

Rigidification of quasi-categories

Daniel Dugger and David Spivak

Sylvain Rossi

1. Introduction
2. Necklaces
3. Categorification of necklaces
4. The categorification functor
5. Flagged necklaces and totally nondegenerate maps
6. An example
7. Homotopy theory
8. Meta-theory via gadgets
9. Properties of categorification

Introduction

Adjunction

We have the following (Quillen) adjunction between the categorification functor and the coherent nerve functor:

$$\mathfrak{C}: \hat{\Delta}_J \rightleftarrows \text{Cat}_\Delta: N^\Delta$$

Introduction

For a simplex Δ^n , Lurie defines $\mathfrak{C}(\Delta^n) := \tilde{P}_n$, where \tilde{P}_n is the simplicial category whose objects are $\{0, 1, \dots, n\}$ and morphisms $\tilde{P}_n(i, j) = N(P_{i,j})$, where $P_{i,j}$ denotes the poset category of subsets of $\{i, i+1, \dots, j\}$ containing i and j (ordered by inclusion).

For $\mathcal{D} \in \text{Cat}_\Delta$, the **coherent nerve** $N^\Delta(\mathcal{D})$ of \mathcal{D} is then defined as the simplicial set given by

$$[n] \mapsto \text{Cat}_\Delta(\mathfrak{C}(\Delta^n), \mathcal{D}).$$

Introduction

For a simplex Δ^n , Lurie defines $\mathfrak{C}(\Delta^n) := \tilde{P}_n$, where \tilde{P}_n is the simplicial category whose objects are $\{0, 1, \dots, n\}$ and morphisms $\tilde{P}_n(i, j) = N(P_{i,j})$, where $P_{i,j}$ denotes the poset category of subsets of $\{i, i+1, \dots, j\}$ containing i and j (ordered by inclusion).

For $\mathcal{D} \in \text{Cat}_\Delta$, the **coherent nerve** $N^\Delta(\mathcal{D})$ of \mathcal{D} is then defined as the simplicial set given by

$$[n] \mapsto \text{Cat}_\Delta(\mathfrak{C}(\Delta^n), \mathcal{D}).$$

Since any simplicial set K may be written as $\text{colim}_{\Delta^n \rightarrow K} \Delta^n$ (in $\hat{\Delta}$), we see that a left adjoint of N^Δ is given by

$$\mathfrak{C}(K) \cong \text{colim}_{\Delta^n \rightarrow K} \mathfrak{C}(\Delta^n).$$

(The colimit is computed in Cat_Δ)

Introduction

The main issue is that finding the mapping spaces of $\mathcal{C}(K)$ via the definition above is a mess... We therefore need to introduce new notions so these colimits become computable.

Necklaces

Necklaces

A **necklace** T is a simplicial set of the form

$$T = \Delta^{n_0} \vee \Delta^{n_1} \vee \dots \vee \Delta^{n_k},$$

where each $n_i \in \mathbb{Z}_{\geq 0}$ and where each final vertex of Δ^{n_i} has been glued to the initial vertex of $\Delta^{n_{i+1}}$.

Each Δ^{n_i} is called a **bead** of T .

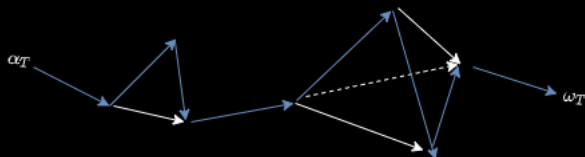
A **joint** is either an initial or final vertex of some bead, the union of which is denoted by J_T . Note that this set is totally ordered. Thus any necklace T comes equipped with a canonical map

$$\partial\Delta^1 \rightarrow T$$

sending 0 to the initial vertex α_T of T and 1 to the final vertex ω_T of T .

Necklaces

e.g. $T = \Delta^1 \vee \Delta^2 \vee \Delta^1 \vee \Delta^3 \vee \Delta^1$ (in blue we have $\text{Spi}[T]$, the spine of T)



Necklaces

A **spine** is a necklace in which every bead is a Δ^1 .

Every necklace T has an associated spine denoted $\text{Spi}[T] \hookrightarrow T$.

A **simplex** is a necklace with one bead.

Every necklace T has an associated simplex $\Delta[T]$ ($T \hookrightarrow \Delta[T]$), namely the simplex with vertex set T_0 .

The assignment $T \rightarrow \Delta[T]$ is functorial, $T \rightarrow \text{Spi}[T]$ is not.

Necklaces

We wish to show that the necklaces form a full subcategory Nec of the category $\hat{\Delta}_{*,*} := (\partial\Delta^1 \downarrow \hat{\Delta})$. To do this, we work in greater generality and make use of **ordered simplicial sets**:

Let $X \in \hat{\Delta}$, for $x, y \in X_0$ we say $x \preceq y$ if there exists a spine T and a map $T \rightarrow X$ such that $\alpha_T \mapsto x$ and $\omega_T \mapsto y$. This relation is clearly reflexive and transitive, but not necessarily antisymmetric.

Necklaces

We wish to show that the necklaces form a full subcategory Nec of the category $\hat{\Delta}_{*,*} := (\partial\Delta^1 \downarrow \hat{\Delta})$. To do this, we work in greater generality and make use of **ordered simplicial sets**:

Let $X \in \hat{\Delta}$, for $x, y \in X_0$ we say $x \preceq y$ if there exists a spine T and a map $T \rightarrow X$ such that $\alpha_T \mapsto x$ and $\omega_T \mapsto y$. This relation is clearly reflexive and transitive, but not necessarily antisymmetric.

Ordered Simplicial Set

A simplicial set X is **ordered** if

- ▶ the relation \preceq is antisymmetric;
- ▶ an n -simplex in X is determined by its sequence of vertices $x(0) \preceq \cdots \preceq x(n)$, i.e. no two distinct n -simplices have the same vertex sequences.

Necklaces

Simple Inclusion

Let $A, X \in \hat{\Delta}$. A map $A \rightarrow X$ is called a **simple inclusion** if it has the RLP with respect to the canonical inclusion $\partial\Delta^1 \hookrightarrow T$ for all necklaces T .

The idea is that if there is a path in X that starts and ends in A , it must entirely lie in A .

Lemma 1

A simple inclusion $A \hookrightarrow X$ has the RLP with respect to the maps $\partial\Delta^k \hookrightarrow \Delta^k$ for all $k \geq 1$.

Necklaces

Lemma 2

Let X, Y be ordered simplicial sets and $f: X \rightarrow Y$ a map.

2. Every necklace is an ordered simplicial set.
5. If T is a necklace and $y: T \rightarrow X$ is map, then the image is a necklace.
6. Suppose that $X \leftarrow A \rightarrow Y$ is a diagram of ordered simplicial sets and both $A \rightarrow X$ and $A \rightarrow Y$ are simple inclusions. Then the pushout $B = X \sqcup_A Y$ is an ordered simplicial set and the inclusions $X \hookrightarrow B$ and $Y \hookrightarrow B$ are simple.

Necklaces

Lemma 2

Let X, Y be ordered simplicial sets and $f: X \rightarrow Y$ a map.

2. Every necklace is an ordered simplicial set.
5. If T is a necklace and $y: T \rightarrow X$ is map, then the image is a necklace.
6. Suppose that $X \leftarrow A \rightarrow Y$ is a diagram of ordered simplicial sets and both $A \rightarrow X$ and $A \rightarrow Y$ are simple inclusions. Then the pushout $B = X \sqcup_A Y$ is an ordered simplicial set and the inclusions $X \hookrightarrow B$ and $Y \hookrightarrow B$ are simple.

Thanks to Lemma 2 2.+ 5., we see that $Nec \subset \hat{\Delta}_{*,*}$ is a full subcategory.

Let $S, T \in Nec$, then the pushout along the canonical maps $S \leftarrow \partial\Delta^1 \rightarrow T$, is denoted by $S \vee T$ and is clearly a necklace again.

Categorification of necklaces

Let T be a necklace, we will now describe $\mathcal{C}(T)$. The object set is T_0 .

For $a, b \in T_0$, let $V_T(a, b)$ denote the set of vertices in T between a and b (included). $J_T(a, b)$ denotes union of $\{a, b\}$ and the set of joints between a and b . There is a unique subnecklace of T (by Lemma 2.2.) with joints $J_T(a, b)$ and vertex set $V_T(a, b)$. Let $\tilde{B}_1, \dots, \tilde{B}_k$ denote its beads. There are canonical inclusions of the \tilde{B}_i 's into T .

Categorification of necklaces

Let T be a necklace, we will now describe $\mathfrak{C}(T)$. The object set is T_0 .

For $a, b \in T_0$, let $V_T(a, b)$ denote the set of vertices in T between a and b (included). $J_T(a, b)$ denotes union of $\{a, b\}$ and the set of joints between a and b . There is a unique subnecklace of T (by Lemma 2.2.) with joints $J_T(a, b)$ and vertex set $V_T(a, b)$. Let $\tilde{B}_1, \dots, \tilde{B}_k$ denote its beads. There are canonical inclusions of the \tilde{B}_i 's into T . Hence, there is a canonical map

$$\mathfrak{C}(\tilde{B}_k)(j_k, b) \times \mathfrak{C}(\tilde{B}_{k-1})(j_{k-1}, j_k) \times \cdots \times \mathfrak{C}(\tilde{B}_0)(a, j_1) \rightarrow \mathfrak{C}(T)(a, b)$$

obtained by first including the \tilde{B}_i 's into T and then using the composition in $\mathfrak{C}(T)$.

Categorification of necklaces

Since the description of $\mathfrak{C}(\Delta^n)(-, -)$ is simple, we obtain a description for $\mathfrak{C}(T)(-, -)$:

Let $C_T(a, b)$ denote the poset whose elements are $V_T(a, b)$ which contain $J_T(a, b)$ (ordered by inclusion). The union of sets yields the following pairing:

$$C_T(b, c) \times C_T(a, b) \rightarrow C_T(a, c).$$

Applying the (regular) nerve functor, we obtain a simplicial category NC_T with object set T_0 . An n -simplex in $NC_T(a, b)$ can be viewed as a flag $\vec{T} := T^0 \subset \dots \subset T^n$, where $J_T(a, b) \subset T^0$ and $T^n \subset V_T(a, b)$.

Categorification of necklaces

Proposition 3

Let $T \in \text{Nec}$. There is a natural isomorphism of simplicial categories $\mathcal{C}(T)$ and NC_T .

Categorification of necklaces

Proposition 3

Let $T \in \text{Nec}$. There is a natural isomorphism of simplicial categories $\mathfrak{C}(T)$ and NC_T .

Idea of the proof:

- ▶ $\mathfrak{C}(T) = \mathfrak{C}(B_1) \sqcup \mathfrak{C}(B_2) \sqcup \cdots \sqcup \mathfrak{C}(B_k)$ and $\mathfrak{C}(B_i) = NC_{B_i}$.
- ▶ \vee factors through the functors $\mathfrak{C}(B_i) \rightarrow NC_{B_i} \rightarrow NC_T$ and thus yields a functor $f: \mathfrak{C}(T) \rightarrow NC_T$.
- ▶ This functor is the identity on object, thus one just needs to show that it is full and faithful.
- ▶ Build an inverse $g: NC_T(a, b) \rightarrow \mathfrak{C}(T)(a, b)$, i.e. take an n -simplex, split it into multiple simplices starting and ending in between joints $J_T(a, b)$, push them forward to $\mathfrak{C}(T)(a, b)$ and glue them back together there.

Categorification of necklaces

Corollary 4

Let $T = B_0 \vee \cdots \vee B_k$ be a necklace. Let $a, b \in T_0$ such that $a < b$. Let $j_r, j_{r+1}, \dots, j_s, j_{s+1}$ be the elements of $J_T(a, b)$ and B_i the bead containing j_i and j_{i+1} for $r \leq i \leq s$. Then the map

$$\mathfrak{C}(B_s)(j_s, j_{s+1}) \times \cdots \times \mathfrak{C}(B_r)(j_r, j_{r+1}) \rightarrow \mathfrak{C}(T)(a, b)$$

is an isomorphism. Therefore $\mathfrak{C}(T)(a, b) \cong (\Delta^1)^N$, with $N = |V_T(a, b)| - |J_T(a, b)|$. In particular, it is contractible if $a \leq b$ and empty otherwise.

The categorification functor

Fix $S \in \hat{\Delta}$ and $a, b \in S_0$. For any necklace T and map $T \rightarrow S_{a,b}$, there is an induced map $\mathfrak{C}(T)(\alpha_T, \omega_T) \rightarrow \mathfrak{C}(S)(a, b)$. We denote by $(Nec \downarrow S)_{a,b}$ the category whose objects are pairs $[T, T \rightarrow S_{a,b}]$ and whose morphisms are maps $T \rightarrow T'$ yielding a commutative triangle over S . Then we obtain a map

$$E_S(a, b) := \operatorname{colim}_{T \rightarrow S \in (Nec \downarrow S)_{a,b}} [\mathfrak{C}(T)(\alpha_T, \omega_T)] \rightarrow \mathfrak{C}(S)(a, b).$$

The categorification functor

Fix $S \in \hat{\Delta}$ and $a, b \in S_0$. For any necklace T and map $T \rightarrow S_{a,b}$, there is an induced map $\mathfrak{C}(T)(\alpha_T, \omega_T) \rightarrow \mathfrak{C}(S)(a, b)$. We denote by $(Nec \downarrow S)_{a,b}$ the category whose objects are pairs $[T, T \rightarrow S_{a,b}]$ and whose morphisms are maps $T \rightarrow T'$ yielding a commutative triangle over S . Then we obtain a map

$$E_S(a, b) := \operatorname{colim}_{T \rightarrow S \in (Nec \downarrow S)_{a,b}} [\mathfrak{C}(T)(\alpha_T, \omega_T)] \rightarrow \mathfrak{C}(S)(a, b).$$

We have a composition law on E_S :

$$E_S(b, c) \times E_S(a, b) \rightarrow E_S(a, c)$$

induced as follows: let $T \rightarrow S_{a,b}$ and $U \rightarrow S_{b,c}$, then $T \vee U \rightarrow S_{a,c}$. Thus, we have the composite

$$\begin{aligned} \mathfrak{C}(U)(\alpha_U, \omega_U) \times \mathfrak{C}(T)(\alpha_T, \omega_T) &\rightarrow \mathfrak{C}(T \vee U)(\omega_T, \omega_U) \times \mathfrak{C}(T \vee U)(\alpha_T, \omega_T) \\ &\rightarrow \mathfrak{C}(T \vee U)(\alpha_T, \omega_U) \end{aligned}$$

which induces the pairing.

The categorification functor

Proposition 5

For every simplicial set S , the map $E_S \rightarrow \mathfrak{C}(S)$ is an isomorphism of simplicial categories.

The categorification functor

Proposition 5

For every simplicial set S , the map $E_S \rightarrow \mathfrak{C}(S)$ is an isomorphism of simplicial categories.

Idea of the proof: playing around with colimits.

The categorification functor

Corollary 6

For any simplicial set S and elements $a, b \in S_0$, the simplicial set $\mathfrak{C}(S)(a, b)$ admits the following description. An n -simplex in $\mathfrak{C}(S)(a, b)$ consists of an equivalence class of triples

$[T, T \rightarrow S, \vec{T}]$, where

- ▶ $T \in \text{Nec}$;
- ▶ $T \rightarrow S$ is a map of simplicial sets, such that $\alpha_T \mapsto a$ and $\omega_T \mapsto b$;
- ▶ \vec{T} is a flag $T^0 \subset \dots \subset T^n$ such that $J_T \subset T^0$ and $T^n \subset V_T$.

The equivalence is generated by maps $f: T \rightarrow U$ over S such that $\vec{U} = f_*(\vec{T})$.

The i^{th} -boundary means deleting the i^{th} -set in the flag, the degeneracy means copying it.

The categorification functor

Corollary 6

For any simplicial set S and elements $a, b \in S_0$, the simplicial set $\mathfrak{C}(S)(a, b)$ admits the following description. An n -simplex in $\mathfrak{C}(S)(a, b)$ consists of an equivalence class of triples

$[T, T \rightarrow S, \vec{T}]$, where

- ▶ $T \in \text{Nec}$;
- ▶ $T \rightarrow S$ is a map of simplicial sets, such that $\alpha_T \mapsto a$ and $\omega_T \mapsto b$;
- ▶ \vec{T} is a flag $T^0 \subset \dots \subset T^n$ such that $J_T \subset T^0$ and $T^n \subset V_T$.

The equivalence is generated by maps $f: T \rightarrow U$ over S such that $\vec{U} = f_*(\vec{T})$.

The i^{th} -boundary means deleting the i^{th} -set in the flag, the degeneracy means copying it.

Proof: proven by considering colimits as coequaliser.

Flagged necklaces and totally nondegenerate maps

Flagged necklaces

A **flagged necklace** is a pair $[T, \vec{T}]$ where $T \in \text{Nec}$ and \vec{T} is a flag of subsets of V_T containing J_T . Its **length** is the cardinality of the set of subset symbols.

A flagged necklace is called **flanked** if $T^0 = J_T$ and $T^n = V_T$.

Flagged necklaces and totally nondegenerate maps

Flagged necklaces

A **flagged necklace** is a pair $[T, \vec{T}]$ where $T \in \text{Nec}$ and \vec{T} is a flag of subsets of V_T containing J_T . Its **length** is the cardinality of the set of subset symbols.

A flagged necklace is called **flanked** if $T^0 = J_T$ and $T^n = V_T$.

Lemma 7

Under the equivalence relation of corollary 6, each triple $[T, T \rightarrow S, \vec{T}]$ is equivalent to one in which the flag is flanked. Moreover, two flanked triples are equivalent if they can be connected by a zig-zag of morphisms of flagged necklaces in which every triple of the zig-zag is flanked.

Flagged necklaces and totally nondegenerate maps

Flagged necklaces

A **flagged necklace** is a pair $[T, \vec{T}]$ where $T \in \text{Nec}$ and \vec{T} is a flag of subsets of V_T containing J_T . Its **length** is the cardinality of the set of subset symbols.

A flagged necklace is called **flanked** if $T^0 = J_T$ and $T^n = V_T$.

Lemma 7

Under the equivalence relation of corollary 6, each triple $[T, T \rightarrow S, \vec{T}]$ is equivalent to one in which the flag is flanked. Moreover, two flanked triples are equivalent if they can be connected by a zig-zag of morphisms of flagged necklaces in which every triple of the zig-zag is flanked.

The proof of this lemma uses a functorial construction called **flankification**: let $[T, \vec{T}]$ be a flagged necklace, then there is a unique necklace $T' \hookrightarrow T$ whose set of joints is T^0 and vertex set is T^n .

Flagged necklaces and totally nondegenerate maps

Totally nondegenerate map

Let $T \in \text{Nec}$ and $S \in \hat{\Delta}$, then a map $T \rightarrow S$ is called **totally nondegenerate** if the image of each bead in T is a nondegenerate simplex of S .

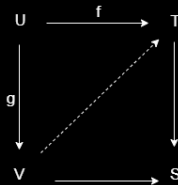
For any necklace T and map $T \rightarrow S$, there exists a necklace \bar{T} , a totally nondegenerate map $\bar{T} \rightarrow S$ and a surjection $T \rightarrow \bar{T}$ making the obvious triangle commute.

Flagged necklaces and totally nondegenerate maps

Proposition 8

Let $S \in \hat{\Delta}$ and $a, b \in S_0$.

- (a) Suppose that T and U are necklaces, $U \rightarrow S$ and $T \rightarrow S$ two maps and the second one totally nondegenerate. Then there is at most one surjection $U \rightarrow T$ making the obvious triangle commute.
- (b) Suppose one has



where T, U and V are flagged necklaces, $T \rightarrow S$ tot.nondegen. and f, g surjections. Then there is a unique map of flagged necklaces $V \rightarrow T$ making the diagram commute.

Flagged necklaces and totally nondegenerate maps

Corollary 9

Let $S \in \hat{\Delta}$ and $a, b \in S_0$. Under the equivalence relation of corollary 6, each triple $[T, T \rightarrow S_{a,b}, \overrightarrow{T}]$ is equivalent to a unique triple $[U, U \rightarrow S_{a,b}, \overrightarrow{U}]$ which is both flanked and totally nondegenerate.

Flagged necklaces and totally nondegenerate maps

Corollary 9

Let $S \in \hat{\Delta}$ and $a, b \in S_0$. Under the equivalence relation of corollary 6, each triple $[T, T \rightarrow S_{a,b}, \vec{T}]$ is equivalent to a unique triple $[U, U \rightarrow S_{a,b}, \vec{U}]$ which is both flanked and totally nondegenerate.

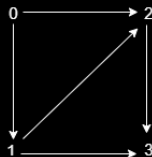
Idea of the proof: via the flankification and tot.nondegen replacement, we know at least one such triple exists. If $[U, U \rightarrow S_{a,b}, \vec{U}]$ and $[V, V \rightarrow S_{a,b}, \vec{V}]$ are both flanked and tot.nondegen. Then by lemma 7 we have a zig-zag of morphisms of flanked necklaces over S connecting U and V , using proposition 8 (b) on the zig-zag, we obtain two maps $U \rightarrow V$ and $V \rightarrow U$. By proposition 8 (a), these two maps must be inverses of each other.

An example

If $D \in \hat{\Delta}$ is ordered, then Dugger-Spivak show that to describe $\mathfrak{C}(D)(a, b)$, one can restrict oneself to flanked triples $[T, T \rightarrow S, \overrightarrow{T}]$ where $T \rightarrow S$ is an injection, i.e. T "lives" in D .

An example

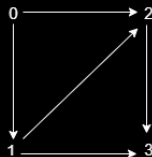
Let us take $S = \Delta^2 \sqcup_{\Delta^1} \Delta^2$:



We will compute $\mathfrak{C}(S)(0, 3)$. There are five necklaces injecting into $S_{0,3}$, namely $T_{1,2} = \Delta^1 \vee \Delta^1$, $U = \Delta^1 \vee \Delta^1 \vee \Delta^1$, $V = \Delta^1 \vee \Delta^2$ and $W = \Delta^2 \vee \Delta^1$.

An example

Let us take $S = \Delta^2 \sqcup_{\Delta^1} \Delta^2$:

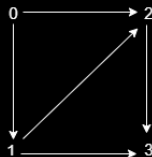


We will compute $\mathfrak{C}(S)(0, 3)$. There are five necklaces injecting into $S_{0,3}$, namely $T_{1,2} = \Delta^1 \vee \Delta^1$, $U = \Delta^1 \vee \Delta^1 \vee \Delta^1$, $V = \Delta^1 \vee \Delta^2$ and $W = \Delta^2 \vee \Delta^1$.

$\mathfrak{C}(S)(0, 3)_0 = \{[T_1, \{0, 1, 3\}], [T_2, \{0, 2, 3\}], [U, \{0, 1, 2, 3\}]\}$,
others are degeneracies.

An example

Let us take $S = \Delta^2 \sqcup_{\Delta^1} \Delta^2$:



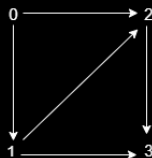
We will compute $\mathfrak{C}(S)(0, 3)$. There are five necklaces injecting into $S_{0,3}$, namely $T_{1,2} = \Delta^1 \vee \Delta^1$, $U = \Delta^1 \vee \Delta^1 \vee \Delta^1$, $V = \Delta^1 \vee \Delta^2$ and $W = \Delta^2 \vee \Delta^1$.

$\mathfrak{C}(S)(0, 3)_0 = \{[T_1, \{0, 1, 3\}], [T_2, \{0, 2, 3\}], [U, \{0, 1, 2, 3\}]\}$,
others are degeneracies.

$\mathfrak{C}(S)(0, 3)_1 = \{[V, \{0, 1, 3\} \subset \{0, 1, 2, 3\}], [W, \{0, 2, 3\} \subset \{0, 1, 2, 3\}]\}$.

An example

Let us take $S = \Delta^2 \sqcup_{\Delta^1} \Delta^2$:



We will compute $\mathfrak{C}(S)(0, 3)$. There are five necklaces injecting into $S_{0,3}$, namely $T_{1,2} = \Delta^1 \vee \Delta^1$, $U = \Delta^1 \vee \Delta^1 \vee \Delta^1$, $V = \Delta^1 \vee \Delta^2$ and $W = \Delta^2 \vee \Delta^1$.

$\mathfrak{C}(S)(0, 3)_0 = \{[T_1, \{0, 1, 3\}], [T_2, \{0, 2, 3\}], [U, \{0, 1, 2, 3\}]\}$,
others are degeneracies.

$\mathfrak{C}(S)(0, 3)_1 = \{[V, \{0, 1, 3\} \subset \{0, 1, 2, 3\}], [W, \{0, 2, 3\} \subset \{0, 1, 2, 3\}]\}$.

Thus $\mathfrak{C}(S)(0, 3)$ looks like $\bullet \leftarrow \bullet \rightarrow \bullet$.

Homotopy theory

We now have a nice description of the mapping spaces, but these are not particularly helpful to do homotopy theory, we will now introduce one type of categorification functor \mathcal{C}^{nec} which has better homotopic properties and another one \mathcal{C}^{hoc} which will enable us to create a zig-zag of weak equivalences between \mathcal{C} and \mathcal{C}^{nec} .

Homotopy theory

For $S \in \hat{\Delta}$, we treat a choice of $a, b \in S_0$ as a map $\partial\Delta^1 \rightarrow S$. We denote by $(Nec \downarrow S)_{a,b}$ the overcategory for the inclusion $Nec \hookrightarrow (\partial\Delta^1 \downarrow S)$ and define

$$\mathfrak{C}^{nec}(S)(a, b) := N(Nec \downarrow S)_{a,b}.$$

The other functor is given by the formula

$$\mathfrak{C}^{hoc}(S)(a, b) := \operatorname{hocolim}_{T \in (Nec \downarrow S)_{a,b}} \mathfrak{C}(T)(\alpha_T, \omega_T).$$

Homotopy theory

Theorem 10

For every simplicial set S , the maps $\mathcal{C}(S) \leftarrow \mathcal{C}^{hoc}(S) \rightarrow \mathcal{C}^{nec}(S)$ are weak equivalences of simplicial categories.

Meta-theory via gadgets

Let $\mathcal{P} \subset \hat{\Delta}_{*,*}$ containing the terminal object. For any $S \in \hat{\Delta}$ and $a, b \in S_0$, let $(\mathcal{P} \downarrow S)_{a,b}$ denote the overcategory whose objects are pairs $[P, P \rightarrow S]$, where $P \in \mathcal{P}$ and $P \rightarrow S$ such that $\alpha_P \mapsto a$ and $\omega_P \mapsto b$. We define an **assignment**:

$$\mathfrak{C}^{\mathcal{P}}(S)(a, b) := N(\mathcal{P} \downarrow S)_{a,b}.$$

It is functorial only if \mathcal{P} is "closed under wedges".

Meta-theory via gadgets

Let $\mathcal{P} \subset \hat{\Delta}_{*,*}$ containing the terminal object. For any $S \in \hat{\Delta}$ and $a, b \in S_0$, let $(\mathcal{P} \downarrow S)_{a,b}$ denote the overcategory whose objects are pairs $[P, P \rightarrow S]$, where $P \in \mathcal{P}$ and $P \rightarrow S$ such that $\alpha_P \mapsto a$ and $\omega_P \mapsto b$. We define an **assignment**:

$$\mathfrak{C}^{\mathcal{P}}(S)(a, b) := N(\mathcal{P} \downarrow S)_{a,b}.$$

It is functorial only if \mathcal{P} is "closed under wedges".

Category of gadgets

We call $\mathcal{G} \subset \hat{\Delta}_{*,*}$ a **category of gadgets** if it satisfies the following properties:

1. $Nec \subset \mathcal{G}$;
2. $\forall X \in \mathcal{G} \forall T \in Nec$, all maps $T \rightarrow X$ are in \mathcal{G} ;
3. $\forall X \in \mathcal{G}$, the simplicial set $\mathfrak{C}(X)(\alpha, \omega)$ is contractible.

The category \mathcal{G} is said to be **closed under wedges** if it is also true that: for any $X, Y \in \mathcal{G}$, the wedge $X \vee Y$ also belongs to \mathcal{G} .

Meta-theory via gadgets

Proposition 11

Let \mathcal{G} be a category of gadgets. Then for any simplicial set S and any $a, b \in S_0$, the natural map

$$\mathfrak{C}^{nec}(S)(a, b) \rightarrow \mathfrak{C}^{\mathcal{G}}(S)(a, b)$$

induced by the inclusion $Nec \hookrightarrow \mathcal{G}$, is a Kan equivalence. If \mathcal{G} is closed under wedges, then the map of simplicial categories $\mathfrak{C}^{nec}(S) \rightarrow \mathfrak{C}^{\mathcal{G}}(S)$ is a weak equivalence.

Meta-theory via gadgets

Proposition 11

Let \mathcal{G} be a category of gadgets. Then for any simplicial set S and any $a, b \in S_0$, the natural map

$$\mathfrak{C}^{nec}(S)(a, b) \rightarrow \mathfrak{C}^{\mathcal{G}}(S)(a, b)$$

induced by the inclusion $Nec \hookrightarrow \mathcal{G}$, is a Kan equivalence. If \mathcal{G} is closed under wedges, then the map of simplicial categories $\mathfrak{C}^{nec}(S) \rightarrow \mathfrak{C}^{\mathcal{G}}(S)$ is a weak equivalence.

The proof relies on [Quillen's theorem A](#), saying that if the overcategories of $Nec \hookrightarrow \mathcal{G}$ are contractible, the induced map is a Kan-equivalence. This is the case by the assumptions on \mathcal{G} and theorem 10.

Properties of categorification

Proposition 12

For any simplicial sets X and Y , both $\mathfrak{C}(X \times Y)$ and $\mathfrak{C}(X) \times \mathfrak{C}(Y)$ are simplicial categories with object set $X_0 \times Y_0$. For any $a_0, b_0 \in X_0$ and $a_1, b_1 \in Y_0$, the natural map

$$\mathfrak{C}(X \times Y)((a_0, a_1), (b_0, b_1)) \rightarrow \mathfrak{C}(X)(a_0, b_0) \times \mathfrak{C}(Y)(a_1, b_1),$$

induced by the projections, is a Kan equivalence. Therefore, the map of simplicial categories

$$\mathfrak{C}(X \times Y) \rightarrow \mathfrak{C}(X) \times \mathfrak{C}(Y)$$

is a weak equivalence in Cat_{Δ} .

Properties of categorification

Proposition 12

For any simplicial sets X and Y , both $\mathfrak{C}(X \times Y)$ and $\mathfrak{C}(X) \times \mathfrak{C}(Y)$ are simplicial categories with object set $X_0 \times Y_0$. For any $a_0, b_0 \in X_0$ and $a_1, b_1 \in Y_0$, the natural map

$$\mathfrak{C}(X \times Y)((a_0, a_1), (b_0, b_1)) \rightarrow \mathfrak{C}(X)(a_0, b_0) \times \mathfrak{C}(Y)(a_1, b_1),$$

induced by the projections, is a Kan equivalence. Therefore, the map of simplicial categories

$$\mathfrak{C}(X \times Y) \rightarrow \mathfrak{C}(X) \times \mathfrak{C}(Y)$$

is a weak equivalence in Cat_{Δ} .

Idea of the proof: take \mathcal{G} the category of objects containing finite product of necklaces, show that it is a category of gadgets. Then use theorem 10 and proposition 11 to work with $\mathfrak{C}^{\mathcal{G}}$ instead of \mathfrak{C} .

It has been a long battle...

It has been a long battle...

By now, most of you are probably as lost as the guy in the middle:



La Liberté guidant le peuple • Eugène Delacroix • 1830

But anyway...

Thank you for sticking with me until the end !
Any questions ?

Rigidification Of Quasi-Categories - Daniel Dugger and David Spivak - 5 October 2009 - arXiv:0910.0814

Lemma 2

Let X, Y be ordered simplicial sets and $f: X \rightarrow Y$ a map.

1. The category of ordered simplicial sets is closed under finite limits.
2. Every necklace is an ordered simplicial set.
3. If $X' \subset X$, then X' is ordered.
4. The map f is entirely determined by $f_0: X_0 \rightarrow Y_0$ on vertices.
5. If f_0 is injective, so is f .
6. The image of an n -simplex $x: \Delta^n \rightarrow X$ is of the form $\Delta^k \hookrightarrow X$ for some $k \leq n$.
7. If T is a necklace and $y: T \rightarrow X$ is map, then the image is a necklace.
8. Suppose that $X \leftarrow A \rightarrow Y$ is a diagram of ordered simplicial sets and both $A \rightarrow X$ and $A \rightarrow Y$ are simple inclusions. Then the pushout $B = X \sqcup_A Y$ is an ordered simplicial set and the inclusions $X \hookrightarrow B$ and $Y \hookrightarrow B$ are simple.