# Ant routing for the Lightning Network (with C. Grunspan)

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Building on Bitcoin

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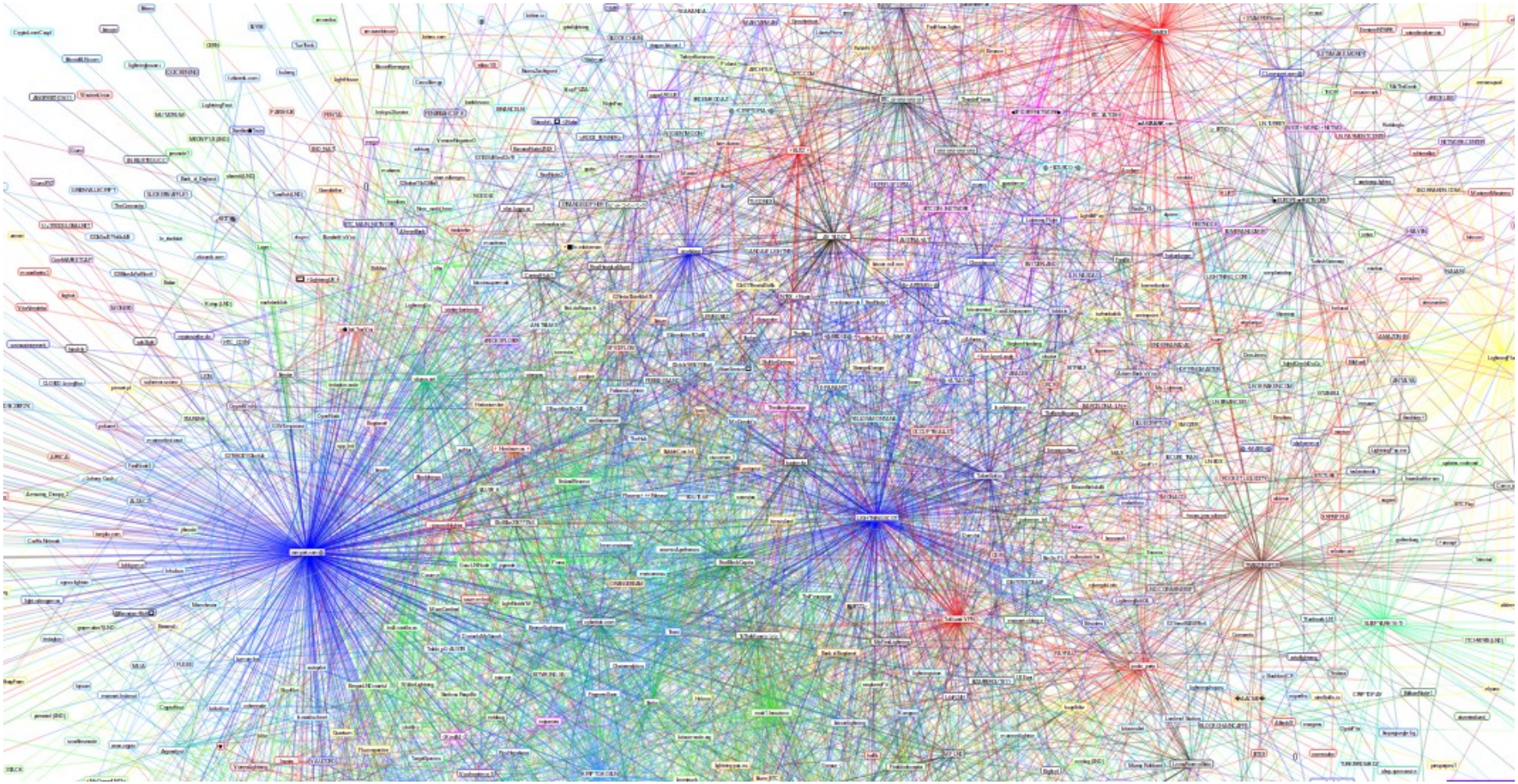

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- Nobody trusts anyone, everyone verifies. Vérificabilité

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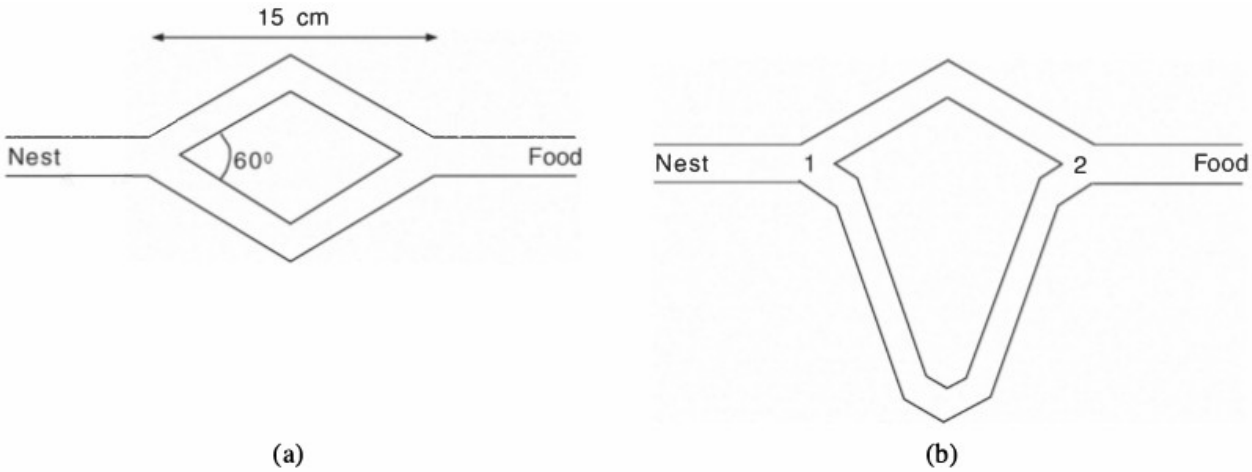
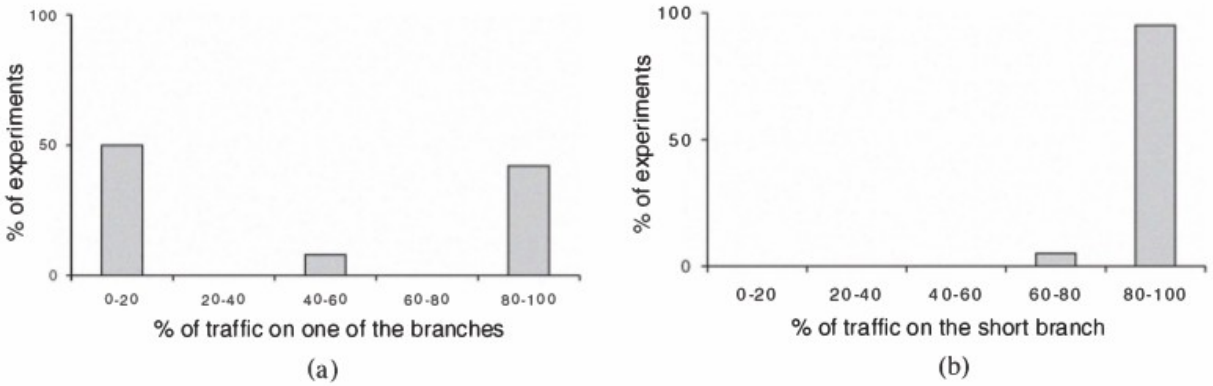



**Figure 1.1**
Experimental setup for the double bridge experiment. (a) Branches have equal length. (b) Branches have different length. Modified from Goss et al. (1989).



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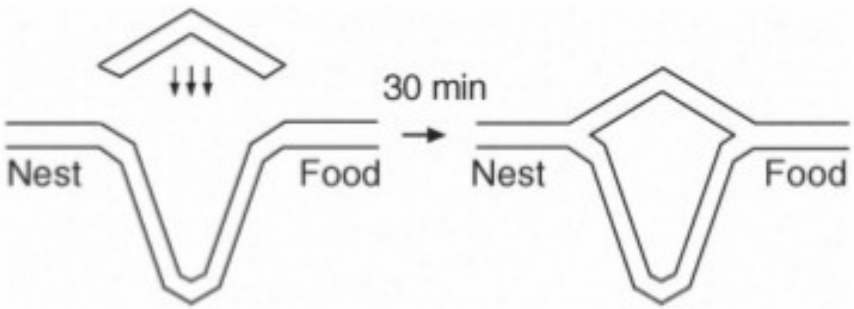


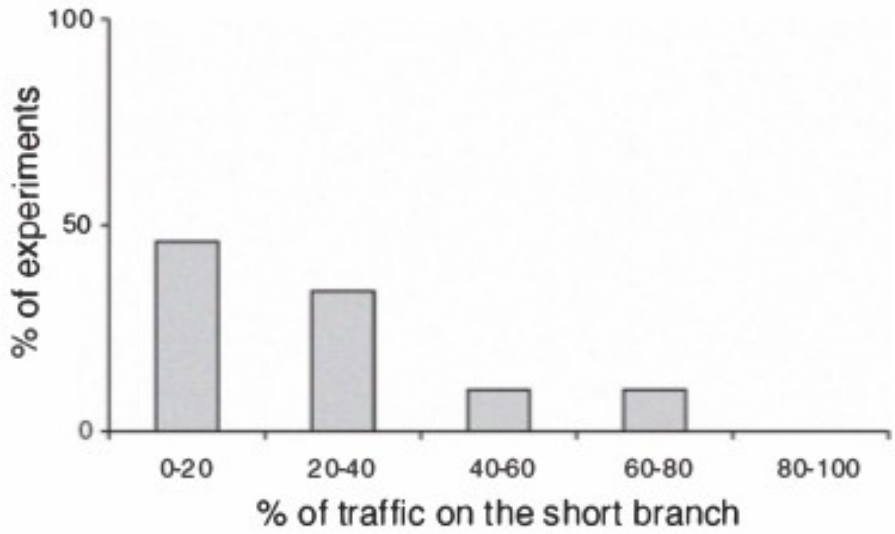

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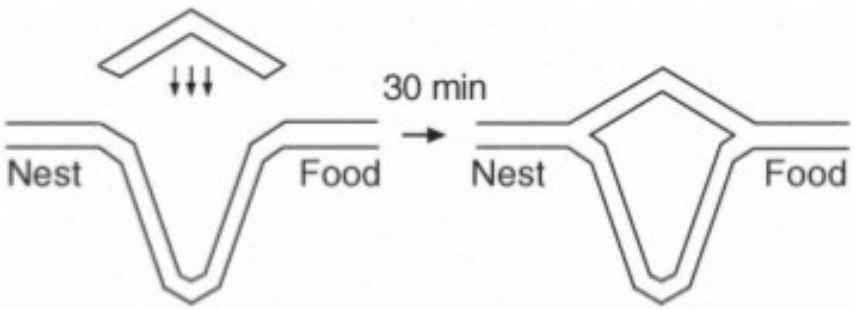


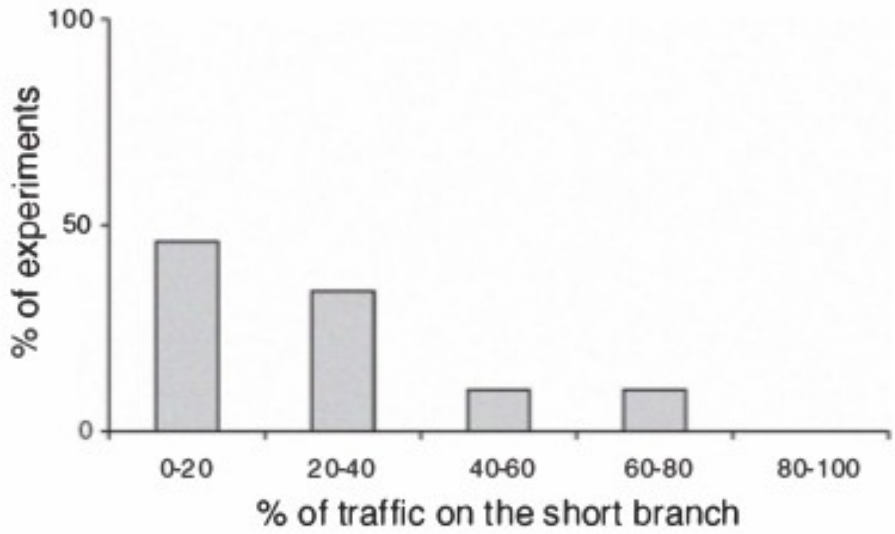


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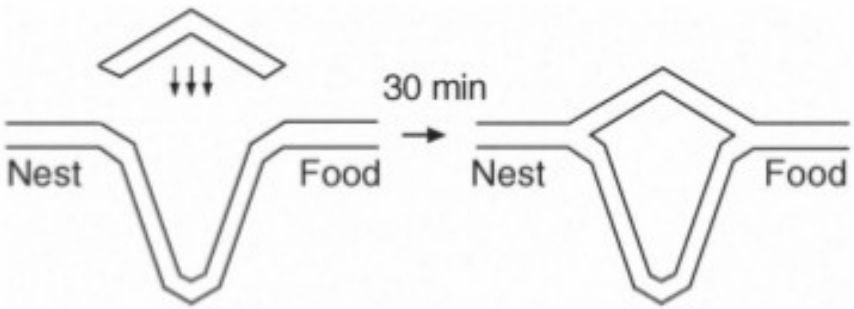


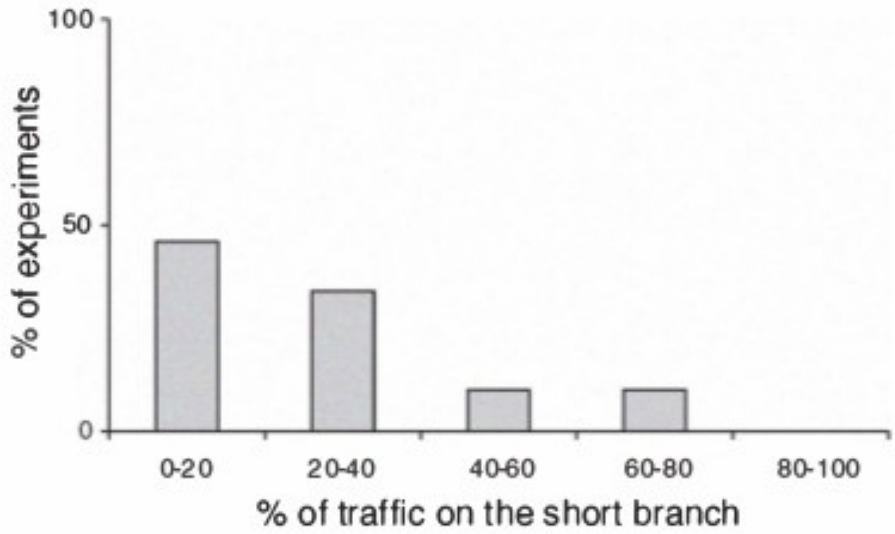


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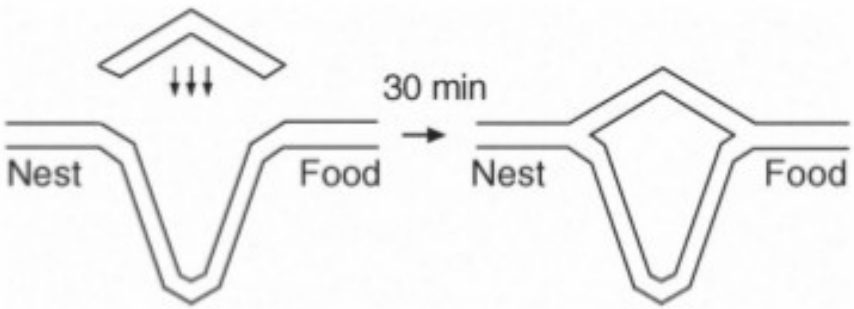


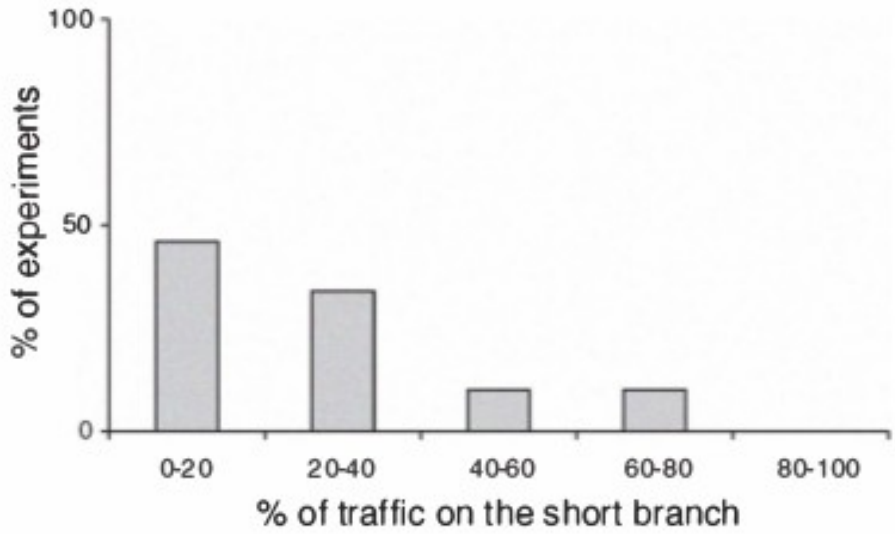


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# Self-improvement features

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# Thank for your attention!