} B$ then $C \xrightarrow{i_C} A \oplus C \xrightarrow{h^{-1}} B \xrightarrow{R} C$
 $\pi_C \circ h^{-1} \circ i_C = Id_C$

3 \Rightarrow 1 $t: B \xrightarrow{h} A \oplus C \xrightarrow{\pi_A} A$

1 \Rightarrow 3

$$\begin{array}{ccccccc}
 0 & \rightarrow & A & \xleftarrow{h} & B & \xrightarrow{R} & C \rightarrow 0 \\
 & & \parallel & & \downarrow \phi & & \downarrow \\
 0 & \rightarrow & A & \xrightarrow{i_A} & A \oplus C & \xrightarrow{\pi_C} & C \rightarrow 0
 \end{array}$$



by construction $\pi_C \circ \phi = R$ and $\pi_A \circ \phi = h$

$$\begin{aligned}
 \phi \circ q &= Id_{A \oplus C} \circ \phi \circ q = (i_A \circ \pi_A + i_C \circ \pi_C) \phi \circ q \\
 &= i_A \pi_A \phi \circ q + i_C \pi_C \phi \circ q \\
 &= i_A \underbrace{\pi_A \phi \circ q}_h + i_C \underbrace{\pi_C \phi \circ q}_R \\
 &= i_A
 \end{aligned}$$

by the short five lemma ϕ is an iso □

Def 5.9 Let \mathcal{E} and \mathcal{D} two abelian categories. Let $F: \mathcal{E} \rightarrow \mathcal{D}$ a functor.

- (1) F is left exact if F preserves finite limits
- (2) F is right exact if F preserves finite colimits
- (3) F is exact if it preserves finite limits and colimits

Lemma 5.10 Let $F: \mathcal{E} \rightarrow \mathcal{D}$ be an additive functor between abelian categories. $\forall F \in \mathcal{E}$

- (1) F is left exact
- (2) F preserves kernels i.e. $F(\ker f) \xrightarrow{\cong} \ker(Ff)$
- (3) $\forall 0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z$ exact sequence in \mathcal{E} , the sequence $0 \rightarrow F(X) \xrightarrow{F(f)} F(Y) \xrightarrow{F(g)} F(Z)$ is exact
- (4) $\forall 0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0$ seq in \mathcal{E} the sequence $0 \rightarrow F(X) \xrightarrow{F(f)} F(Y) \xrightarrow{F(g)} F(Z)$ is exact.

Proof

(1) \Rightarrow (2) clear kernels are limits

(2) \Rightarrow (3) ~~seq~~ sequence $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z$ exact // $f = \ker(g)$

so $F(f)$ is a kernel of $F(g) \Rightarrow$ second sequence is exact

(3) \Rightarrow (4) $f = \ker(g) \Rightarrow$ same

(3) \Rightarrow (2) ~~$0 \rightarrow \ker f \xrightarrow{i} X \xrightarrow{f} Y \xrightarrow{g} Z \rightarrow 0$~~ clear

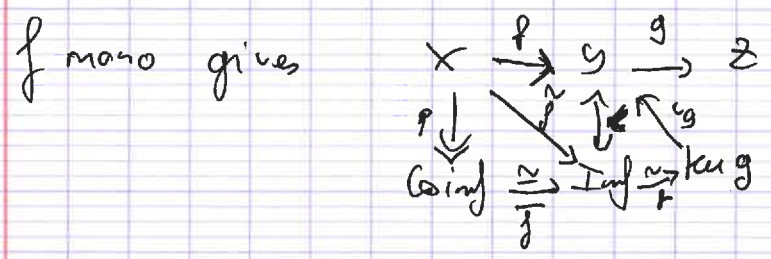
(2) \Rightarrow (1) also since any finite limits can be expressed using kernels and products

we are missing (4) \Rightarrow (3) $0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z$ then consider

$$0 \rightarrow X \xrightarrow{f} Y \rightarrow \text{coker}(f) \rightarrow 0$$

$0 \rightarrow X \xrightarrow{f} Y \xrightarrow{g} Z$ exact iff $f = \ker g$
 \Leftarrow clear

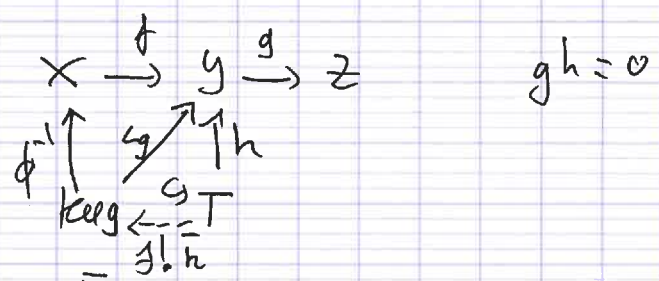
\Rightarrow exactness $\Rightarrow \text{Im } f \cong \ker g + f$ mono



$\tilde{g} \circ \tilde{f} \circ \rho = f$ f mono $\Rightarrow \rho$ mono so $\tilde{f} = \tilde{f} \circ \rho$ iso

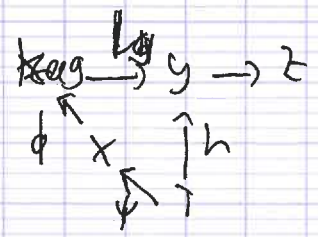
so $\phi = \ker \tilde{f}$ is an iso s.t. $c_g \circ \phi = f \Rightarrow f \circ \phi^{-1} = c_g$

Let us prove that f is $\ker g$ by checking univ property



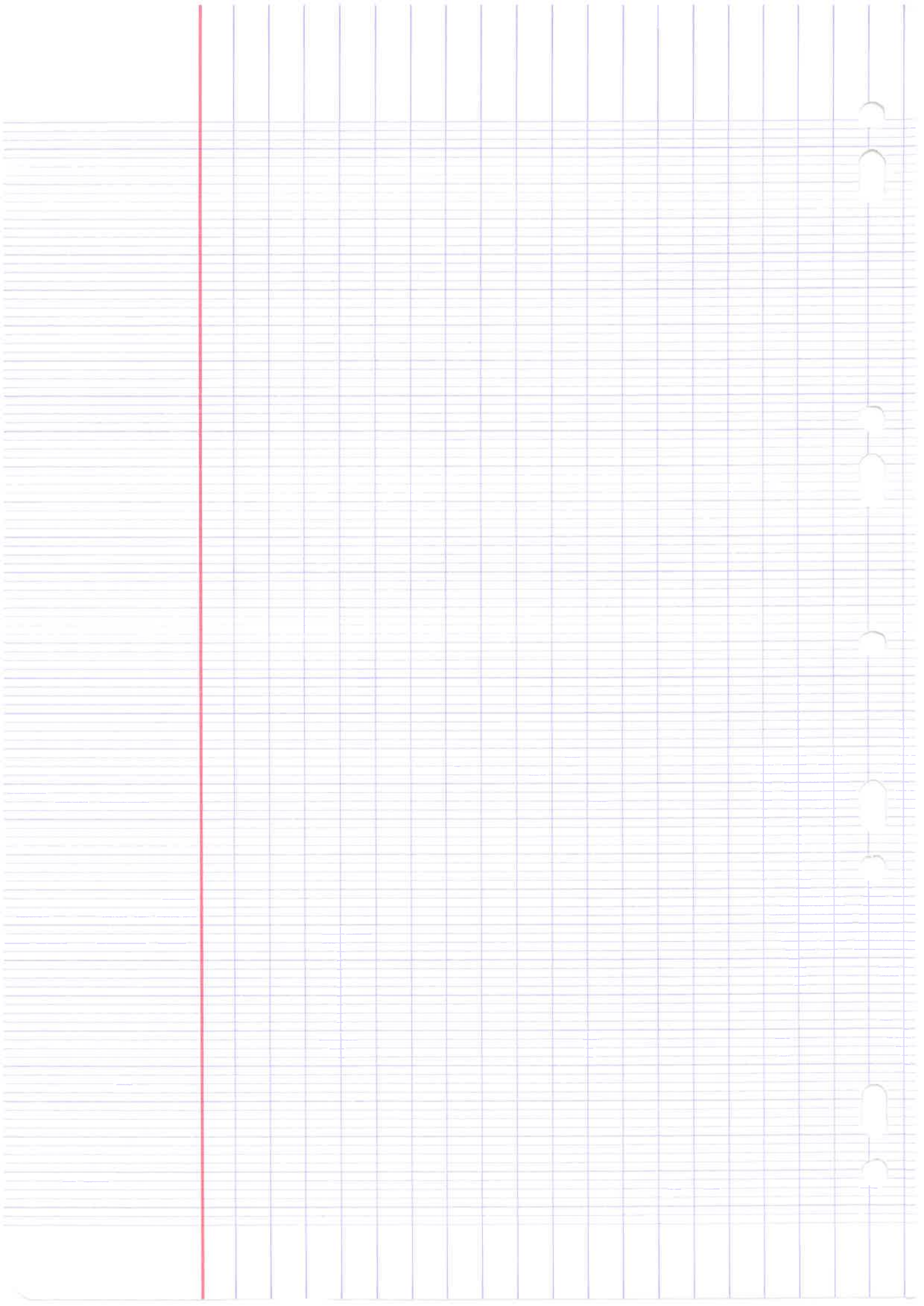
$f \circ \phi^{-1} \circ \tilde{h} = c_g \tilde{h} = h$

if $\begin{array}{ccc} X & \xrightarrow{f} & Y & \xrightarrow{g} & Z \\ & \nearrow \psi & \uparrow h & & \\ & & T & & \end{array}$ is another factorisation then



$c_g \circ \phi \circ \psi = f \circ \psi = h$

so $\phi \circ \psi = \tilde{h}$ so $\psi = \phi^{-1} \circ \tilde{h}$



applying F gives $F(f)$ mono. So F preserves mono.

Moreover we have $0 \rightarrow x \xrightarrow{f} y \xrightarrow{g} \text{Im}(g) \rightarrow 0$ exact so

$$0 \rightarrow F(x) \xrightarrow{F(f)} F(y) \xrightarrow{F(g)} F(\text{Im}(g)) \rightarrow 0 \text{ exact}$$

since $i: \text{Im}(g) \hookrightarrow Z$ is mono + F preserves mono we have $F(i): \text{Im}(g) \rightarrow F(Z)$ mono so does not change the kernel!

$$\hookrightarrow 0 \rightarrow F(x) \xrightarrow{F(f)} F(y) \xrightarrow{F(g)} F(Z) \rightarrow 0 \text{ exact} \quad \square$$

Coro 5.11 TFAE for an additive functor between abelian categories

(1) F is exact

(2) \forall seq $0 \rightarrow A \xrightarrow{\alpha} B \xrightarrow{\beta} C \rightarrow 0$ the sequence $0 \rightarrow F(A) \xrightarrow{F(\alpha)} F(B) \xrightarrow{F(\beta)} F(C) \rightarrow 0$ is exact.

Prop 5.12 (1) $\text{Hom}_E(-, -): E^{\text{op}} \times E \rightarrow \text{Mod } \mathbb{Z}$ is left exact in each variable

(2) $- \otimes_A -: \text{Mod } A \times A \text{Mod} \rightarrow \mathbb{Z} \text{Mod}$ is right exact in each variable

(3) $F + G$, then F is right exact and G is left exact

Proof (3) clear since ~~RAP~~ ~~LAP~~ adjoint preserves colim

2 - Chain complexes in abelian categories

Def 5.13 \mathcal{A} be an abelian category $(X_\bullet, d_\bullet) \in \text{Ch}_\bullet(\mathcal{A})$

For $n \in \mathbb{Z}$ we set $Z_n(X) = \ker(d_n)$ "n-zycles"
 $B_n(X) = \text{Im}(d_{n+1})$ "n-boundaries"
 $H_n(X) = Z_n(X) / B_n(X)$ "nth homology of X"
 \triangleq in an arbitrary abelian category = $\text{coker of } B_n \hookrightarrow Z_n$

when Cochain : speaks about cocycles, coboundaries, cohomology.

$f: X_\bullet \rightarrow Y_\bullet$ chain morphism:

$$\begin{array}{ccc}
 \ker d_n^X \xrightarrow{i} X_n \xrightarrow{d_n^X} X_{n-1} & & 0 = f_{n-1} d_n^X i = d_n^Y \circ f_n \circ i \\
 \downarrow \beta_n & \downarrow \beta_n & \downarrow \beta_{n-1} \\
 \ker d_n^Y \rightarrow Y_n \xrightarrow{d_n^Y} Y_{n-1} & & \text{so } \exists ! z(f_n)
 \end{array}$$

(*) similarly $X_{n-1} \xrightarrow{f_{n-1}} Y_{n-1}$ induces

$$\begin{array}{ccc}
 \text{Im}(d_n^X) \xrightarrow{i} X_{n-1} \xrightarrow{\pi} \text{Cok}(d_n^X) & & \\
 \downarrow \beta_{n-1} & \downarrow \beta_{n-1} & \downarrow \beta_{n-1} \\
 \text{Im}(d_n^Y) \xrightarrow{i} Y_{n-1} \xrightarrow{\pi} \text{Cok}(d_n^Y) & &
 \end{array}$$

$$\text{so } \pi f_{n-1} i = f_{n-1} \pi i = 0$$

So we have

$$\begin{array}{ccc}
 \text{Im } d_{n+1}^X \xrightarrow{i} \ker d_n \rightarrow \text{Cok}(i) & & \\
 \downarrow \beta_n & \downarrow \beta_n & \downarrow \beta_n \\
 \text{Im } d_{n+1}^Y \xrightarrow{i} \ker d_n \xrightarrow{\pi} \text{Cok}(i) & &
 \end{array}$$

In Mod A $H_n(X) = \ker d_n / \text{Im } d_{n+1}$ $H_n(f)([x]) = [f_n(x)]$

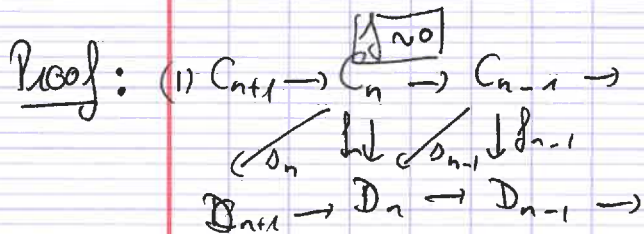
We get $H_n: \text{Ch}_\bullet(A) \rightarrow A$ a functor called the n th homology functor. Moreover it is an additive functor.

Def 5.14 $f: X_\bullet \rightarrow Y_\bullet$ is a quasi-isomorphism if $H_n(f)$ is an isomorphism $\forall n \in \mathbb{Z}$.

Prop 5.15 $f: X_* \rightarrow Y_* \in \text{class}(C_*(A))$ for A an abelian cogroup.

(1) If $f \sim g$ then $H_n(f) = H_n(g)$

(2) If f is a homotopy equivalence, then it is a quasi-isomorphism



$$f_n = s_{n-1} d_n^C + d_{n+1}^D s_n$$

$$H_n(f) : H_n(C) = \frac{\ker(d_n)}{\text{Im}(d_{n+1})} \rightarrow \frac{\ker(d_n^D)}{\text{Im}(d_{n+1}^D)} = H_n(Y) = \frac{\ker(d_{n+1}^D)}{\text{Im}(d_{n+1}^D)} = 0$$

due $f \sim g \Rightarrow f \sim g \sim 0 \Rightarrow H_n(f) = H_n(g) = 0$

(2) $f \sim g \sim \text{Id} \Rightarrow H_n(f)H_n(g) = H_n(\text{Id}) = \text{Id} \quad \square$

Def 5.16 (1) C_* is contractible if C is homotopy equivalent to 0

($\text{Id} \sim 0$)

(2) C_* is acyclic if C is qis to 0

\hookrightarrow Contractible \Rightarrow acyclic

Thm 5.16 [long exact sequence] A short exact sequence

$$0 \rightarrow C' \xrightarrow{\alpha} C \xrightarrow{\beta} C'' \rightarrow 0 \text{ gives rise to a long exact sequence}$$

$$\rightarrow H_n(C') \xrightarrow{H_n(\alpha)} H_n(C) \xrightarrow{H_n(\beta)} H_n(C'') \xrightarrow{\delta_n} H_{n-1}(C) \rightarrow \dots$$

where the δ_n (to be defined) are called the "connecting" homomorphisms.

Proof (1) In Mod A \mathcal{S} is commutative via diagram chasing

$$\begin{array}{ccccccc}
 0 & \rightarrow & C_n' & \xrightarrow{\alpha_n} & C_n & \xrightarrow{\beta_n} & C_n'' \rightarrow 0 \\
 & & \downarrow d_n' & & \downarrow d_n & & \downarrow d_n'' \\
 0 & \rightarrow & C_{n-1}' & \xrightarrow{\alpha_{n-1}} & C_{n-1} & \xrightarrow{\beta_{n-1}} & C_{n-1}'' \rightarrow 0 \\
 & & \downarrow d_{n-1}' & & \downarrow d_{n-1} & & \downarrow d_{n-1}''
 \end{array}$$

$$H_n(C'') = \ker(d_n'') / \text{Im}(d_n'') \quad H_{n-1}(C) = \ker(d_{n-1}') / \text{Im}(d_{n-1}')$$

$$x \in \ker(d_n') \exists y \in C_n; \beta_n(y) = x \text{ and}$$

$$0 = d_n'' \beta_n(y) = \beta_{n-1} \circ d_n(y) \Rightarrow d_n(y) \in \ker(\beta_{n-1}) \Rightarrow \exists z \in C_{n-1}' \text{ by}$$

$$d_n(y) = \alpha_{n-1}'(z)$$

Set $\mathcal{S}_n([x]) := [z]$

well-defined? - if y' is another lift of x

$$\begin{aligned}
 & \alpha_n(c) \\
 & y - y' \in \ker \beta_n = \text{Im } d_n \\
 & d_n \alpha_n(c) = \alpha_{n-1}' d_n'(c) \\
 & \text{so } z \text{ differs by an element} \\
 & \text{of } \text{Im } d_n' \\
 & \text{can be chosen}
 \end{aligned}$$

- if $x' = x + t \in \text{Im}(d_{n+1}')$ then lift of $t \in \text{Im}(d_{n+1}')$ so $d_n(y) = 0$

$$\begin{aligned}
 - z \in \ker(d_{n-1}')? \quad \alpha_{n-2}' \circ d_{n-1}'(z) &= d_{n-1} \circ \alpha_{n-1}'(z) \\
 &= d_{n-1} d_n'(z) = 0
 \end{aligned}$$

(2) Using snake lemma exercise

It remains to see that the long sequence is exact

$$(*) \quad H_n(C') \xrightarrow{H_n(\alpha)} H_n(C) \xrightarrow{H_n(\beta)} H_n(C'')$$

$$[x] \in \ker(H_n(\beta)) \quad \ker \beta_n(x) \in \text{Im } d_{n+1}''$$

$$\begin{array}{ccccccc} 0 & \rightarrow & C'_{n+1} & \xrightarrow{d_{n+1}} & C_{n+1} & \xrightarrow{\beta_{n+1}} & C''_{n+1} \rightarrow 0 \\ & & \downarrow & & \downarrow d_{n+1} & & \downarrow d''_{n+1} \\ 0 & \rightarrow & C'_n & \xrightarrow{\alpha_n} & C_n & \xrightarrow{\beta_n} & C''_n \rightarrow 0 \end{array}$$

$$\Rightarrow \exists c \in C_{n+1}; \beta_n(x) = d_{n+1}(c) \quad \text{so } \exists c \in C_{n+1}; \beta_n(x) = d_{n+1} \beta_{n+1}(c) = d_{n+1}(c)$$

$$\Rightarrow x - d_{n+1}(c) \in \ker(\beta_n) = \text{Im } \alpha_n$$

$$\Rightarrow x - d_{n+1}(c) = \alpha_n(c')$$

$$\text{we have } [x] = [x - d_{n+1}(c)] = [\alpha_n(c')] \quad \text{and}$$

$$d_n(x) = 0 \Rightarrow c \in \ker d_n'$$

$$\text{so } \ker(H_n(\beta)) \subseteq \text{Im}(H_n(\alpha))$$

Conversely we have $\text{Im}(\alpha) \subseteq \ker(\beta)$ so also $\text{Im}(H_n(\alpha)) \subseteq \ker H_n(\beta)$.

• $\boxed{\text{Im}(H_n(\beta)) = \ker(\delta_n)}$ if $c'' \in \text{Im}(\beta_n)$ then c is a lift of $\beta_n(c'')$
 $c'' = \beta_n(c)$

$$\text{so } d_n(c) = 0 \text{ because } c \in \ker(d_n)$$

$$\text{so } \text{Im}(H_n(\beta_n)) \subseteq \ker(\delta_n)$$

Conversely if $x \in \ker(\delta_n)$ we have

$$\begin{array}{l} y \text{ s.t. } x \\ \downarrow \\ \delta_n(y) = x \end{array} \quad \text{so } d_n(y) = 0 \quad \text{so } y \in \ker d_n$$

~~we can so $x = \alpha_n(y)$ with $y \in \ker d_n$~~

$$[\delta(x)] = 0 \Rightarrow \delta(x) \in \text{Im}(d_n) \quad \text{so } \exists c' \in C'_n; \delta(x) = d'_n(c')$$

$$\text{so } \alpha_{n-1} d'_n(c') = \alpha_{n-1}(\delta(x)) = d_n(y)$$

$$d_n \alpha_n(c') \quad \text{so } y - \alpha_n(c') \in \ker d_n \quad \text{and } \beta_n(y - \alpha_n(c')) = x!$$

$$\textcircled{*} \boxed{\text{Ker}(H_{n-1}(\alpha)) = \text{Im}(\delta_n)}$$

$$\textcircled{*} [x] \in \text{Ker}(H_{n-1}(\alpha)) \Rightarrow \alpha_{n-1}(x) \in \text{Im}(d_{n-1}) \\ \Rightarrow \alpha_{n-1}(x) = d_n(y) \text{ and } \beta_n(y) \in \mathbb{C}_n^2$$

y is a lift so $\delta_n([x]) = x$

$\textcircled{*} \alpha_n(\delta_n(x)) \in \text{Im} d_n$ by construction □

Remark (1)
$$\boxed{\delta_n = \alpha_{n-1}^{-1} \circ d_n \circ \beta_n^{-1}}$$

(2) δ is natural in the following sense

if

$$\begin{array}{ccccccc} 0 & \rightarrow & A & \xrightarrow{\alpha} & B & \xrightarrow{\beta} & C \rightarrow 0 \\ & & \downarrow f & & \downarrow g & & \downarrow h \\ 0 & \rightarrow & A' & \xrightarrow{\alpha'} & B' & \xrightarrow{\beta'} & C' \rightarrow 0 \end{array}$$

are two ses then

We have a morphism of long exact sequence

$$\begin{array}{ccccccc} \dots & \rightarrow & H_n(A) & \xrightarrow{H_n(\alpha)} & H_n(B) & \rightarrow & H_n(C) \xrightarrow{\delta_n} H_{n-1}(A) \rightarrow H_{n-1}(B) \rightarrow \dots \\ & & \downarrow H_n(f) & \hookrightarrow & \downarrow H_n(g) & \hookrightarrow & \downarrow H_n(h) \\ & & H_n(A') & \xrightarrow{H_n(\alpha')} & H_n(B') & \xrightarrow{H_n(\beta')} & H_n(C') \xrightarrow{\delta_n'} H_{n-1}(A') \rightarrow H_{n-1}(B') \rightarrow \dots \\ & & & & & & \boxed{\text{Need a proof}} \end{array}$$

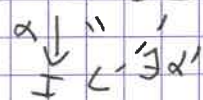
just because H_n is a functor

3. Projective and injective objects

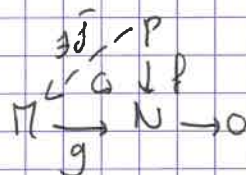
Def 5.17 Let \mathcal{E} be an abelian category. An object

- (1) $I \in \mathcal{O}(\mathcal{E})$ is injective if $\text{Hom}_{\mathcal{E}}(\cdot, I) : \mathcal{E}^{\text{op}} \rightarrow \text{Ab}$ is exact
- (2) $P \in \mathcal{O}(\mathcal{E})$ is projective if $\text{Hom}_{\mathcal{E}}(P, \cdot) : \mathcal{E} \rightarrow \text{Ab}$ is exact
- (3) \mathcal{E} has enough projectives (injectives) if $\forall X \in \mathcal{O}(\mathcal{E}) \exists f : P \rightarrow X$ an epi with P projective ($g : X \rightarrow I$ a mono with I injective),

Prop 5.18 (1) $I \in \text{ob}(\mathcal{E})$ is injective iff $\forall 0 \rightarrow X \xrightarrow{f} Y$



(2) $P \in \mathcal{E}$ is projective iff \forall



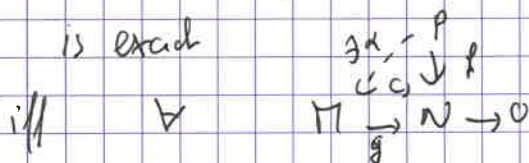
Proof P is projective iff $\text{Hom}_{\mathcal{E}}(P, -)$ is exact

iff $\text{Hom}_{\mathcal{E}}(P, -)$ is right exact

iff $\forall M \xrightarrow{g} N \rightarrow 0$ exact

The induced sequence $\text{Hom}(P, M) \rightarrow \text{Hom}(P, N) \rightarrow 0$

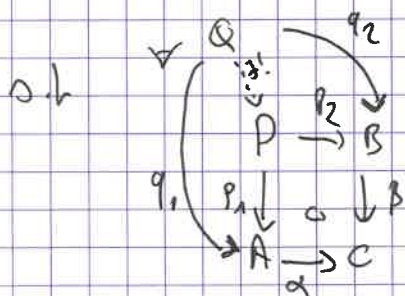
is exact



(1) is dual □

Def 5.19 Let \mathcal{E} be an abelian category. A pullback of $A \xrightarrow{\alpha} C$ $B \xrightarrow{\beta} C$ is an object P together with two morphisms $p_1: P \rightarrow A$ $p_2: P \rightarrow B$

$p_2: P \rightarrow B$



Dually we have notion of pushout.

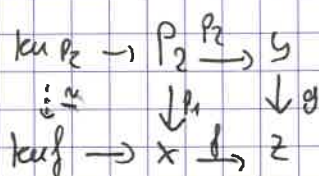
Prop 5.20 Let \mathcal{A} be an abelian category

(1) Any pair of morphism has a pullback in \mathcal{A} and a pushout in \mathcal{A}

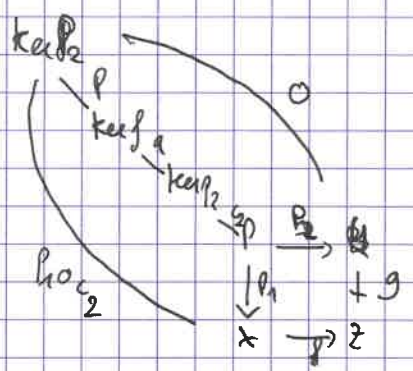
(2) If β is a mono then p_1 is a mono: pullback preserves mono

(3) Pushout preserves epi

(4) Pullback preserves kernels:



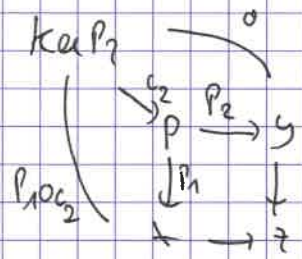
Now consider the Pullback diagram:



$$p_1 \circ c_2 \circ q \circ p = c_1 \circ p = p_1 \circ c_2$$

(1)

Hence

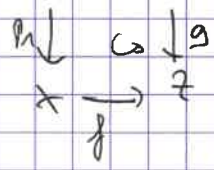


are the same pullback so we have

$$c_2 = c_2 \circ q \circ p \text{ so } q \circ p = \text{Id since}$$

c_2 is a monomorphism.

(5) Let $P \xrightarrow{p_2} Y$ be a pull back. Then we have a



exact sequence

$$P \xrightarrow{\begin{pmatrix} p_1 \\ p_2 \end{pmatrix}} X \oplus Y \xrightarrow{(-f, g)} Z$$

Moreover saying that P is a pull back is equivalent to $\begin{pmatrix} p_1 \\ p_2 \end{pmatrix} = \ker(-f, g)$

indeed

$$P \xrightarrow{\begin{pmatrix} p_1 \\ p_2 \end{pmatrix}} X \oplus Y \xrightarrow{(-f, g)} Z$$

$\uparrow \begin{pmatrix} a \\ b \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix}$
 $\exists! \alpha \quad \exists \alpha \quad \exists \alpha$

$$(-f, g) \begin{pmatrix} a \\ b \end{pmatrix} = 0 \Leftrightarrow af = gb$$

$$\exists \alpha \text{ s.t. } \begin{pmatrix} p_1 \\ p_2 \end{pmatrix} \alpha = \begin{pmatrix} a \\ b \end{pmatrix}$$

is equivalent to $\begin{cases} \alpha p_1 = a \\ \alpha p_2 = b \end{cases}$

If f is an epimorphism then $(-f, g)$ is an epimorphism so we have

a short exact sequence $0 \rightarrow P \xrightarrow{\begin{pmatrix} p_1 \\ p_2 \end{pmatrix}} X \oplus Y \xrightarrow{(-f, g)} Z \rightarrow 0$ saying $\text{Coker} \begin{pmatrix} p_1 \\ p_2 \end{pmatrix} = (-f, g)$.

